

HYDRO-UNIVERSITY COMPUTING CENTRE

ALGOL PROCEDURES REFERENCE MANUAL

This document contains details of Algol procedures which have been written for and tested on the Elliott 503. Tapes for these procedures are available in the Computing Centre. Users wishing to incorporate these procedures in programs submitted for punching should indicate their requirements on the Algol program sheet thus:

TITLE:

begin integer; real.....;

LIBRARY LEO4

LIBRARY NCO1

comment then follows the program text;

⋮
⋮

end;

Such instructions will be sufficient to ensure that the appropriate procedures are copied at the places indicated. It will be left to programmers to ensure that the order of copying of requested procedures is a valid order of declaration. This is important in those cases where one procedure uses another as indicated in the procedure descriptions.

The manual also contains an index of HUCC Library procedures, together with an index of other known Algol procedures and their sources. In most cases copies of these are available in the Computing Centre. Procedure tapes, however, are available only for HUCC procedures.

CLASSIFICATIONS

C	COMMERCIAL
	CA INTEREST
D	INFORMATION HANDLING-DATA PROCESSING
	DH SORTING
	DP PLOTTING
	DZ MISCELLANEOUS
F	FUNCTIONS-EVALUATION OF-ETC
	FB BESSEL
	FC COMPLEX
	FD POWERS AND EXPONENTIALS
	FE ELLIPTIC INTEGRALS
	FP POLYNOMIALS-INC CHEBYSHEV ETC
	FS SPECIAL
	FT TRIGONOMETRICAL FUNCTIONS
	FZ MISCELLANEOUS
G	DIFFERENTIAL EQUATIONS
	GA ORDINARY-NOT LINEAR OR 1ST ORDER
	GL LINEAR
	GP PARTIAL
L	LINEAR ALGEBRA
	LA CHANGE FORM OF MATRIX
	LB BOOLEAN MATRICES
	LE LINEAR EQUATIONS AND INVERSION
	LF FORM SPECIAL MATRIX
	LG ARITHMETIC FUNCTIONS-ONE MATRIX
	LH ARITHMETIC FUNCTIONS-TWO MATRICES
	LL LATENT ROOTS
	LO DETERMINANTS
	LR READ OR INPUT
	LZ MISCELLANEOUS
M	MATHEMATICAL METHODS
	MC CURVE AND SURFACE FITTING
	MD DIFFERENTIATION-MAX AND MIN
	ME ERROR ANALYSIS
	MG GEOMETRY
	MH HARMONIC ANALYSIS
	MR ROOTS OF EQUATIONS
	MS INTEGRATION AND SUMMATION
	MT INTERPOLATION AND DIFFERENCES
	MZ MISCELLANEOUS
N	INTEGERS AND NUMBER THEORY
	NC PERMUTATIONS AND COMBINATIONS
	NP PRIME NUMBERS
	NR PARTITIONS
	NZ MISCELLANEOUS
O	OPERATIONAL RESEARCH
	OL LINEAR PROGRAMMING
	OP PERT-CRITICAL PATH ANALYSIS
P	PHYSICS
	PH HEAT
	PN NUCLEAR ENGINEERING
	PQ QUANTUM MECHANICS
S	STATISTICS
	SA ANALYSIS
	SB SMOOTHING
	SC CORRELATION
	SM MOMENTS
	SR REGRESSION
	SS STOCHASTIC PROCESSES
	SV ANALYSIS OF VARIANCE
Z	MISCELLANEOUS
	ZA COMPILER TECHNIQUES
	ZZ MISCELLANEOUS

ABBREVIATIONS

ACM COMMUNICATIONS OF ACM
AIS ALGOL INFORMATION SHEET
ALY ALGORYTMY
BIT NORDISK TIDSKRIFT FOR INF BEHANDLG
CJ COMPUTER JOURNAL
JCM JOURNAL OF THE A C M
MCA MATHEMATISCH CENTRUM AMSTERDAM
NUM NUMERISCHE MATHEMATIK
ORN OAK RIDGE NATIONAL LABORATORY
SUC STANFORD UNIVERSITY CALIFORNIA
BOO1 SELECTED NUMERICAL MTHS BY C. GRAM
NPL NATIONAL PHYSICAL LABORATORY

INDEX of HUCC ALGOL LIBRARY PROCEDURES.

DH 01 SORT REAL NUMBERS INTO ASCENDING ORDER
DH 02 LOCATE ELEMENT IN LIST
DH 03 SORT REAL NUMBERS IN ASCENDING ORDER
DH 04 SORT ROWS OR COLUMNS OF MATRIX

DP 01 OUTPUT GRAPH OF VECTOR ELEMENTS
DP 02 OUTPUT PLOT OF ELEMENTS OF TWO VECTORS

FC 01 COMPLEX ASSIGNMENT
FC 02 MULTIPLE COMPLEX ASSIGNMENT
FC 03 COMPLEX MULTIPLICATION
FC 04 COMPLEX DIVISION
FC 05 ASSIMILATE REAL IN COMPLEX OPERATION
FC 06 ASSIMILATE IMAGINARY IN COMPLEX OPERATION
FC 07 COMPLEX CONJUGATE
FC 08 MODULUS OF COMPLEX NUMBER
FC 09 ARGUMENT OF COMPLEX NUMBER
FC 10 POLAR FORM
FC 11 MULTIPLICATION BY IMAGINARY OPERATOR
FC 12 COMPLEX SQUARE
FC 13 COMPLEX RECIPROCAL
FC 14 COMPLEX ROOT
FC 15 COMPLEX LOGARITHM
FC 16 COMPLEX EXPONENTIAL
FC 17 HYPERBOLIC FUNCTIONS
FC 18 COMPLEX SINE
FC 19 COMPLEX COSINE
FC 20 COMPLEX TANGENT
FC 21 COMPLEX INVERSE SINE
FC 22 COMPLEX INVERSE COSINE
FC 23 COMPLEX INVERSE TANGENT
FC 24 COMPLEX POWER OF COMPLEX VARIABLE
FC 25 TEST FOR EVEN INTEGER
FC 26 LOGARITHM OF COMPLEX NUMBER

FP 01 SUM SERIES OF CHEBYSHEV POLYNOMIALS

FS 01 LOGARITHM OF FACTORIAL(n)

GL 01 FOURTH ORDER RUNGE-KUTTA

LE 01 MATRIX INVERSION
LE 02 SOLVE LINEAR EQUATIONS - ONE R.H. SIDE
LE 03 MATRIX DECOMPOSITION $Ax=b$
LE 04 FORWARD AND BACK SOLUTION $Ax=b$
LE 05 MATRIC DECOMPOSITION BY CROUTS METHOD
LE 06 TRIANGULAR MATRIX PRODUCT $LU=A$
LE 07 INVERT QUASI LOWER TRIANGULAR MATRIX
LE 08 INVERT QUASI UPPER TRIANGULAR MATRIX
LE 09 TRIANGULAR MATRIX PRODUCT $UL=A$
LE 10 procedure ATOLIU

LE 11	procedure LIUSOL
LE 12	procedure INV1
LE 13	procedure INVU
LE 14	procedure UL1TOA
LE 15	SOLVE TRIDIAGONAL LINEAR EQUATIONS
LE 16	DECOMPOSE SYMMETRIC POSITIVE DEFINITE MATRIX
LE 17	INVERT LOWER TRIANGULAR MATRIX IN SITU
LE 18	MATRIX MULTIPLICATION
LE 19	DECOMPOSE BANDMATRIX INTO LOWER AND UPPER TRIANGLES
LE 20	FORWARD SOLUTION AND BACK SUBSTITUTION FOR BAND EQUATION $Ax=b$
LE 21	SOLVE BAND EQUATIONS $Ax=b$
LE 22	REARRANGE PERMUTED TRIANGULAR MATRIX
LE 23	procedure MOVL
LE 24	DECOMPOSE SYMMETRIC POSITIVE DEFINITE MATRIX
LE 25	FORWARD AND BACK SOLUTION $Ax=b$, SYMMETRIC A
LE 26	DECOMPOSE BAND MATRIX WITH PIVOTING
LE 27	FORWARD AND BACK SOLUTION $Ax=b$, BANDMATRICES
LE 28	DECOMPOSE SYMMETRIC POSITIVE DEFINITE MATRIX
LE 29	FORWARD AND BACKWARD SOLUTION OF $Ax=b$, SYMMETRIC
LL 01	EIGENVALUES AND EIGENVECTORS OF SYMMETRIC MATRIX
LL 02	SOLUTION OF THE EIGENPROBLEM $(A-\lambda B)x=0$
LZ 01	GENERATE UNSYMMETRIC TES MATRIX
MC 01	LEAST SQUARES POLYNOMIAL FIT
MD 01	LOCATE MINIMUM OF FUNCTION $f(x)$
MG 01	CONVERT SEXAGESIMAL ANGLES TO RADIANS
MG 02	CONVERT RADIAN ANGULAR MEASURE TO SEXAGESIMAL UNITS
MG 03	CONVERT CENTESIMAL ANGLES TO RADIANS
MG 04	CONVERT RADIAN ANGULAR MEASURE TO CENTESIMAL UNITS
MG 05	COMPUTE DISTANCE AND GRID BEARING
MG 06	COMPUTE POLAR CO-ORDINATES OF POINT X, Y
MG 07	COMPUTE POINT OF INTERSECTION FROM ANGLES
MG 08	COMPUTE POINT OF INTERSECTION FROM BEARINGS
MG 09	THREE POINT RESECTION BY ANGLES
MG 10	MULTIPLE POINT RESECTION FROM ANGLES
MG 11	THREE POINT RESECTION FROM DISTANCES
MG 12	MULTIPLE POINT RESECTION USING DISTANCES
MG 13	THE INACCESSIBLE BASE PROBLEM
MH 01	COMPUTE FOURIER COEFFICIENTS
MH 02	SUM FOURIER SERIES
MR 01	LOCATE ROOT OF $f(x)=0$
MS 01	GENERAL SUM SERIES
MS 02	ADAPTIVE SIMPSON INTEGRATION
MS 03	EVALUATE DEFINITE INTEGRAL
MT 01	LAGRANGE INTERPOLATION
MT 02	AITKEN UNEQUAL INTERVAL INTERPOLATION
MT 03	AITKEN EQUAL INTERPOLATION
MT 04	AITKEN Nth ORDER INTERPOLATION
MT 05	ESTIMATION OF DERIVATIVE OF GIVEN FUNCTION - UNEQUAL INTERVALS
MT 06	ESTIMATION OF DERIVATIVE OF GIVEN FUNCTION EQUAL INTERVALS

NC 01 REARRANGE ELEMENTS OF VECTOR
NC 02 PERMUTE ROWS OR COLUMNS OF MATRIX
NC 03 PERMUTATION OF ELEMENTS OF A VECTOR
NC 04 PERMUTE ROWS OR COLUMNS OF MATRIX
NC 05 PRE OR POST MULTIPLY MATRIX BY PERMUTED IDENTITY MATRIX
NC 06 INVERSEPERMUAION OF INTEGER VECTOR

OP 01 CRITICAL PATH SCHEDULING

SS 01 REAL RANDOM NUMBER GENERATOR
SS 02 REAL RANDOM NUMBER GENERATOR
SS 03 RANDOM NUMBER GENERATOR - POISSON DISTRIBUTION
SS 04 RANDOM NUMBER GENERATOR - STANDARD NORMAL DISTRIBUTION
SS 05 GENERATE LARGE INTEGER RANDOM NUMBERS
SS 06 GENERATE LARGE INTEGER RANDOM NUMBERS
SS 07 GENERATE SMALL INTEGER RANDOM NUMBERS
SS 08 GENERATE SMALL INTEGER RANDOM NUMBERS

ZM 01 OUTPUT BINARY VALUE
ZM 02 OUTPUT OCTAL VALUE
ZM 03 SENSE NUMBER GENERATOR KEY
ZM 04 INPUT NEWLINE STRINGS

ZP 02 PRINT MATRIX
ZP 03 CONTROL FLEXOWRITER PAGE PRINTING

TITLE	REFERENCES
PERIODIC PAYMENT AS FN OF UNPAID PRINC	CA ACM45 4-61 9-63
SORTING-MATHSORT	DH ACM23 11-60 5-61
PARTITION-USED BY QUICKSORT ACM64	ACM63 7-61 8-62 8-63
QUICKSORT	ACM64 7-61 8-62 8-63
FIND	ACM65 7-61 8-62 8-63
SORTING PROCEDURES	ACM76 1-62 6-62
TREESORT-SUPERSEDED BY TREESORT 1 ACM143	ACM113 8-62
TREESORT 1	ACM143 12-62
TREESORT 2	ACM144 12-62
SHUTTLE SORT-EXCHANGING PAIRS	ACM175 6-63 10-63 12-63 5-64
SHELLSORT-SHELLS METHOD	ACM201 8-63 6-64
STRINGSORT	ACM207 10-63 10-64
RADIX EXCHANGE ETC	ACM 5-63 P209
SORTING BY DISTRIBUTION COUNTING	ORNL3 148
HEAP SORT	ACM 23 2 6-64
TREESORT 3	ACM245 12-64
SEARCH IN A LIST	JACM 1962(23)
INSERTION IN A LIST	JACM 1962(23)
DELETION FROM A LIST	JACM 1962(24)
SORTING WITH MIN STORAGE	JACM 1962(27)
	DP
COMPUTE CODE STRING TO MOVE PLOTTING PEN	ACM162 4-63 8-63 8-64
PROCEDURES FOR GRAPH PLOTTING	NPL 0003
	DZ
ADD ITEM TO CHAIN LINKED LIST	ACM100 6-62
REMOVE ITEM FROM CHAIN-LINKED LIST	ACM101 6-62
ENLARGEMENT OF A GROUP	ACM136 11-62
LOCATE VECTOR IN LEXICOGR ORDERED LIST	ACM151 2-63
	FB
BESSEL FN 1-SERIES EXPANSION	ACM5 4-60
BESSEL FN 1-ASYMPTOTIC EXPANSION	ACM6 4-60
BESSEL FN FOR SET OF INTEGER ORDERS	ACM21 11-60
RICCATI-BESSEL FNS OF 1ST AND 2ND KIND	ACM22 11-60
BESSEL FUNCTIONS BY RECURSION	ACM44 4-61
EVALUATE BER OR BEI FN	ACM57 4-61 7-62 8-62
Q BESSEL FUNCTIONS	ACM214 11-63 6-64
BESSEL FUNCTIONS-CHEBYSHEV-CLENSHAW	AIS13
BESSEL FCNS 1 ST KIND	ACM236 8-64
Q BESSEL FCN	ACM228 5-64
	FC
COMPLEX EXPONENTIAL INTEGRAL	ACM13 7-60 2-61
EXPONENTIAL OF COMPLEX NUMBER	ACM46 4-61 6-62
LOG OF COMPLEX NUMBER	ACM48 4-61 6-62 7-62 8-64
NTH ROOTS OF A COMPLEX NUMBER	ACM53 4-61 7-61
COMPLEX NUMBER TO A REAL POWER	ACM106 7-62 11-62
COMPLEX DIVISION	ECM116 8-62
COMPLEX ARITHMETIC	ACM186 7-63
COMPLEX POWER	ACM190 7-63
HYPERGEOMETRIC FN-COMPLEX PARAMETERS	ACM191 7-63 4-64
CONFLUENT HYPERGEOM FN-COMPLEX PARAMS	ACM192 7-63 4-64
FREQUENCY RESPONSE	AIS16
ARSENAL OF COMPLEX ALGOL PROCEDURES	BIT 2 4 P237
COMPLEX SQUARE ROOT	MCA X1 AP216
ZERO OF COMPLEX FUNCTION	MCA X1 AP217
MULTIPLY COMPLEX VALUES	MCA X1 AP218
COMPLEX PART OF POLYNOMIAL	MCA X1 AP219
LOG OF COMPLEX NUMBER	ACM243 11-64

GENERAL ORDER ARITHMETIC
EXPONENTIATION OF SERIES -COEFFS
EXPONENTIATION OF SERIES-COEFFS

COMPLETE ELLIPTIC INTEGRAL 1ST KIND
COMPLETE ELLIPTIC INTEGRAL 2ND KIND
INCOMPLETE ELLIPTIC INTEGRALS
COMPLETE ELLIPTIC INTEGRAL
COMPLETE ELLIPTIC INTEGRALS
PERIOD OF ELLIPTIC FUNCTION
ARSENAL ELLIPTIC INTEGRALS AND FUNCTIONS

CHEBYSHEV POLYNOMIAL BY RECURSION
HERMITE POLYNOMIAL BY RECURSION
LAGUERRE POLYNOMIAL BY RECURSION
LEGENDRE POLYNOMIAL BY RECURSION
POLYNML ORIGIN AND AXIS TRANSFORMER
EVALUATE TABLE OF CHEBYSHEV POLYNMLS
REDUCE DEGREE OF APPROXG POL-TELESCOPE 1
REDUCE DEGREE OF APPROXG POL-TELESCOPE 2
COEFFS OF QUOTIENT OF 2 POWER SERIES
REVERSION OF SERIES
CHEBYSHEV COEFFS
CHEBYSHEV POLY COEFF
POISSON-CHARLIER POLY

COMPLEX EXPONENTIAL INTEGRAL
GAMMA FUNCTION
GAMMA FUNCTION
ASSOC LEGENDRE FNS 1ST KIND-REAL OR IMAG
SPHERICAL NEUMANN FUNCTION
GAMMA FN-RANGE 1 TO 2
SET ASSOC LEGENDRE POLYNMLS 2ND KIND
RECIPRCCAL GAMMA FN-REAL ARGUMENT
JACOBI SYMBOL BY QUADRATIC RECIPROCITY
REAL ERROR FUNCTION
HANKEL FUNCTION 1ST KIND
PSIF-LOG DERIV OF FACTORIAL FUNCTION
MODIFIED HANKEL FUNCTION
INCOMPLETE BETA RATIO
ERROR FUNCTION-LARGE X
COMPLEMENTARY ERROR FN-LARGE X
HYPERGEOMETRIC FN-COMPLEX PARAMETERS
CONFLUENT HYPERGEOFN FN-COMPLEX PARAMS
CALENDAR DATE TO AND FROM JULIAN DAY
GAUSS FUNCTION
FRESNEL INTEGRALS
CALCULATION OF EASTER
COMPUTING OF ROMAN FUNCTION
CALCN OF NORMAL DISTRIBUTION FUNCTION
COMPLEMENTARY FRESNEL INTEGRALS
COMPUTATION OF THE FERMI FUNCTION
CONTD FRCTN EXPNSN FOR BINML QUADTC SURDS
EXPONENTIAL INTEGRAL BY EPSILON ALGTHM
GRAMMA FCN
INCOMPLETE BETA FCN
GAMMA FCN CONTROLLED ACCY
NORMAL DIST FCN

FD
ACM93 6-62 10-62
ACM134 11-62 7-63
ACM158 3-63 7-63 9-63

FE
ACM55 4-61 4-63
ACM56 4-61
ACM73 12-61 10-62 2-63
ACM149 12-62 4-63
ACM165 4-63
NUM 4-5 P396
NUM 5-4 P296

FP
ACM10 6-60 4-61
ACM11 6-60
ACM12 6-60
ACM13 6-60 4-61
ACM29 11-60
ACM36 3-61
ACM37 3-61 8-62 8-63
ACM38 3-61 8-63
ACM131 11-62
ACM193 7-63 12-63
ACM 2-62 P93
ACM227 5-64
ACM234 7-64

FS
ACM13 7-60 2-61
ACM31 2-61 12-62 1-63
ACM34 2-61 7-62
ACM47 4-61 8-63
ACM49 4-61
ACM54 4-61
ACM62 7-61 12-61
ACM80 3-62
ACM99 6-62 11-62
ACM123 9-62 6-63 10-63 3-64
ACM124 9-62
ACM147 12-62 4-63
ACM163 4-63 9-63
ACM179 6-63
ACM180 6-63
ACM181 6-63
ACM191 7-63
ACM192 7-63
ACM199 8-63
ACM209 10-63 3-64
ACM213 10-63
ACM4-62 P209 11-62 P556
ALY 1 1 P25
ALY 1 1 P33
BIT 2 3 P192
BIT 3 2 P141
NUM 5-2 P115
MCA MR60
ACM 221 3-64 10-64
ACM 222 3-64 4-64
ACM 225 5-64 10-64
ACM226 5-64

ARCCOSSIN-ARCCOS AND ARCSIN	FT
TAN-LAMBERTS CONTINUED FRACTION	ACM206 9-63
ARCTAN ALD	ACM 2-62 P89
ARCTAN ALS	MCA X1 AP116
SELECTIVE SUMMATION OF FOURIER SERIES	MCA X1 AP117
ELEMENTARY FCNS CONT. FRAC.	CJ 6-3 P248
ARCTAN(Z)	ACM229 5-64
	ACM241 9-64
PRODUCT OF TERMS OF A FN-CONTINUED PROD	FZ
ALGORITHMS FOR SPECIAL FUNCTIONS 1	MCA X1 AP202
	NUM 4-5 P409
SECOND ORDER EQN-ADAMS METHOD	GA
INTEGRATION OF ORDINARY DIFFERENTL EQNS	ACM 2-62 P88
	MCA R743
RUNGE KUTTA	GL
ZERSOL-ZEROS OF SOLN OF 1ST EQUNS	ACM9 5/60
KUTTA MERSON	ACM194 8-63
FREQUENCY RESPONSE	ACM218 12-63 10-64
RUNGE KUTTA	AIS16
	NUM 2 2 P131
MULTIREGION TWO GROUP PROGRAMS	GP
USE OF METHOD OF KERNEL FUNCTION	AIS14
SOLVE 5 PT APPRXN DIRICHLET PROBLEM	NUM 3 3 P219
SOLVE 9 PT APPRXN DIRICHLET PROBLEM	B0001 P49
SOLVE 5 PT APPRXN DIRICHLET PROBLEM 2	B0001 P49
INTEGRATE HEAT EQUATION	B0001 P55
CANONICAL COEFFICIENTS FOR DIRICHLET PRO	B0001 P83
	NPL0001
REDUCTION OF ARRAY TO JACOBI FORM	LA
REAL SYMM MATRIX TO TRIDIAG FORM	ACM104 7-62
ORTHONORMALISING PROCEDURE	ACM122 9-62 3-64
SYMM BANDMATRIX TO TRIDIAG FORM	ACM127 10-62
ORTHONORMALISE COLUMNS	ACM183 6-63
REDUCTION OF A MATRIX TO CODIAGONAL FORM	ACM 2-62 P91
TRANSFORM MATRIX SYMMETRIC TO TRIDIAG	CJ 4 P175
	MCA X1 AP210
ANCESTOR	LB
SHORTEST PATH	ACM96 6-62 3-63
PATH MATRIX	ACM97 6-62
	ACM141 11-62

SIM EQNS-CROUT WITH PIVOTING	LE
TRIDIAGONAL SIM EQNS	ACM16 9-60 10-60 3-61
TRIDIAGONAL LINEAR EQUATIONS	ACM17 9-60
INVERT MATRIX	ACM24 11-60
SIM EQNS-CROUT WITH PIVOTING 2	ACM42 4-61 11-61 1-63 8-63
INVERSE OF FINITE SEG OF HILBERT MATRIX	ACM43 4-61 8-63
ADJUST INVERSE OF PERTURBED MATRIX	ACM50 4-61 1-62 1-63
MATRIX INVERSION-GAUSS-JORDAN	ACM51 4-61 7-62
INVERSION-SQ RT METHOD	ACM58 5-61 6-62 8-62 12-62
SIM EQUNS AND INVERSION - BY ITERATION	ACM66 7-61 1-62 6-62
SIM EQNS-GAUSS METHOD	ACM92 5-62
MATRIX INVERSION 2-GAUSS JORDAN	ACM107 7-62 1-63 8-63
SIM EQNS-GAUSS METHOD	ACM120 8-62 1-63 8-63
CROUT WITH EQUILIBRATION AND ITERATION	ACM126 10-62
MATRIX INVERSION	ACM135 11-62 7-64
INVERSE OF SYMMETRIC MATRIX 2	ACM140 11-62 8-63
1 ROW OF INVERSE MATRIX-MONTE CARLO	ACM150 2-63 7-63 3-64
BANDSOLVE-BAND EQUATIONS	ACM166 4-63 9-63
GAUSS SEIDEL	ACM195 8-63
INVERSION-GAUSS JORDAN	ACM220 12-63 6-64
BAND MATRICES	ACM 2-62 P94
GAUSSIAN SOLUTION OF LINEAR EQUATIONS	AIS11
SOLVE LINEAR EQUATIONS	BIT 2 4 P256
INVERT	MCA X1 AP205
DETERMINANT AND INVERSE	MCA X1 AP206
SOLVE USING AP220 EVAL OF DETERMINANT	MCA X1 AP208
INVERT USING AP220	MCA X1 AP221
INVERT USING AP220 2	MCA X1 AP222
CONJUGATE GRADIENT METHOD	NUM5-2 P194
SOLVE USING AP224	NUM MR63 AP225
INVERT USING AP224	MCA MR63 AP225
INVERT USING AP224 2	MCA MR63 AP226
SOLVE USING AP228	MCA MR63 AP227
GAUSSIAN ELIMINATION	MCA MR63 AP229
SOLVE BY OVERRELAXATION	B0001 P25
SOLVE BANDMATRIX	B0001 P26
SOLVE DETERMINANT INVERT SYMMETRIC	BIT 3-3 P207
SOLVE DETERMINANT SYMMETRIC BAND	NPL 0004
MATRIX PERMUTATION	NPL 0005
MATRIX INVERSION COMP. DIV.	ACM230 6-64
CONT. GRAD. METHOD	ACM231 6-64
	ACM238 8-64
FORM SET OF TEST MATRICES	LF
FORM SET OF TEST MATIECES	ACM52 4-61 8-61 11-61 8-62
	ACM52 1-63 8-63
SYMMETRIC MATRICES-MATRIX SCHEME	LG
MATRIX DIVISION-SQUARE ROOT METHOD	AIS21
TERM BY TERM ARITHUETRIC PROCEDURES	
SYMMETRIC MATRICES -MATRIX SCHEME	
MATRIX EIGEN VALUES AND VECTORS-JACOBI	LH
PREPARE FOR GIVENS METHOD AP212	ACM197 8-63 3-64
GIVENS METHOD	AIS20
COMPUTE EIGENVECTOR	AIS21
TRANSFORM EIGENVECTOR	
EIGENVALUES AND VECTORS OF SYM MATRIX	LL
HOUSEHOLDERS METHOD FOR SYMMETRIC MATRIX	ACM85 4-62 8-62 8-63
EVALUES OF SYM TRIDIAG MAT BY BISECTION	MCA X1 AP211
EVECS OF SYM TRIDIAG MAT BY INV ITERATN	MCA X1 AP212
EVALUES OF SYM BAND MATRIX BY CHOLESKI	MCA X1 AP213
EIGENVALUES AB MINUS K1-SYMMETRIC	MCA X1 AP214
EIGENVALUES A MINUS KB-SYMMETRIC	MCA X1 AP215

DETERMINANT BY TRIANGULARISATION
DETERMINANT-COMBINATORIAL METHOD
REDUCTION OF MATRIX WITH POLYNML ELEMS
DETERMINANT OF SYMMETRIC POSITIVE MATRIX
FORM DETERMINANT
DETERMINANT AND SOLUTION
DETERMINANT AND SOLUTION
DETERMINANT AND INVERSE
DETERMINANT OF A BAND MATRIX
DETERMINANT OF SYMMETRIC POSITIVE MATRIX
DETERMINANT OF SYMMETRIC POSITIVE MATX 2
EVALUATION OF DETERMINANT

CRAM-READ SQUARE SYMM MATRIX AS VECTOR

SYMMETRIC MATRICES-MATRIX SCHEME

LEAST SQUARE FIT BY ORTHOGONAL POLYNOMIALS
CURVE FITTING WITH CONSTRAINTS
CHEBYSHEV CURVE FIT
SURFACE FIT -ORTHOGL POLYNOMIAL LST SQS
LEAST SQUARES SURFACE FIT
LEAST SQUARES SOLUTION WITH CONSTRAINTS
ERLANG PROBABILITY FOR CURVE FITTING
NORMAL PROBABILITY FOR CURVE FITTING
LEAST SQUARES SOLUTION
SPLINE CURVE

INTERPGLATN-DIFFERENTIATN AND INTEGRATN
MIN OF FN OF N VARABLES-STEEPEST DESCENT
DIRECT SEARCH-MIN OF FN OF N VARIABLES
DIFFERENCES AND DERIVATIVES
STEEP 1-MIN BY STEEPEST DESCENT
STEEP 2-MIN STEEPEST DESCENT
ATIVE-USED BY ACM 203
MAXIMUM OF A FUNCTION

PROCEDURES FOR RANGE ARITHMETIC
CALCN OF NORMAL DISTRIBUTION FUNCTION
ALGOL PROCEDURES FOR RANGE ARITHMETIC

POSITION OF POINT WRT POLYGON
COMPUTE CODE STRING TO MOVE PLOTTING PEN
PERSPECTIVE DRAWING PROGRAM-MARK 3

SUMMATION OF FOURIER SERIES
FOURIER SERIES APPROXIMATION
EPSILON ALGORITHM
EPSILON ALGORITHM

LQ
ACM41 4-61 9-63 3-64
ACM159 3-63 12-63
ACM170 4-63 8-63 7-64
MCA MR 63 AP224
MCA X1 AP 204
MCA X1 AP207
MCA X1 AP207
MCA X1 AP208
MCA X1 AP209
MCA X1 AP220
MCA MR63 AP228
ACM224 4-64 12-64

LR
ACM67 7-61 6-62

LZ
AIS21

MC
ACM28 11-60 12-61
ACM74 1-62 6-63
ACM91 5-62 4-63 5-64
ACM164 4-63 8-63
ACM176 6-63
ACM177 6-63 7-63
ACM184 7-63
ACM185 7-63
ACM 2-62 P92
BIT 2 2 P81

MD
ACM77 2-62 6-62 8-63 11-63
ACM129 11-62 9-63
ACM178 6-63
ACM187 7-63
ACM203 9-63 10-63
ACM204 9-63
ACM205 9-63
MCA X1 AP201

ME
ACM61 7-61
ALY 1 1 P83
SUC 15

MG
ACM112 8-62 12-62
ACM162 4-63 8-63
AIS4

MH
ACM128 10-62 7-64
ACM157 3-63 9-63 10-63
11-63 5-64
NUM MATH V6

ROOTFINDER-FAST ITERATION METHOD
ROOTS OF POLY-BAIRSTOW-HITCHCOCK METHOD
ROOTS BY ITERATED BISECTION
ROOTFINDER 2-ITERATION
REAL ZEROS OF AN ARBITRARY FUNCTION
ROOTFINDER 3
ROOTS OF POLY-NEWTON BAIRSTOW ETC
ZEROS REAL POLYNML BY RESULTNT PROCEDURE
FACTORS OF POLYNOMIALS
ROOTS OF POLYNOMIALS-INT-COEFFS
ZEROS OF POLYNOMIALS-NEWTON MAEHLY
BOUNCS ON A ZERO OF A POLYNOMIAL
ROOTS OF ARBITRARY FN-MULLERS METHOD
ROOTFINDERS-LINEAR AND QUADRATIC
ZERO OF COMPLEX FUNCTION
GRAIFEE PROCESS
ZERO OF A FUNCTION
SMALL INTEGERS AS PRIMITIVE ROOTS
QUOTIENT DIFFERENCE METHOD
ZERO OF AN ARBITRARY FUNCTION

MR
ACM2 2-60 6-60 8-60 3-61
ACM3 2-60 6-60 2-61 3-61
ACM4 3-60 3-61
ACM15 8-60 11-60 3-61
ACM25 11-60 3-61
ACM26 11-60 3-61
ACM30 12-60 5-61 1-62
ACM59 5-61
ACM75 1-62 7-62 8-62
ACM78 2-62 3-62 8-62
ACM105 7-62 7-63
ACM174 6-63
ACM196 8-63
AIS15
MCA X1 AP217
JCM 7 P93 83
MCA X1 AP200
NUM 4-4 P338
NUM 5-2 P96
BIT 3-2 P205

QUADRATURE INT-SEVERAL FNS-SAME LIMITS
EULER SUMMATION OF SERIES
REAL EXPONENTIAL INTEGRAL
MYLTIPLE INTEGRATION-GAUSS
ROMBERG INTEGRATION
INTERPOLATN-DIFFERENTIATN AND INTEGRATN
SIMPSONS INTEGRATION
FRESNEL SIN AND COS INTEGRALS-ASYMP EXPN
FRESNEL SIN INTEGRAL
FRESNEL COS INTEGRAL
DEFINITE COMPLEX LINE INTEGRALS
SIMPSONS RULE INTEGRATOR
DEFINITE EXPONENTIAL INTEGRALS A
DEFINITE EXPONENTIAL INTEGRALS B
WEIGHT COEFFS FOR GAUSS QUADRATURE
ADAPTIVE INTN BY SIMPSONS RULE
MULTIPLE INTEGRATION
NONRECURSIVE ADAPTIVE INTEGRATION
ADAPTIVE AND MULTIPLE INTEGRATION
GAUSS FUNCTION
FRESNEL INTEGRALS
COMPLEMENTARY FRESNEL INTEGRALS
THE LIMIT OF A CONVERGING SEQUENCE
SIMPSON INTEGRATION VARIABLE H
INTEGRATION BY SIMPSONS RULE
SIMPSONS RULE IN ALGOL 58
EULER-SUMMATION OF SERIES
TRANSFORMATION OF SLOWLY CONVERGR SERIES
MULTIPLE INT. SIMPS. RULE
ROMBERG INTEGRATION
FRESNEL INTEGRAL

MS
ACM 2-60
ACM8 5-60 11-63
ACM20 10-60 2-61 4-61
ACM32 2-61 2-63
ACM60 6-61 3-62 5-62 7-64
ACM77 2-62 6-62 8-63 11-63
ACM84 4-62 7-62 8-62 11-62
ACM88 5-62 10-63
ACM89 5-62 10-63
ACM90 5-62 10-63
ACM103 6-62
ACM103 6-62
ACM108 7-62
ACM109 7-62
ACM125 10-62
ACM145 12-62 4-63 5-64
ACM146 12-62
ACM182 6-63 4-64
ACM195 8-63
ACM209 10-63
ACM213 10-63 11-64
BIT 2 3 P192
BIT 1 1 P64
BIT 1 4 P290
MCA X1 AP203
NUM 1 1 P60
NUM 2 2 P130
NUM 4-1 P10
ACM233 6-64
NUM MATH V6(15)
ACM244 11-64

INTERPOLATION BY CONTINUED FRACTIONS
INTERPOLATION BY AITKEN ITERATION
INTERPOLATN-DIFFERENTIATN AND INTEGRATION
DIFFERENCE EXPRESSION COEFFICIENTS
CONFLUENT DIVIDED DIFFERENCES
BACKWARD DIVIDED DIFFS
FORWARD DIVIDED DIFFS
DIFFERENCES AND DERIVATIVES
LAGRANGIAN INTERPOLATION
HERMITE INTERPOLATION
NEVILLE INTERPOLATION
INTERPOLATION WITH RATIONAL FUNCTIONS

MT
ACM18 9-60 3-62
ACM70 11-61 7-62
ACM77 2-62 6-62 8-63 11-63
ACM79 2-62 3-63
ACM167 4-63 9-63
ACM168 4-63 9-63
ACM169 4-63 9-63
ACM187 7-63
ACM210 10-63 10-63 P619
ACM211 10-63
AIS17
NUM 3-3 P302

ORTHONORMALISING PROCEDURE
INOUT-ALGEBRA OF SETS-SPECIAL SUM
CALENDAR DATE TO AND FROM JULIAN DAY
SHANKS
ORTHONORMALISE COLUMNS
CALCULATION OF EASTER
SUM
INNERPRODUCT
FACTOR
REMAINDER
SOLVE LICHTENSTEIN-GERSHGORIN EQUATION
CONFORMAL MAPPING OF NEARBY CIRCULAR REG
CONVERGING FACTOR FOR CONTINUED FRACTION
DISPLAY EPSILON ALGORITHM
CALCULATE EPSILON ALGORITHM ARRAY
DISPLAY SINGULAR EPSILON ALGORITHM
CALC SINGULAR EPSILON ALGORITHM ARRAY

MZ
ACM127 10-62
ACM156 3-63 8-63
ACM199 8-63
ACM215 11-63
ACM 2-62 P91
ACM 4-62 P209 11-62 P556
MCA X1 AP119
MCA X1 AP120
MCA X1 AP125
MCA X1 AP126
B0001 P242
B0001 P256
NUM 5-4 P341
BIT 3-3 P183
BIT 3-3 P184
BIT 3-3 P191
BIT 3-3 P192

BINOMIAL COEFFICIENTS BY RECURSION
FACTORIAL
PERMUTATIONS OF INTEGERS 0 TO N
COMPOSITION GENERATOR
PERMUTE FIRST N COMPONENTS OF VECTOR X
PERMUTATION GENERATOR-LEXICOGRAPHIC ORDER
COMBINATION GENERATOR
PERMUTATION IN LEXICOGRAPHICAL ORDER
PERMUTE FIRST N COMPONENTS OF VECTOR X
PERMUTATION GENERATOR
LOCATE VECTOR IN LEXICOOG ORDERED LIST
NEXCOM-COMBS OF ELEMNTS IN COLUMN VECTOR
COMBINATION IN LEXICOGRAPHICAL ORDER
COMBINATION IN ANY ORDER
CALC COMBS OF M THINGS N AT A TIME
COMBS FROM 1 TO N AT A TIME OF M THINGS
PERMUTATIONS IN LEXICOGRAPHICAL ORDER
RANDOM PERMUTATION
PERMS WITH REPETITION

SIEVE-PRIME NUMBERS UP TO N
PRIME TWINS
GREATEST COMMON DIVISOR

NC
ACM19 10-60 6-62 8-62
ACM33 2-61
ACM71 11-61 4-62 8-62
ACM72 11-61 8-62
ACM86 4-62 8-62
ACM87 4-62 8-62 10-62
ACM94 6-62 11-62 12-62
ACM102 6-62
ACM115 8-62 10-62 12-62
ACM130 11-62
ACM151 2-63
ACM152 2-63 7-63
ACM154 3-63 8-63
ACM155 3-63 8-63
ACM160 4-63 8-63 10-63
ACM161 4-63 8-63 10-63
ACM202 9-63
ACM235 7-64
ACM242 10-64
NP
ACM35 3-61 4-62 8-62
ACM223 4-64
ACM237 3-64 12-64

GENERATE PARTITIONS IN PART COUNT FORM
GENERATE PARTITIONS WITH CONSTRAINTS

NR
ACM95 6-62
ACM114 8-62

EUCLIDEAN ALGORITHM	NZ
AUGMENTATION OF X BY Y	ACM7 4-60
CHAIN TRACING	ACM68 8-61 11-61
JACOBI SYMBOL BY QUADRATIC RECIPROCITY	ACM69 9-61
MAGIC SQUARE-EVEN ORDER	ACM99 6-62 11-62
MAGIC SQUARE-ODD ORDER	ACM117 8-62 1-63 3-63
DIOPHANTINE EQUATION	ACM118 8-62 12-62 1-63
TERM OF MAGIC SQUARE	ACM139 11-62
CALENDAR DATE TO AND FROM JULIAN DAY	ACM148 12-62 4-63
CALCULATION OF EASTER	ACM199 8-63
	ACM 4-62 P209 11-62 P556
ASSIGNMENT	OL
ECONOMISING A SEQUENCE 1	ACM27 11-60 10-63 12-63
ECONOMISING A SWQUENCE 2	ACM81 3-62
OPTIMAL CLASSIFICATION OF OBJECTS	ACM 82 3-62
SHORTEST PATH	ACM83 3-62
INTEGER SOLNS OF LINEAR PROG PROBLEMS	ACM97 6-62
DISCRETE CONVOLUTION	ACM153 2-63 8-63
	ACM208 10-63
CRITICAL PATH SCHEDULING	OP
EVALUATION OF A PERT NETWORK	ACM40 3-61 9-61 10-62 6-64
MINIMUM EXCESS COST CURVE	ACM119 8-62
TOPOLOGICAL ORDERING FOR PERT NETWORKS	ACM217 12-63
A PERT PROGRAM IN ALGOL 60	ACM219 21-63
	MCA MR56
CARBON DIOXIDE AND AIR FUNCTIONS	PH
THERMODYNAMIC PROPS OF STEAM AND WATER	AIS12
	AIS19
MULTIREGION TWO GROUP PROGRAMS	PN
	AIS14
QUANTUM MECH INTS-SLATER TYPE ORBITALS	PQ
MOLECULAR ORBITAL CALCULATION	ACM110 7-62
QUANTUM MECH INTS-SLATER TYPE INTEGRALS	ACM111 7-62
	ACM132 11-62
CHI SQUARED	SA
STATISTICAL ANALYSIS-CARD INPUT	AIS8
	AIS22
SMOOTHING BY GRAMS FORMULAS-1	SB
SMOOTHING BY GRAMS FORMULAS-2	ACM188 7-63
SMOOTH-FOURTH ORDER SMOOTHING	ACM189 7-63
	ACM216 11-63
CORRELATION COEFS WITH MATRIX MULT	SC
	ACM39 3-61
FREQUENCY DISTRIBUTION	SM
MEAN STANDARD DEVIATION AND VARIANCE	ACM212 10-63
MEAN STANDARD DEV VAR AND DEG OF FREEDOM	AIS6
CALCN OF NORMAL DISTRIBUTION FUNCTION	AIS7
	ALY 1 1 933
TRIANGULAR REGRESSION	SR
MULTIPLE REGRESSION-MARK 1	ACM142 12-62
	AIS18

POS NORMAL DEVIATE FROM RANDOM NUMBER	SS
GENERATE RANDOM NUMBER	ACM121 9-62
ERLANG PROBABILITY FOR CURVE FITTING	ACM133 11-62 12-62 3-63
NORMAL PROBABILITY FOR CURVE FITTING	ACM184 7-63
NORMAL RANDOM	ACM185 7-63
DISCRETE CONVOLUTION	ACM200 8-63
QUASI-RANDOM POINT SEQUENCE	ACM208 10-63
	ACM247 12-64
	SV
VARIANCE HOMOGENEITY TEST-BARTLETT	AIS5
MEAN STANDARD DEVIATION AND VARIANCE	AIS6
MEAN STANDARD DEV VAR AND DEG OF FREEDOM	AIS7
TWO WAY ANALYSIS OF VARIANCE	AIS9
2 TO THE K FACTORIAL ANALYSIS-YATES	AIS10
	ZA
NESTING OF FOR STATEMENT 1	ACM137 11-62
NESTING OF FOR STATEMENT 2	ACM138 11-62
DIAGRAM	ACM 1-61 P55
FUSBUDGET	ACM 1-61 P64
COLLAPSE OWN STORAGE	ACM 1-61 P64
BOOLEAN OPERATOR NOT FLAG	ACM 1-61 P71
STORAGE ALLOCATION PROCEDURES	ACM 10-61 P444
REDUNDANCY CHECK	ACM 6-62 P340
ALGOL 60 SYNTAX CHECKER FOR THE IBM 7090ORN13399	
	ZZ
ASSIGN	ACM173 6-63 10-63
CALENDAR DATE TO AND FROM JULIAN DAY	ACM199 8-63 11-64
CALCULATION OF EASTER	ACM 4-62 P209 11-62 P556
ALGOL INFORMATION SHEETS-THEIR PURPOSES	AIS1
ALGOL ON PEGASUS	AIS2
ALGOL ON DEUCE	AIS3
FREE FIELD READ	ACM239 8-64
COORDS FO ELLIPSOID	ACM240 9-64
GRAYCODE	ACM246 12-64

LOCATE ELEMENT IN LIST

procedure search(p) lowerbound:(ℓ)upperbound:(u)relation:(b);

integer p, ℓ ,u; boolean b;

comment HUCC LIBRARY PROCEDURE DHO2:

AUTHOR: J. BOOTHROYD

A Jensen procedure which searches an ordered array a between the inclusive limits a[ℓ] and a[u] for an element of given value x, say. For arrays sorted in ascending order the relations

$b = (x \leq a[p])$ and $b = (x < a[p])$ yield exit values of ℓ and u satisfying $a[\ell] < x \leq a[u]$ and $a[\ell] \leq x < a[u]$ respectively.

To locate an element of value y in the third column of a matrix A[1:10,1:5] for example, use the call

$\ell := 1; \quad u := 10;$

search(p, ℓ ,u,y \leq A[p ,3])

and to find an element of value y in the second row of the same matrix perform

$\ell := 1; \quad u := 5;$

search(p, ℓ ,u,y \leq A[2, p])

Note from these examples that ℓ and u specify the extent of the search in the call as well as serving as result variables on exit from the procedure.

SORT REAL NUMBERS INTO ASCENDING ORDER

procedure shellsort (a, n); value n; real array a; integer n;

comment HUCC LIBRARY PROCEDURE DH01:

AUTHOR J. BOOTHROYD : :

Elements a[1] through a[n] of real array a[1:n] are rearranged in ascending order. The method is that of D.A. SHELL (A High Speed Sorting Procedure. COMM. A.C.M. 2 (1959) 30-32) with subsequences chosen such that m_1 , the first value of the partition parameter, is given by

$$m_1 = 2^k - 1 \text{ for } 2^k \leq n < 2^{k+1}$$

To implement Shell's original choice of $m_1 = [n/2]$ change the first statement of the procedure to $m_1 = n$. Note that shellsort specifies the array as type real. It may thus not be used to reorder elements of integer arrays unless the specification of a is changed to integer array a and in this case the local working variable w should preferably be declared as integer w;

SORT REAL NUMBERS IN ASCENDING ORDER

procedure exsort(a,n); value n; integer n; array a;

comment HUCC LIBRARY PROCEDURE DH03:

AUTHOR J. BOOTHROYD :

A procedure which sorts the elements a[1] through a[n] of real array a[1:n] into ascending order. The method used is that of exchanging element pairs. After the first pass, which ranges over all element positions a[1] to a[n], each pass is restricted to the range bounded by the first and last exchanges which occurred on the previous pass. This significantly improves the efficiency of the method.

The procedure is simple and may be easily modified to sort numbers into decending order by changing the Boolean expression a[i+1] < a[i] to a[i+1] \geq a[i].

A version of exsort (under the identifier jensort) suitable for sorting rows or columns of an $m \times n$ matrix is included in the library as DH04;

SORT ROWS OR COLUMNS OF MATRIX

procedure jensort(ai,aiplus1,i,n); value n; real ai,aiplus1;
integer i,n;

comment HUCC LIBRARY PROCEDURE DHO4:

AUTHOR J. BOOTHROYD :

A Jensen modification of exsort (DHO3) which permits the sorting into ascending order of rows or columns of arrays of any number of dimensions. To sort the 3rd row of an array declared as A[1:n,1:m] use the call

jensort(A[3,i],A[3,i+1],i,m). To sort each column of B[1:10,1:20] use the construction

for p:= 1 step 1 until 20 do jensort(B[i,p],B[i+1,p],i,10).

These examples, and a study of the differences between DHO3 and DHO4 should make clear the significance of the parameters ai and aiplus1;

DH01

```
procedure shellsort(a,n); comment ACM 201; value n; real array a; integer n;
  begin integer i,j,k,m; real w; switch s:=endj;
    for i:=1 step 1 until n do m:=2*i-1;
    for m:=m div 2 while m ≠ 0 do
      begin k:=n-m;
        for j:=1 step 1 until k do
          begin for i:=j step -m until 1 do
            begin if a[i+m] > a[i] then goto endj;
              w:=a[i]; a[i]:=a[i+m]; a[i+m]:=w
            end i;
          end j;
        end m;
  end shellsort;
```

DH02

```
procedure search(p) lowerbound: (1)upperbound: (u)relation:(b);
  integer p,l,u; boolean b;
  comment a Jensen procedure which searches an ordered array a between the inclusive limits a[1] and a[u] for
  an element of given value, say x. For arrays sorted into ascending order the relations
  b=(x<=a[p]) and b=(x<=a[p]) produce exit values of 1 and u satisfying a[1] < x ≤ a[u] and
  a[1] ≤ x < a[u] respectively;
  begin for p:=(u+1) div 2 while u ≥ 1 do if b then u:=p-1 else l:=p+1;
    u:=u+1; l:=l-1
  end search;
```

DH03

```
procedure exsort(a,n);  value n;  integer n;  array a;
begin integer i,k,r;  switch s:=L1,L2;  real temp;
L1: k:=2;
L2: r:=n-1;  n:=0;
  for i:=k-1 step 1 until r do
    if a[i+1]<a[i] then
      begin temp:=a[i];  a[i]:=a[i+1];  a[i+1]:=temp;
        if n=0 then k:=i;  n:=i
      end;
    if n#0 then goto if k#1 then L2 else L1
end exsort;
```

DH04

```
procedure jensort(ai,aiplus1,i,n);  value n;  real ai,aiplus1;  integer n,i;
begin integer k,r;  switch s:=L1,L2;  real temp;
L1: k:=2;
L2: r:=n-1;  n:=0;
  for i:=k-1 step 1 until r do
    if aiplus1 < ai then
      begin temp:=ai;  ai:=aiplus1;  aiplus1:=temp;
        if n=0 then k:=i;  n:=i
      end;
    if n # 0 then goto if k # 1 then L2 else L1
end jensort;
```

OUTPUT GRAPH OF VECTOR ELEMENTS

```
procedure Graphic(a,m,n);
value m,n; array a; integer m,n;
comment HUCC LIBRARY PROCEDURE DP01:
AUTHOR: W. G. WARNE:
```

Plots the values of the elements of the nth order vector a on an $(m+1) \times n$ grid. Each element is plotted on a new line as a + character. Scale is chosen so that the maximum of a_i occurs m spaces from the left, and the minimum zero spaces from the left. Three newline characters are output before the first line marked by a row of $m+1$ dots and two newlines are output after the last line marked by a row of $m+1$ dots. A column of vertical bar characters marks the baseline (left margin).

Use is made of tabs to ensure optimum print speed.
Tabs must be set at 12 spaces.

```
procedure Graphic(a,n,m);
value m,n;  array a;  integer m, n;
begin      integer j, k, h, r;  real max, min, x, y;
print #13??;
for j := 0 step 1 until m do print #.?;
k := 0;
for j := 1 step 1 until n do
  k := if a[k] > a[j] then k else j;
max := a[k];
k := 0;
for j := 1 step 1 until n do
  k := if a[k] < a[j] then k else j;
min := a[k];
x := m / (max-min);
for h := 0 step 1 until n do
begin      print #1?P;
  k := x*(a[h]-min);
  r := k div 12;
  for j := 1 step 1 until r do print #t??
  r := k - r*12;
  for j := 1 step 1 until r do print #c??
  print #+?;
end;
print #1??
for j := 0 step 1 until m do print #.?;
print #12??
end Graphic;
```

OUTPUT PLOT OF ELEMENTS OF TWO VECTORS

```
procedure Bigraphic(a,b,n,m);  
value n,m; array a,b; integer m,n;  
comment HUCC LIBRARY PROCEDURE DPO2
```

AUTHOR: W. G. WARNE:

Plots values of the elements of two nth order vectors
a, and b on an $(m+1) \times n$ grid.

Each pair of elements is plotted on a new line,
 a_j being represented by (underline character) and b_j
by a dot. Scale is chosen so that the maximum of a_j or b_j
occurs m spaces from the left and the minimum of a_j , b_j
at zero spaces, and both elements are plotted with the same
scale and origin.

Otherwise format is as for Graphic(qv).

```

procedure Bigraphic(a,b,n,m);
value n,m; array a,b; integer m,n;
begin integer j,ka,kb,h,sp,l;
  real min, max,x;
  procedure position;
    begin integer i;
    if l+sp<12 then
      begin for i:=1 step 1 until sp do print ff12?;
      l:=l+sp;
      sp:=0
    end
    else if l<10 then
      begin for i:=(l+sp)div 12 step -1 until 1 do print ff12?;
      l:=l+sp-(l+sp) div 12 * 12; sp:=0;
      for i:=1 step 1 until 1 do print ff12??
    end else
      begin for i:=12-1 step -1 until 1 do print ff12??
      sp := sp-12+1; l:=0;
      position
    end
  end position;
  print ff12??
  for j:=0 step 1 until m do print f.?;
  max:= if a[0]>b[0] then a[0] else b[0];
  for j:=1 step 1 until n do
    max:= if a[j]>max or b[j]>max then (if a[j]>b[j] then
    a[j] else b[j]) else max;
  min:= if a[0]<b[0] then a[0] else b[0];
  for j:=1 step 1 until n do
    min:= if a[j]<min or b[j]<min then (if a[j]<b[j] then
    a[j] else b[j]) else min;
  x:=m/(max-min);
  for h:=0 step 1 until n do
begin print ff12??
  ka:=x*(a[h]-min);
  kb:=x*(b[h]-min);
  sp:=0;
  l:=0;

```

```
for j:=0 step 1 until m do
  begin if ka=j then
    begin position;
      print f?
    end;
    if kb = j then
      begin position;
        l:=l+1;
        print f.?
      end
    else sp:=sp+1;
  end;
print ff1??
for j:=0 step 1 until m do print f.?;
print ff1??
end Bigraphic;
```

DZ01

CHARACTER PACK AND UNPACK

```
procedure pack(x) in position:(n)error exit:(label);  
comment global integer MODE equal to the number of bits in x  
must be declared and assigned before procedure is called;  
value x,n;  
integer x,n;  
label label;
```

HUCC LIBRARY PROCEDURE DZ01:

AUTHOR: P. W. FORD:

DZ02

```
procedure unpack(x) from position:(n);  
comment as for pack procedure;  
value n;  
integer x,n;
```

HUCC LIBRARY PROCEDURE DZ02:

AUTHOR: P. W. FORD:

This pair of procedures packs and unpacks sets of words containing MODE binary digits into the available ALGOL free store. In any one program, MODE must be constant and less than 39. The packing is dense, that is, each location in store is completely filled. If more than MODE binary digits are presented to be packed, the superfluous bits are ignored. n must not be less than 1. Storing starts at location 316 and works up the store until the program's data space is reached. If current data space is about to be overwritten, the procedure pack outputs on the typewriter "STORE FULL" followed by the machine location about to be overwritten and the number (n) of the word which would be placed in the location. The procedure then exits through the label without overwriting the store. Procedure unpack does not check the upper bound since, although rubbish may be extracted, the program itself cannot be destroyed by this procedure.

The machine is not aware of any packed information, so that subsequent data space allocation by program could overwrite some of the packed store. Thus the rule is, do not declare any variables, arrays, etc. after procedure pack is called, at least until the packed data has been processed. To pack and unpack one word takes of the order of two milliseconds.

ALGOL Tape 1 is destroyed by these procedures, but the owncode system is unaffected.

DZ01

```

procedure pack(x) in position:(n) error exit:(label);
  comment global MODE equal to the number of bits in x must be
  declared and assigned before procedure is called;
  value x,n;
  integer x,n;
  label label;
begin integer a,b,c,divcount,shift,count;
array A[1:1];
switch s:=error,end	unset;
a:=address(A)-1;
b:=316;
c:=39;
elliott(~,2,0,0,2,7,n);
elliott(3,0,MODE,0,5,2,n);
elliott(5,7,0,0,0,0,0);
elliott(1,6,count,0,5,6,c);
elliott(0,4,b,0,2,0,divcount);
elliott(0,7,a,0,4,1,error);
elliott(3,0,divcount,0,0,5,b);
elliott(5,2,c,0,5,7,0);
elliott(~,7,count,0,0,7,c);
elliott(2,0,shift,0,6,7,divcount);
elliott(3,0,0,~,6,7,shift);
elliott(5,0,0,0,1,0,x);
elliott(6,7,MODE,0,5,0,0);
elliott(3,0,x,~,6,7,shift);
elliott(5,4,0,0,6,7,divcount);
elliott(2,0,0,0,0,2,shift);
elliott(0,7,MODE,0,4,1	unset);
elliott(5,4,39,0,2,0,x);
elliott(0,2,divcount,0,0,0,0);
elliott(1,0,x,1,2,0,0);
unset: elliott(4,3,end,0,4,0,end);
error:  print punch(3),2
STORE FULL?, sameline,divcount,n;
  goto label;
end:  end pack;

```

DZ02

```

procedure unpack(x) from position:(n);
  value n;
  integer n,x;
begin integer a,b,c,divcount,shift,count,diff,temp;
switch ss:=rt,pu,collate,next;
b:=316;
c:=39;
elliott(~,2,0,0,2,1,temp);
elliott(3,0,c,0,0,5,MODE);
elliott(1,0,temp,1,5,1,0);
elliott(2,0,a,0,3,0,n);
elliott(5,2,MODE,0,5,7,0);
elliott(1,6,count,0,0,0,0);
elliott(5,6,c,0,0,4,b);
elliott(2,0,divcount,0,0,5,b);
elliott(5,2,c,0,5,7,0);
elliott(0,7,count,0,4,2,pu);
elliott(2,0,diff,0,~,7,c);
elliott(2,0,shift,0,6,7,divcount);
elliott(3,0,0,0,2,0,temp);
elliott(0,2,diff,0,0,7,MODE);
elliott(4,1,rt,0,3,0,temp);
elliott(5,0,39,0,6,7,divcount);
elliott(2,7,8191,0,6,7,diff);
elliott(5,4,0,0,4,0,collate);
rt:  elliott(3,0,temp,0,6,7,shift);
elliott(5,1,0,0,4,0,collate);
pu:  elliott(0,0,divcount,1,2,7,8191);
collate: elliott(2,3,a,0,4,3,next);
next: x:=a
end unpack;

```

Double Length Floating Point Arithmetic Package.

Author: W.G. Warne

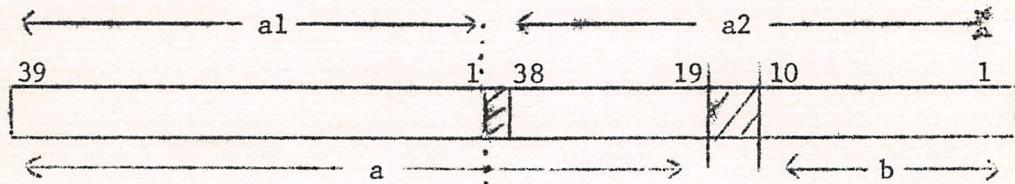
Procedures DZ03 to DZ11 provide facilities for performing real arithmetic and yield results which have twice the accuracy of the standard Algol and arithmetic operations.

Within this scheme a double length floating point number is represented by a number pair each of whose members is of type integer. For the number pair (a1, a2), for example, the floating point number $a*2^b$ is distributed within a1, a2 as follows :-

a is held as bits 39 to 1 of a1 and bits 38 to 19 of a2.

b is held as bits 10 to 1 of a2.

bits 39 and 18 to 11 of a2 are unused.



With this representation the double length floating number range is

$-2^{511} \leq x < -2^{-512}$, 0 , $2^{-512} \leq x < 2^{511}$ with a mantissa precision of 57 binary digits, roughly 17 decimal digits.

No provision is made for printing the decimal equivalent of a double length floating point number. The procedures are intended for special programming occasions when it is essential to minimise round-off errors. These procedures will be available only to those who convince either the officer-in-charge or the second officer-in-charge that their particular problem requires such precision. This restriction is imposed for the following reasons :-

- (a) TIME An average time for all procedures is 1 to 1.5 milliseconds.
- (b) STORAGE Total storage for the package is 659 locations.

ERROR INDICATIONS

Where the result of an operation exceeds the stated range, and also when the single length result of a package operation exceeds normal single length range, an error indication is given and the program is stopped. An error indication is also given if single length underflow is encountered ($0 < |x| < 2^{-256}$).

DZ03

procedure expand (s) single length real to double : (d1,d2);
value s; real s; integer d1,d2;

DZ04

procedure intexp (i) becomes : (d1,d2); value i; integer i, d1, d2;

DZ05

procedure contract (d1,d2) double to single : (s);
value d1,d2; integer d1,d2; real s;
comment error is "contract error";

DZ06

procedure mpyssd (s1) and :(s2) multiplied to give double :(d1,d2);
value s1,s2; integer d1,d2; real s1,s2;

DZ07

procedure divdss (d1,d2) divided by: (s) gives: (r);
value d1,d2,s; real s,r; integer d1,d2;
comment error is "divdss error";

DZ08

procedure negd (d1,d2) has complement :(n1,n2);
value d1,d2; integer d1,d2,n1,n2;

DZ09

procedure addddd (a1,a2) plus: (b1,b2) gives: (r1,r2);
value a1,a2,b1,b2; integer a1,a2,b1,b2,r1,r2;

DZ10

procedure mpyddd (a1,a2) multiplied by :(b1,b2) gives :(r1,r2);
value a1,a2,b1,b2; integer a1,a2,b1,b2,r1,r2;
comment error is "dfptoobig";

DZ11

procedure divddd(a1,a2) divided by: (b1,b2) gives :(t1,t2);
value a1,a2,b1,b2; integer a1,a2,b1,b2,t1,t2;
comment in the result (t1,t2) of this procedure the mantissa accuracy is 38
binary digits, (about 12 decimal digits).

BESSEL FUNCTION FOR SET OF INTEGER ORDERS

procedure BESSEL(x,n,eps,J); value x,n,eps; real x,eps;
integer n; real array J;

comment HUCC LIBRARY PROCEDURE FB01:

SOURCE : ACM 21:

Computes the values of the Bessel functions $J_p(x)$
for real x for the set of all integer orders $0 \leq p \leq n$.
These values are stored as $J[p]$ in array $J[0:n]$. See
also Comm ACM 3 (1960) p 600 and 8 (1964) p 219;

```

procedure BESSEL(x,n,eps,J);  value x,n,eps;  real x,eps;  integer n;  real array J;
comment Based on Algorithm 21, Communications of the A.C.M., 3,(1960),page 600, as
corrected and certified Vol.8(1964),page 219;
begin real dist,rec0,rec1,rec2,sum,max,MAX,abx,err,z;
integer m,k,p,q;  boolean s;  real array Jbar[0:n];
switch SS:=exit,rec,norm;
if x=0.0 then begin J[0]:=1.0;
for p:=1 step 1 until n do J[p]:=0.0;  goto exit
end;
abx:=abs(x);
dist:= if abx > 8.0 then 5.0 *abx↑(1.0/3.0)
else 10.0;
k:=entier(( if abx > n then abx else n)+dist)+1;
s:= false;  MAX:=0.2510+74;
rec0:=sum:=0.0;  rec1:=1.0;
z:=MAX*abs(x/k);
for p:=k step -1 until 1 do
begin if p>n+1 then q:=n else q:=p-1;
J[q]:=rec2:=2.0*p/x*rec1-rec0;
if abs(rec2)>z then
begin rec1:=rec1/z;  rec2:=rec2/z;  sum:=sum/z;
for m:=n step -1 until p-1 do J[m]:=J[m]/z
end;
if p=1 then sum:=sum+rec2
else if (p div 2*2)≠p then sum:=sum+2.0*rec2;
rec0:=rec1;  rec1:=rec2;
end recursion;
for p:=0 step 1 until n do J[p]:=J[p]/sum;
if s then begin max:=0;
for p:=0 step 1 until n do
begin err:=abs(J[p]-Jbar[p]);
if err>max then max:=err
end maximum error;
if max < eps then goto exit;
end
else s:= true;
for p:=0 step 1 until n do Jbar[p]:=J[p];
k:=entier(k+dist);  goto rec;
exit:
end BESSEL;

```

P. WYNN'S ARSENAL OF COMPLEX ARITHMETIC PROCEDURES

Procedures FC01 to FC26 are taken from BIT Vol 2, No.4, p.237.

Within these procedures any complex number written z in ordinary algebraic notation is represented by an ALGOL array of two elements, its real and imaginary parts.

It should be declared thus array z[0:1];

The elements $z[0]$ and $z[1]$ are respectively the real and imaginary parts. Outside the Wynn procedures the complex number is referred to as $z[i]$. The subscript i must be declared as an integer and the declaration must be valid in all blocks in which statements of procedures FC01 to FC26 occur. Inside the procedures i is assigned the value 0 before operations on real parts, 1 before operations on imaginary parts. Outside the procedures i should not be used.

Vectors, matrices and arrays with more dimensions having complex elements may be accommodated within the general scheme. For example an array with typical element $A_{p,q,r}$ might be declared as

array A[$\ell_1:u_1, \ell_2:u_2, \ell_3:u_3, 0:1$] and
referred to by A[p,q,r,i]

or alternatively array A[0:1, l1:u1, l2:u2, l3:u3]
appropriately referenced by A[i,p,q,r]

Call by name is used extensively in these procedures permitting the use of linear expressions of complex variables as actual parameters, e.g. $a_1z_1 + a_2z_2 + a_3z_3 + \dots$ where a_1, a_2, a_3 are real numbers and z_1, z_2, z_3, \dots complex scalars or complex subscripted variables.

FC01

COMPLEX ASSIGNMENT

procedure eq(one,other); real one, other;

Examples: `eq(z2[i],z1[i])`

effect is " $z_2 := z_1$ "

eq(z3[i],z1[i]+z2[i])

" $z_3 := z_1 + z_2$ "

eq(z3[i],z1[i]-z2[i])

" $z_3 := z_1 - z_2$ "

MULTIPLE COMPLEX ASSIGNMENT

procedure seqeq(third,second,first); real third,second,first;
Example: seqeq(z4[i],z3[i],z2[i]+z1[i]) "z₄ := z₃ := z₂+z₁"

FC03

COMPLEX MULTIPLICATION

procedure cm(res,one,other); real res,one,other;
Example: cm(z3[i],z2[i],z1[i]) "z₃ := z₂*z₁"

FC04

COMPLEX DIVISION

procedure cd(res,one,other); real res,one,other;
Example: cd(z3[i],z2[i],z1[i]) "z₃ := z₂/z₁"

FC05

ASSIMILATE REAL IN COMPLEX OPERATION

real procedure real(variable); real variable;
Example: eq(z2[i],z1[i]+real(a)) "z₂ := z₁+a"

FC06

ASSIMILATE IMAGINARY IN COMPLEX OPERATION

real procedure imaginary(variable); real variable;
Example: eq(z2[i],z1[i]+imaginary(a)) "z₂ := z₁+j*a"
(where, here, j² = -1)

COMPLEX CONJUGATE

real procedure cxconj(it); real it;

Example: eq(z2[i],cxconj(z1[i]))

if $z_1 = a + j * b$ then $z_2 = a - j * b$ $(j^2 = -1)$

MODULUS OF COMPLEX NUMBER

real procedure mod(it); real it;

Example: $r := \text{mod}(z1[i])$

evaluates $r = |z|$ where $|z| = (a^2+b^2)^{\frac{1}{2}}$ for $z = a+j*b$

ARGUMENT OF COMPLEX NUMBER

real procedure arg(it); real it;

Example: $\theta := \text{arg}(z1[i])$ assigns to theta the argument of z

for $z = r e^{j\theta}$ the argument lies in the range $-\pi \leq \theta \leq \pi$

POLAR FORM

procedure polar form(res,r,theta);

Example: $\text{polar form}(z1[i],R,\theta)$ assigns to the elements of $z1$ the values $r \cos \theta, r \sin \theta$.

MULTIPLICATION BY IMAGINARY OPERATOR

procedure imult(res,it); real res,it;

Example: $\text{imult}(z1[i],z2[i])$

" $z_1 := j * z_2$ "

Procedures FC12 to FC23 (except FC17) are all of the form

procedure function(res,it) and effect the operation
res:= function(it).

FC12

COMPLEX SQUARE

procedure compsq(res,it); real res,it;

FC13

COMPLEX RECIPROCAL

procedure comprecip(res,it); real res,it;

FC14

COMPLEX ROOT

procedure cxsqrt(res,it); real res,it;

N.B. USES FC08, FC09 and FC14.

FC15

COMPLEX LOGARITHM

procedure compln(res,it); real res,it;

FC16

COMPLEX EXPONENTIAL

procedure compexp(res,it); real res,it;

FC17

HYPERBOLIC FUNCTIONS

procedure hyp(sinh,cosh,y); value y; real sinh,cosh,y;

evaluates coshy=(e^y+e^{-y})/2 and sinh=(e^y-e^{-y})/2

Special care with precision is taken in evaluating sinh for small values of y.

COMPLEX SINE

procedure compsin(res,it); real res,it;

N.B. USES FC17.

COMPLEX COSINE

procedure compcos(res,it); real res,it;

N.B. USES FC17.

COMPLEX TANGENT

procedure comptan(res,it); real res,it;

N.B. USES FC04 and FC17.

COMPLEX INVERSE SINE

procedure cxarcsin(res,it); real res,it;

COMPLEX INVERSE COSINE

procedure cxarccos(res,it); real res,it;

N.B. USES FC01, FC05, FC21.

COMPLEX INVERSE TANGENT

procedure cxarctan(res,it); real res,it;

N.B. FC01, FC04, FC05, FC12 and FC14.

COMPLEX POWER OF COMPLEX VARIABLE

procedure onechother(res,one,other); real res,one,other;

Example: onechother(z3[i],z2[i],z1[i]) has the effect of

$z_3 := z_2 \dagger z_1$

TEST FOR EVEN INTEGER

boolean procedure even(integer); integer integer;

This procedure takes the value true if 2 is a factor of integer.

LOGARITHM OF COMPLEX NUMBER

procedure LOGC(a,b,c,d,FAIL); value a,b; real a,b,c,d; label FAIL;

comment HUCC LIBRARY PROCEDURE FC26:

AUTHOR ACM 243 : :

The procedure computes the number $c + di$ which is equal to the principal value of the natural logarithm of $a + bi$, i.e. such that $-\pi < d < +\pi$. A label parameter FAIL permits exit from the procedure in the event that the real part of the result approaches minus infinity. Where required in the procedure the numerical values of π , $\pi/2$ and $\ln(\sqrt{3})$ are provided;

```
procedure eq(one,other);  real one,other;  
for i:=0,1 do one:=other;
```

FC01

```
procedure seqeq(third,second,first);  real third,second,first;  
for i:=0,1 do third:=second:=first;
```

FC02

```
procedure cm(res,one,other);  real res,one,other;  
begin real Reone,Imone,Reother,Imother;  
  i:=0; Reone:=one; Reother:=other;  
  i:=1; Imone:=one; Imother:=other;  
  res:=Reone*Imother+Imone*Reother; i:=0;  
  res:=Reone*Reother-Imone*Imother  
end;
```

FC03

```
procedure cd(res,one,other);  real res,one,other;  
begin real Reone,Imone,Reother,Imother,denom;  
  i:=0; Reone:=one; Reother:=other;  
  i:=1; Imone:=one; Imother:=other;  
  denom:=Reother*Reother+Imother*Imother;  
  res:=(Imone*Reother-Reone*Imother)/denom;  
  i:=0; res:=(Reone*Reother+Imone*Imother)/denom  
end;
```

FC04

```
real procedure real(variable);  real variable;  
real :=(if i=0 then variable else 0.0);
```

FC05

should
not be
underlined.

FC 06

```
real procedure imaginary(variable); real variable;
  imaginary:=(if i=0 then 0.0 else variable);
```

FC 07

```
real procedure cxconj(it); real it;
  cxconj:=(if i=0 then it else -it);
```

FC 08

```
real procedure mod(it); real it;
  begin real Reit,Imit;
    i:=0; Reit:=it; i:=1; Imit:=it;
    mod:=sqrt(Reit*Reit+Imit*Imit)
  end;
```

FC 09

```
real procedure arg(it); real it;
  begin real Reit,Imit;
    i:=0; Reit:=it; i:=1; Imit:=it;
    arg:=(if Reit>0.0 then arctan (Imit/Reit) else
      if abs(Imit)>0.0 then
        sign(Imit)*1.57079633 -
        arctan(Reit/Imit))
      else 3.14159265)
```

*N.B. error.
bracket is wrong place*

ISSUE TWO

```
procedure polar form(res,r,theta);  real res,r,theta;
begin real r1,theta1;
  r1:=r;  theta1:=theta;  i:=0;  res:=r1*cos(theta1);
  i:=1;  res:=r1*sin(theta1)
end;
```

FC10

```
procedure imult(res,it);  real res,it;
begin real aux;
  i:=0;  aux:=it;  i:=1;  res:=aux;  aux:=it;
  i:=0;  res:=-aux
end;
```

FC11

```
procedure compsq(res,it);  real res,it;
begin real Reit,Imit;
  i:=0;  Reit:=it;  i:=1;  Imit:=it;
  res:=2.0*Reit*Imit;
  i:=0;  res:=Reit*Reit-Imit*Imit
end;
```

FC12

```
procedure comprecip(res,it);  real res,it;
begin real Reit,Imit,denom;
  i:=0;  Reit:=it;  i:=1;  Imit:=it;
  denom:=Reit*Reit+Imit*Imit;
  res:=-Imit/denom;  i:=0;  res:=Reit/denom
end;
```

FC13

FC 14

```
procedure exsqrt(res,it); real res,it;
  polar form(res,sqrt(mod(it)),0.5*arg(it));
```

FC 15

```
procedure compln(res,it); real res,it;
begin real aux;
  aux:=ln(mod(it)); i:=0; res:=aux;
  aux:=arg(it); i:=1; res:=aux
end;
```

FC 16

```
procedure compexp(res,it); real res,it;
begin real aux1,aux2;
  i:=0; aux1:=exp(it); i:=1; aux2:=it;
  res:=aux1*sin(aux2);
  i:=0; res:=aux1*cos(aux2)
end;
```

ISSUE TWO

```
procedure hyp(sinh,cosh,y);  value y;  real sinh,cosh,y;
begin real y1;
  y1:=exp(y);  cosh:=0.5*(y1+1.0/y1);
  if abs(y)>1.0 then sinh:=0.5*(y1-1.0/y1)
  else begin integer r;  real br,brplus1,brplus2;
    array CWC[0:5];
    CWC[0]:=1.13031 821;
    CWC[1]:=4.43368 498;
    CWC[2]:=5.42926 312;
    CWC[3]:=3.19843 64610-6;
    CWC[4]:=1.10367 710-8;
    CWC[5]:=2.49810-11;
    brplus1:=brplus2:=0.0;
    y1:=2.0*(2.0*y*y-1.0);
    for r:=5 step -1 until 0 do
      begin br:=y1*brplus1-brplus2+CWC[r];
      if r#0 then begin brplus2:=brplus1;
        brplus1:=br
        end
      end;
    sinh:=y*(br-brplus1)
    end
  end;
```

```
procedure compsin(res,it);  real res,it;
begin real Reit,Imit,sinhImit,coshImit;
    i:=0; Reit :=it; i:=1; Imit:=it;
    hyp(sinhImit,coshImit,Imit);
    res:=cos(Reit)*sinhImit; i:=0;
    res:=sin(Reit)*coshImit
end;
```

FC18

```
procedure compcos(res,it);  real res,it;
begin real Reit,Imit,sinhImit,coshImit;
    i:=0; Reit:=it; i:=1; Imit:=it;
    hyp(sinhImit,coshImit,Imit);
    res:=-sin(Reit)*sinhImit; i:=0;
    res:=cos(Reit)*coshImit
end;
```

FC19

```
procedure comptan(res,it);  real res,it;
begin real Reit,Imit,sinhImit,coshImit,sinReit,cosReit;
  array aux1;  aux2[0:1];
  i:=0;  Reit:=it;  i:=1;  Imit:=it;
  hyp(sinhImit,coshImit,Imit);
  sinReit:=sin(Reit);  cosReit:=cos(Reit);
  aux1[1]:=cosReit*sinhImit;
  aux2[1]:=-sinReit*sinhImit;
  aux1[0]:=sinReit*coshImit;  aux2[0]
  aux2[0]:=cosReit*coshImit;
  cd(res,aux1[i],aux2[i])
end;
```

```
procedure cxarcsin(res,it);  real res,it;
begin real x,y,d1,d2,d3,d4;
  i:=0;  x:=it;  i:=1;  y:=-it;
  d3:=x*x;  d4:=y*y;  d1:=d3+d4;  d2:=d3-d4;
  d3:=d1*d1-2.0*d2;  d4:=sqrt(d3+1.0);
  res:=sign(y)*0.5*
  ln(d1+d4+sqrt(d3+d1*(d1+2.0*d4)));
  i:=0;  res:=arctan(x*sqrt(2.0/(1.0-d2+d4)))
end;
```

FC21

```
procedure cxarccos(res,it);  real res,it;
begin array aux[0:1];
  cxarcsin(aux[i],it);
  eq(res,real(1.57079 633)-aux[i])
end;
```

FC22

```
procedure cxarctan(res,it);  real res,it;
begin array aux1,aux2[0:1];
  eq(aux1[i],it);  compsq(aux2[i],aux1[i]);
  cxsqrt(aux2[i],real(1.0)+aux2[i]);
  cd(aux2[i],aux1[i],aux2[i]);
  cxarcsin(res,aux2[i])
end;
```

FC23

```
procedure onehochother(res,one,other);  real res,one,other;
begin array aux1[0:1];
  compln(aux1[i],one);  cm(aux1[i],other,aux1[i]);
  compexp(res,aux1[i])
end;
```

FC24

```
boolean procedure even(integer);  integer integer;
even:=(integer=(integer div 2)*2);
```

FC25

```
procedure LOGC(a,b,c,d,FAIL);  value a,b;  real a,b,c,d;  label FAIL;
  if a=0 and b=0 then goto FAIL else
    begin real e,f;
      e:=0.5*a;  f:=0.5*b;
      if abs(e) < 0.5 and abs(f) < 0.5 then
        begin c:=abs(2*a)+abs(2*b);
          d:=8*a/c*a+8*b/c*b;
          c:=0.5*(ln(c)+ln(d))-1.03972077
        end
      else begin c:=abs(0.5*e)+abs(0.5*f);
        d:=0.5*e/c*e+0.5*f/c*f;
        c:=0.5*(ln(c)+ln(d))+1.03972077
      end;
      d:=if a ≠ 0 and abs(e) > abs(f) then arctan(b/a)+(if sign(a) ≠ -1 then 0 else
        if sign(b) ≠ -1 then 3.14159265 else -3.14159265) else -arctan(a/b)+1.57079633 * sign(b)
    end LOGC;
```

SUM SERIES OF CHEBYSHEV POLYNOMIALS

real procedure Chebsum(x,P,m,s);
value x,m,s; real x; integer m,s; array P;
comment HUCC LIBRARY PROCEDURE FP01
 AUTHOR: W. G. WARNE:

Evaluates the sums:-

$$\sum_{i=0}^m P[i]T_i(x) \quad \text{if } s = 1$$

$$\sum_{i=0}^m P[i]T_{2i}(x) \quad \text{if } s = 2$$

$$\sum_{i=0}^m P[i]T_{2i+1}(x) \quad \text{if } s = 3$$

where $T_i(x)$ is the i th Chebyshev polynomial defined by
 $T_i(x) = \cos(i \arccos(x))$.

The evaluation uses Chenshaw's method involving the recurrence relation for Chebyshev polynomials.

see Chenshaw, C.W.: A note on the summation of
 Chebyshev Series M.T.A.C. 9 188-120 (1955).

```

real procedure Chebsum (x,P,m,s);
value x,m,s; real x;integer m,s; array P;
comment This procedure evaluates the sum
 $\sum_{i=0}^m P_i T_i(x)$  where  $T_i(x) = \cos(i \arccos(x))$  and the prime denotes
halving of the first term, when  $s=1$ .

```

For $s=2$ and $s=3$, the sums $\sum_{i=0}^m P_{2i} T_{2i}(x)$ and $\sum_{i=0}^m P_{2i+1} T_{2i+1}(x)$ are evaluated

respectively.;

```

begin      real C0,C1,C2,y,mult;
           integer i;
           switch S1:= L1,L2,L6;
           switch S2:= L3,L7,L4,L5;
           y:=x;
           goto S1[s];
L2:L6:    y:= 2*x*x - 1.0;
L1:        C1:=C0:=C2:=0.0;
           mult := 2.0*y;
           for i:=m step -1 until 0 do
           begin      C2:= C1; C1:= C0;
           C0:= mult*C1 -C2 + P[i]
           end;
           goto S2[s];
L3:L7:    Chebsum := 0.5 * (C0 - C2); goto L5;
L4:        Chebsum := x* (C0 - C1);
L5: end of Chebsum;

```

LOGARITHM OF FACTORIAL(n)

real procedure logfac(a); integer a;
comment HUCC LIBRARY PROCEDURE FS01:
AUTHOR COMPUTER BULLETIN :

Evaluates $\ln(\text{factorial}(a))$. For values of $a \leq 7$ the procedure computes $\text{factorial}(a)$ and uses the standard function \ln to evaluate $\ln(\text{factorial}(a))$. For values of $a > 7$ the procedure uses the approximation

$$\begin{aligned}\ln(\text{factorial}(a)) &= \ln(\text{gamma}(a+1)) \\ &\approx \ln(\sqrt{2\pi}) + (a+0.5)\ln(a) - a + B1/(2a) - B2/(12a)\end{aligned}$$

where $B1(=1/6)$ and $B2(=1/30)$ are the first two Bernoulli numbers.

```
real procedure logfac(a); integer a;
comment evaluates ln(factorial(a));
begin real aa,k,b;
aa:=a;
if aa>7.5 then begin k:=1.0/aa;
logfac:=(aa+0.5)*ln(aa)+k*(0.0833333333-k*k/360.0)-aa+0.9189385333
end
else begin aa:=aa+0.5; b:=1.0;
for k:=2.0 step 1.0 until aa do b:=b*k;
logfac:=ln(b)
end
end logfac;
```

FOURTH ORDER RUNGE-KUTTA

```
procedure RKFOUR(x,y,f,h,n); value h,n; integer n;
real x,h; array y,f;
comment HUCC LIBRARY PROCEDURE GLO1:
```

AUTHOR: J. BOOTHROYD :

A procedure for solving n first order linear differential equations of the form

$$\frac{dy_1}{dx} = f_1(x, y_1, y_2, \dots, y_n)$$

$$\frac{dy_2}{dx} = f_2(x, y_1, y_2, \dots, y_n)$$

$$\frac{dy_n}{dx} = f_n(x, y_1, y_2, \dots, y_n)$$

for known starting values $(y_1)_0, (y_2)_0, \dots, (y_n)_0$.

The correct use of RKFOUR requires the following preliminary operations.

1. Array y must be initialised so that, for $y[1:n]$,

$$y[i] = (y_i)_0 \quad i = 1, 2, \dots, n$$

2. The right-hand sides of the equations must be specified by the declaration of a procedure function such that the operation of function places in array $f[1:n]$ the appropriate $f[i]$, $i = 1, 2, \dots, n$.

Examples: A $\frac{dy_1}{dx} = y_2$

$$\frac{dy_2}{dx} = -k^2 y_1$$

```
procedure function(x,y,f); real x; array y,f;
```

```
begin x:=x; f[1]:=y[2]; f[2]:=-k*k*y[1] end;
```

B $\frac{dy_1}{dx} = y_2 + x$

$$\frac{dy_2}{dx} = y_3$$

$$\frac{dy_3}{dx} = -y_2 y_3 - y_1 x^2$$

```
procedure function(x,y,f);  real x;  array y,f;  
begin f[1]:= y[2]+x;  f[2]:= y[3];  
f[3]:= -y[2]*y[3]-y[1]*x*x end;
```

A single call of RKFOUR replaces the values $y[i]$, $i=1,2,\dots,n$ appropriate to some x by the corresponding values at $x+h$. The user must, therefore, arrange to call RKFOUR successively the appropriate number of times consistent with h , the interval of interest and the values of the argument at which values of the solution are required.

Suppose values of $y[1]$ are required at intervals of 0.2 from $x=0$ to $x=5$ inclusive and the computation step interval h is 0.05.

```
x:= 0;  
for i:= 1 step 1 until 26 do  
begin print x,y[1];  
for j:= 1 step 1 until 4 do RKFOUR(      )  
end;
```

```
procedure RKFOUR(x,y,f,h,n);  value h,n;  integer n;  real x,h;  array y,f;  
begin real hby2,hby6;  integer i;  array ybar,p[1:n];  
procedure step(a,b,k1,k2);  value k1,k2;  real k1,k2;  array a,b;  
for i:=1 step 1 until n do  
begin p[i]:=p[i]+k1*f[i];  a[i]:=y[i]+k2*b[i]  end step;  
hby2:=h/2.0;  hby6:=h/6.0;  
for i:=1 step 1 until n do p[i]:=0.0;  
function(x,y,f);  step(ybar,f,1.0,hby2);  x:=x+hby2;  
function(x,ybar,f);  step(ybar,f,2.0,hby2);  
function(x,ybar,f);  step(ybar,f,2.0,h);  x:=x+hby2;  
function(x,ybar,f);  step(y,p,1.0,hby6)  
end RKFOUR;
```

MATRIX INVERSION

```
procedure mxinvert(a,n,eps,singular);  value n,eps;  
    array a;  integer n;  real eps;  label singular;  
comment  HUCC LIBRARY PROCEDURE LEO1:  
AUTHOR:  J. BOOTHROYD
```

Inverts a matrix in its own space using the Gauss-Jordan method with complete matrix pivoting. I.e., at each stage the pivot has the largest absolute value of any element in the remaining matrix. The coordinates of the successive matrix pivots used at each stage of the reduction are recorded in the successive element positions of the row and column index vectors *r* and *c*. These are later called upon by the procedure mxperm which rearranges the rows and columns of the matrix. If the matrix is singular the procedure exits to an appropriate label in the main program.

This procedure uses procedure mxperm, and for correct operation either NC02 or NC04 must previously have been declared.

```

procedure mxinvert (a,n,eps,singular); value n,eps; array a; integer n; real eps; label singular;
begin integer i,j,k,pivi,pivj,p,ri,ci,rk,cj,iless1; real pivot; integer array r,c[1:n];
  comment set row and column index vectors;
  for i:=1 step 1 until n do r[i]:=c[i]:=i;
  comment find initial pivot; pivi:=pivj:=1; pivot:=a[1,1];
  for i:=1 step 1 until n do for j:=1 step 1 until n do
    if abs(a[i,j])>abs(pivot) then begin pivi:=i; pivj:=j; pivot:=a[i,j] end;
  comment start reduction;
  for i:=1 step 1 until n do
    begin ri:=r[pivi]; r[pivi]:=r[i]; r[i]:=ri; ci:=c[pivj]; c[pivj]:=c[i]; c[i]:=ci; iless1:=i-1;
      if eps > abs(a[ri,ci]) then
        begin print punch(3), sameline, digits(3), $21?MATRIX SINGULAR?, $21?i=? ,i,$21?PIVOTS FOLLOW?;
          for i:=1 step 1 until n do
            print punch(3), sameline, digits(3), $21??,r[i], $2s4??,c[i];
          goto singular
        end;
      for j:=1 step 1 until iless1,i+1 step 1 until n do
        begin cj:=c[j]; a[ri,cj]:=a[ri,cj]/pivot end;
        a[ri,ci]:=1.0/pivot; pivot :=0;
      for k:=1 step 1 until iless1,i+1 step 1 until n do
        begin rk:=r[k];
          for j:=1 step 1 until iless1,i+1 step 1 until n do
            begin cj:=c[j]; a[rk,cj]:=a[rk,cj]-a[ri,cj]*a[rk,ci];
              if k>i and j>i and abs(a[rk,cj]) >abs(pivot) then
                begin pivi:=k; pivj:=j; pivot:=a[rk,cj] end conditional
            end jloop;
            a[rk,ci]:=-a[ri,ci]*a[rk,ci]
          end kloop
        end iloop and reduction;
  comment rearrange rows; mxperm(a[j,p],a[k,p],j,k,r,c,n,p);
  comment rearrange columns; mxperm(a[p,j],a[p,k],j,k,c,r,n,p)
end mxinvert;

```

SOLVE LINEAR EQUATIONS - ONE R.H. SIDE

procedure SOLVEQ(a,n); value n; integer n; real array a;
comment HUCC LIBRARY PROCEDURE LEO2;
AUTHOR: J. BOOTHROYD

A procedure which solves $Ax = b$ where A is $[1:n, 1:n]$

and b is $[1:n]$. The elements of A and b must
occupy the elements of $a[1:n, 1:n+1]$ as follows:-

$a[i, j] = A[i, j]$ $i = 1, 2, \dots, n, j = 1, 2, \dots, n$.

$a[i, n+1] = b[i]$ $i = 1, 2, \dots, n$.

i.e. the R.H.S. vector b occupies column $n+1$ of the
augmented matrix a. On exit from the procedure the
solutions $x[i]$ occupy $a[i, n+1]$ $i = 1, 2, \dots, n$.

The method is Gauss-Jordan reduction to diagonal form
with partial pivoting. No row exchanges are performed
(the matrix is reduced in-situ), but a final permutation
of $a[i, n+1]$ is performed to reorder the solution vector.
The procedure does not include a test for singularity or
ill-condition, and should not be used if the equations
may be suspected of either attribute.

```

procedure SOLVEQ(a,n);  value n;  integer n;  real array a;
begin integer i,j,k,piv,ri,rk,n1;  real pivot,arki;  integer array r[1:n];  switch S:=L;
  for i:=1 step 1 until n do r[i]:=i;  n1:=n+1;
  for i:=1 step 1 until n do
    begin piv:=i;  ri:=r[i];  pivot:=a[ri,i];
      for k:=i+1 step 1 until n do
        if abs(a[r[k],i])>abs(pivot) then
          begin piv:=k;  pivot:=a[r[k],i] end;
        ri:=r[piv];  r[piv]:=r[i];  r[i]:=ri;
        for j:=i+1 step 1 until n1 do a[ri,j]:=a[ri,j]/pivot;
        for k:=1 step 1 until i-1,i+1 step 1 until n do
          begin rk:=r[k];  arki:=a[rk,i];
            for j:=i+1 step 1 until n1 do a[rk,j]:=a[rk,j]-arki*a[ri,j]
          end for k;
        end reduction;
      for i:=n step -1 until 2 do
        begin k:=r[i];
        L:  if k
```

MATRIX DECOMPOSITION Ax = b

procedure ATOLUL(a,r,n,eps,singular); value n,eps;
 integer n; real array a; integer array r; label singular;
comment HUCC LIBRARY PROCEDURE LEO3
AUTHOR: J. BOOTHROYD

Performs an in situ decomposition of a matrix A into quasi triangular matrices L, a lower triangle, and U a strictly upper triangle, by Gaussian elemination with partial pivoting. The successive pivots of the reduction constitute the quasi diagonal elements of L. These are chosen to be the elements of maximum modulus in the leading column of successive reduced matrices. The pivotal row subscripts are recorded in the non-local vector r[1:n], the ith pivotal row subscript occupying r[i]. If the matrix is singular (eps is the appropriate tolerance) the procedure records the ith stage of the reduction at which singularity is detected by changing the sign of r[i] before it exists to the main program via the label parameter singular.

```
procedure ATCLU1(a,r,n,eps,singular);  value n,eps;  integer n;  real eps;
real array a;  integer array r;  label singular;
begin integer i,j,k,piv,ri,rk;  real pivot,arki;
  for i:=1 step 1 until n do r[i]:=i;
  for i:=1 step 1 until n do
    begin piv:=i;  ri:=r[i];  pivot:=a[ri,i];
    for k:=i+1 step 1 until n do
      if abs(a[r[k],i])>abs(pivot) then begin piv:=k; pivot:=a[r[k],i] end;
      if abs(pivot)<eps then
        begin print punch(3),#1?MATRIX IS SINGULAR?; r[piv]:=-r[piv];
          goto singular
        end;
      ri:=r[piv];  r[piv]:=r[i];  r[i]:=ri;
      for j:=i+1 step 1 until n do a[ri,j]:=a[ri,j]/pivot;
      for k:=i+1 step 1 until n do
        begin rk:=r[k];  arki:=a[rk,i];
          for j:=i+1 step 1 until n do a[rk,j]:=a[rk,j]-arki*a[ri,j]
        end for k
      end for i
  end ATCLU1;
```

FORWARD AND BACK SOLUTION Ax = b

```

procedure LU1SOL(a,b,r,n); value n; real array a,b;
integer array r; integer n;
comment HUCC LIBRARY PROCEDURE LEO4
AUTHOR: J. BOOTHROYD

```

Solves the equation $Ax=b$ by the successive operations $Lf=b$ and $Ux=f$ where L and U are respectively a lower and strictly upper triangle such that $L(U+I)=A$. The elements of L and U share the array a in accordance with the pivoting strategy of procedure AT01U1.

i.e. $L[i,j]$ occupy $a[r[i],j]$ $j \leq i$ $i=1,2,\dots,n$
 and $U[i,j]$ occupy $a[r[i],j]$ $j > i$ $i=1,2,\dots,n$.

$r[1:n]$ is a non-local permutation vector in which is stored the a array row subscripts of successive rows of L . The procedure uses procedure sigma (MS01).

Notes: For a matrix $A[1:n,1:n]$, vector $b[1:n]$, the procedure

calls ATOLU1(A,r,n,eps,bad) LEO3

LU1SOL(A,b,r,n) LEO4

PERMB(b,r,n) NCO1

Solve $Ax=b$ for one right hand side.

For large sets of equations with many right hand sides use:-

ATOLU1(A,r,n,eps,bad) LEO3

→ Read b

LU1SOL(A,b,r,n) LEO4

PERMB(b,r,n) NCO1

Print b

Alternatively the output procedure ZP01 may be used in place of PERMB, Print b.

procedure LU1SOL(a,b,r,n); value n; real array a,b; integer array r; integer n;
comment solves the equation $Ax=b$ by the successive operations $Lf=b$ and $Ux=f$
where L and U are respectively a lower and strictly upper triangle such that
 $L(U+I)=A$. The elements of L and U share the array a in accordance with the
pivoting strategy of procedure ATOLU1.

i.e. $L[i,j]$ occupy $a[r[i],j]$ $j \leq i$ $i = 1, 2, \dots, n$

and $U[i,j]$ occupy $a[r[i],j]$ $j > i$ $i = 1, 2, \dots, n$.

$r[1:n]$ is a non-local permutation vector in which is stored the a array row
subscripts of successive rows of L . The procedure uses procedure sigma;

begin integer i,j,ri,iless1;
for i:=1 step 1 until n do
 begin ri:=r[i]; iless1:=i-1;
 b[ri]:=(b[ri]-sigma(a[ri,j]*b[r[j]],j,1,iless1))/a[ri,i]
 end;
 for i:=n-1 step -1 until 1 do
 begin ri:=r[i];
 b[ri]:=b[ri]-sigma(a[ri,j]*b[r[j]],j,i+1,n)
 end
end LU1SOL;

MATRIC DECOMPOSITION BY CROUT'S METHOD

```
procedure crout(a,r,n,eps,singular); value n,eps; integer n;
real eps; real array a; integer array r; label singular;
comment HUCC LIBRARY PROCEDURE LE05:
AUTHOR: J. BOOTHROYD : 
```

Performs an in situ decomposition of a matrix A into quasi triangular matrices L, a lower triangle and U, a strictly upper triangle, by Crout's method with partial pivoting. The successive pivots of the reduction constitute the quasi diagonal elements of L. These are chosen to be the elements of maximum modulus in the leading column of successive reduced matrices. The pivotal row subscripts are recorded in the non-local vector r[1:n], the ith pivotal row subscript occupying r[i]. If the matrix is singular (eps is the appropriate tolerance) the procedure records the ith stage of the reduction at which singularity is detected by changing the sign of r[i] before it exits to the main program via the label parameter singular. This procedure uses the procedure sigma.

The results of this procedure are identical with those of LE03. LE05 is not a direct replacement for LE03 since its identifier is different and LE05 uses MS01.

```
procedure crout(a,r,n,eps,singular);  value n,eps;  integer n;  real eps;
real array a;  integer array r;  label singular;
begin integer i,j,k,piv,ri,rk,iless1;  real pivot,arki;
  for i:=1 step 1 until n do r[i]:=i;
  for i:=1 step 1 until n do
    begin pivot:=0.0;  piv:=i;  iless1:= i-1;
    for k:=i step 1 until n do
      begin rk:=r[k];
        arki:=a[rk,i]:=a[rk,i]-sigma(a[rk,j]*a[r[j],i],j,1,iless1);
        if abs(arki)>abs(pivot) then begin piv:=k; pivot:=arki end;
      end;
    if abs(pivot)<eps then begin print punch(3),££1?MATRIX IS SINGULAR?;
      r[piv]:=-r[piv];  goto singular
    end;
    ri:=r[piv];  r[piv]:=r[i];  r[i]:=ri;
    for j:=i+1 step 1 until n do a[ri,j]:=(a[ri,j]-sigma(a[ri,k]*a[r[k],j],k,1,iless1))/pivot
  end for i
end crout;
```

TRIANGULAR MATRIX PRODUCT LU=A

```
procedure LULTOA(a,r,n);  value n;  integer n;  real array a;  
integer array r;  
comment  HUCC LIBRARY PROCEDURE LEO6:  
AUTHOR:  J. BOOTHROYD      :
```

Performs an in situ matrix multiplication of quasi-triangular matrices L, a lower triangle, and U, a strictly upper triangle to form a matrix A given, in effect, by $A=L(U+I)$. The elements of L and U are distributed within the array a in accordance with the pivoting strategy of procedure ATOLU1, i.e. elements $L[i,j]$ occupy positions $a[r[i],j]$ $j \leq i$, $i=1,2,\dots,n$ elements $U[i,j]$ occupy positions $a[r[i],j]$ $j > i$, $i=1,2,\dots,n$. $r[1:n]$ is a non-local permutation vector in which is recorded the a array row subscripts of successive rows of L.

Note: LEO6 is the inverse of LEO3. For a matrix A the operations

ATOLU1(A,r,n,eps,singular)

LULTOA(A,r,n)

will leave A unchanged except for corruptions to the elements of A caused by computational errors.

```
procedure LU1TOA(a,r,n);  value n;  integer n;  real array a;  integer array r;
comment performs an in situ matrix multiplication of quasi-triangular matrices L,
a lower triangle, and U, a strictly upper triangle to form a matrix A given, in
effect, by A=L(U+I).  The elements of L and U are distributed within the array
a in accordance with the pivoting strategy of procedure ATCLU1, i.e.
elements L[i,j] occupy positions a[r[i],j] j≤i, i=1,2,....., n
elements U[i,j] occupy positions a[r[i],j] j>i, i=1,2,....., n.
r[1:n] is a non-local permutation vector in which is recorded the a array row
subscripts of successive rows of L;
begin integer i,j,k,ii,ri;  real sum;
  for i:=n step -1 until 1 do
    begin ii:=i+1;  ri:=r[i];
      for j:=n step -1 until ii do
        begin sum:=a[ri,j]*a[ri,i];
          for k:= i-1 step -1 until 1 do sum:=sum+a[ri,k]*a[r[k],j];
          a[ri,j]:=sum
        end;
      for j:= i step -1 until 1 do
        begin sum:=a[ri,j];
          for k:= j-1 step -1 until 1 do sum:=sum+a[ri,k]*a[r[k],j];
          a[ri,j]:=sum
        end
    end
  end LU1TOA;
```

INVERT QUASI LOWER TRIANGULAR MATRIX

```
procedure INVL(a,r,n);  value n;  integer n;  real array a;  
integer array r;  
comment  HUCC LIBRARY PROCEDURE LE07:  
AUTHOR:  J. BOOTHROYD  :
```

Inverts a quasi triangular matrix L in its own space.

The elements of the lower triangle L are distributed within the array a in accordance with the pivoting strategy of procedure ATOLU1, i.e.

elements $L[i,j]$ occupy positions $a[r[i],j]$ $j \leq i, i=1,2,\dots,n$
 $r[1:n]$ is a non-local permutation vector in which is recorded the a array row subscripts of successive rows of L.
The procedure uses procedure sigma (MS01).

INVERT QUASI UPPER TRIANGULAR MATRIX

```
procedure INVUL(a,r,n);  value n;  integer n;  real array a;  
integer array r;  
comment  HUCC LIBRARY PROCEDURE LE08:  
AUTHOR:  J. BOOTHROYD      :
```

Inverts a quasi-triangular matrix in its own space.

The elements of U, a strictly upper triangle, are distributed within the array a in accordance with the pivoting strategy of procedure ATOLU1, i.e.

elements $U[i,j]$ occupy positions $a[r[i],j]$ $j > i$ $i=1,2,\dots,n$.
 $r[1:n]$ is a non-local permutation vector in which is recorded the a array row subscripts of successive rows of U. The procedure uses procedure sigma (MS01).

LE07

```
procedure INV(a,r,n); value n; integer n; real array a; integer array r;
begin integer i,j,k,ri,iless1; real arii;
for i:=1 step 1 until n do
begin ri:=r[i]; iless1:=i-1;
arii:=a[ri,1];
for j:=1 step 1 until iless1 do a[ri,j]:=-sigma(a[ri,k]*a[r[k],j],k,j,iless1)/arii;
a[ri,i]:=1.0/arii
end
end INV;
```

LE08

```
procedure INVU1(a,r,n); value n; integer n; real array a; integer array r;
begin integer i,j,k,iplus1,jless1,ri;
for i:=n-1 step -1 until 1 do
begin ri:=r[i]; iplus1:=i+1;
for j:=n step -1 until iplus1 do
begin jless1:=j-1;
a[ri,j]:=(a[ri,j]+sigma(a[ri,k]*a[r[k],j],k,iplus1,jless1))
end
end
end INVU1;
```

TRIANGULAR MATRIX PRODUCT UL=A

procedure U1LTOA(a,r,n); value n; integer n; integer array r;

array a;

comment HUCC LIBRARY PROCEDURE LEO9:

AUTHOR: J. BOOTHROYD : :

Performs an in-situ matrix product $(U+I)L=A$ where U and L are respectively a strictly upper triangle and a lower triangle occupying the array a in accordance with the pivoting strategy of procedure ATOLU1

i.e. $L[i,j]$ occupy $a[r[i],j]$ $j \leq i$ $i=1,2,\dots,n$

and $U[i,j]$ occupy $a[r[i],j]$ $j > i$ $i=1,2,\dots,n$.

$r[1:n]$ is a non-local permutation vector in which is recorded the a array row subscripts of successive rows of L . The procedure uses procedure sigma (MS01).

```
procedure U1LTOA(a,r,n); value n; integer n; integer array r; array a;
comment performs an in- situ matrix product (U+I)L=A where U and L are respectively
a strictly upper triangle and a lower triangle occupying the array a in accordance
with the pivoting strategy of procedure ATOLU1
    i.e. L[i,j] occupy a[r[i],j]  j<=i  i=1,2,...,n
    and   U[i,j] occupy a[r[i],j]  j>i  i=1,2,...,n.
r[1:n] is a non-local permutation vector in which is recorded the a array row
subscripts of successive rows of L. The procedure uses procedure sigma;
begin integer i,j,k,nless1,ri;
    nless1:=n-1;
    for i:= 1 step 1 until nless1 do
        begin ri:= r[i];
            for j:= 1 step 1 until i do a[ri,j]:=a[ri,j]+sigma(a[ri,k]*a[r[k],j],k,i+1,n);
            for j:=i+1 step 1 until n do a[ri,j]:= sigma(a[ri,k]*a[r[k],j],k,j,n)
        end
    end U1LTOA;
```

LIBRARY PROCEDURES LEO3 to LEO9 produce and perform operations on quasi triangular matrices L and U1 where L is a lower triangle and U a strictly upper triangle. LIBRARY PROCEDURES LE10 to LE14 correspond to some, but not all of LEO3 to LEO9 for matrices L1 and U respectively a strictly lower and an upper triangle.

LE10

procedure ATOL1U(a,r,n,eps,singular); value n,eps; integer n;
real eps; integer array r; label singular;
comment HUCC LIBRARY PROCEDURE LE10:
AUTHOR: J. BOOTHROYD :
see note above and LEO3.

LE11

procedure L1USOL(a,b,r,n); value n; real array a,b;
integer array r; integer n;
comment HUCC LIBRARY PROCEDURE LE11:
AUTHOR: J. BOOTHROYD :
see note above and LEO4.

LE12

procedure INV1L(a,r,n); value n; real array a; integer array r;
comment HUCC LIBRARY PROCEDURE LE12:
AUTHOR: J. BOOTHROYD :
see note above and LEO7.

LE10

to

LE14

cont.

procedure INVU(a,r,n); value n; integer n; real array a;
integer array r;

comment HUCC LIBRARY PROCEDURE LE13:

AUTHOR: J. BOOTHROYD :

see note overleaf and LE08.

LE14

procedure UL1TOA(a,r,n); value n; integer n; array a;

comment HUCC LIBRARY PROCEDURE LE14:

AUTHOR: J. BOOTHROYD :

see note overleaf and LE09.

```

procedure ATOL1U(a,r,n,eps,singular);  value n,eps;  integer n;  real eps;
real array a;  integer array r;  label singular;
begin
    integer i,j,k,ri,rk,piv;  real pivot,arki;
    for i:=1 step 1 until n do r[i]:=i;
    for i:=1 step 1 until n do
        begin piv:=i;  pivot:=a[r[i],i];
        for k:=i+1 step 1 until n do
            if abs (a[r[k],i])>abs(pivot)then begin piv:=k;  pivot:=a[r[k],i] end;
            if abs(pivot)<eps then
                begin print punch(3),££1?MATRIX IS SINGULAR?;  r[piv]:=-r[piv];
                goto singular
                end;
            ri:=r[piv];  r[piv]:=r[i];  r[i]:=ri;
            for k:=i+1 step 1 until n do
                begin rk:=r[k];  arki:=a[rk,i]:=a[rk,i]/pivot;
                for j:=i+1 step 1 until n do a[rk,j]:=a[rk,j]-arki*a[ri,j];
                end
            end
        end
end ATOL1U;

```

LE11

```
procedure L1USOL(a,b,r,n); value n; real array a,b; integer array r; integer n;
begin integer i,j,ri,iless1;
  for i:=2 step 1 until n do
    begin ri:=r[i]; iless1:=i-1;
      b[ri]:=b[ri]-sigma(a[ri,j]*b[r[j]],j,1,iless1);
    end;
  for i:=n step -1 until 1 do
    begin ri:=r[i];
      b[ri]:=(b[ri]-sigma(a[ri,j]*b[r[j]],j,i+1,n))/a[ri,i]
    end
end L1USOL;
```

LE12

```
procedure INVL1(a,r,n); value n; integer n; real array a; integer array r;
begin integer i,j,k,ri,iless1,jplus1;
  for i:=2 step 1 until n do
    begin ri:=r[i]; iless1:=i-1;
      for j:=1 step 1 until iless1 do
        begin jplus1:=j+1;
          a[ri,j]:=-(a[ri,j]+sigma(a[ri,k]*a[r[k],j],k,jplus1,iless1))
        end
    end
end INVL1;
```

LE13

```
procedure INVU(a,r,n);  value n;  integer n;  real array a;  integer array r;
begin      integer i,j,k,ri,iplus1;  real arri;
    for i:=n step -1 until 1 do
        begin ri:=r[i];  iplus1:=i+1;  arri:=a[ri,i];
            for j:=n step -1 until iplus1 do a[ri,j]:=-sigma(a[ri,k]*a[r[k],j],k,iplus1,j)/arri;
            a[ri,i]:=1.0/arri
        end
    endINVU;
```

LE14

```
procedure UL1TOA(a,r,n);  value n;  integer n;  integer array r;  array a;
begin integer i,j,k,nless1,iless1,ri;
    nless1:=n-1;
    for i:=1 step 1 until n do
        begin ri:=r[i];  iless1:=i-1;
            for j:=1 step 1 until iless1 do a[ri,j]:=sigma(a[ri,k]*a[r[k],j],k,i,n);
            for j:=i step 1 until nless1 do a[ri,j]:=a[ri,j]+sigma(a[ri,k]*a[r[k],j],k,j+1,n)
        end
    end UL1TOA;
```

SOLVE TRIDIAGONAL LINEAR EQUATIONS

procedure tridiag(a,b,c,d,lo,hi,eps,fail); value lo,hi,eps;
integer lo,hi; real eps; label fail; real array a,b,c,d;
comment HUCC LIBRARY PROCEDURE LE15:

AUTHOR J. BOOTHROYD :

Solves the equations

$$\begin{bmatrix} b_1 & c_1 & & & & \\ a_2 & b_2 & c_2 & & & \\ a_3 & b_3 & c_3 & & & \\ & & & & & \\ & & a_{n-1} & b_{n-1} & c_{n-1} & \\ & & a_n & b_n & & \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_{n-1} \\ x_n \end{bmatrix} = \begin{bmatrix} d_1 \\ d_2 \\ d_3 \\ \vdots \\ d_{n-1} \\ d_n \end{bmatrix}$$

by the recurrence formulae

$$x_n = d_n' \quad x_r = d_r' - c_r' x_{r+1}$$

$$d_1' = d_1/b_1 \quad c_1' = c_1/b_1$$

$$d_r' = \frac{d_r - a_r d_{r-1}'}{b_r - a_r c_{r-1}'} \quad c_r' = \frac{c_r}{b_r - a_r c_{r-1}'} \quad r=2,3,\dots,n.$$

The coefficients occupy the elements of arrays a,b,c,d as follows:

a_2 to a_n occupy $a[lo]$ to $a[hi-1]$ of $a[lo:hi-1]$

b_1 to b_n occupy $b[lo]$ to $b[hi]$ of $b[lo:hi]$

c_1 to c_{n-1} occupy $c[lo]$ to $c[hi-1]$ of $c[lo:hi-1]$

d_1 to d_n occupy $d[lo]$ to $d[hi]$ of $d[lo:hi]$

eps is a singularity tolerance. If $b_r - a_r c_{r-1}' < \text{eps}$ for any r the set is presumed singular and exit occurs through the label parameter fail.

The values of $b_r - a_r c_{r-1}'$ and r are recorded on the control typewriter.

The solution is delivered in array d and array c is overwritten in the process;

```
procedure tridiag(a,b,c,d,lo,hi,eps,fail);  value lo,hi,eps;  integer lo,hi;
real eps;  label fail;  real array a,b,c,d;
begin          integer i,iless1;  real w,ailess1;
  d[lo]:=d[lo]/b[lo];  w:=b[lo];
  for i:=lo+1 step 1 until hi do
    begin iless1:=i-1;  ailess1:=a[iless1];
    c[iless1]:=c[iless1]/w;
    w:=b[i]-ailess1*c[iless1];
    if w<eps then begin print punch(3),scaled(3),digits(3),
                      ffl1?TRIDIAG MATRIX SINGULAR w=? ,w,fft?i=? ,i;
                      goto fail
    end;
    d[i]:=(d[i]-ailess1*d[iless1])/w
  end;
  for i:=hi-1 step -1 until lo do d[i]:=d[i]-c[i]*d[i+1]
end tridiag;
```

LE 16

DECOMPOSITE SYMMETRIC POSITIVE DEFINITE MATRIX

procedure choleski (a,n,fail); value n; integer n; array a; label fail;

comment HUCC LIBRARY PROCEDURE LE 16 (revised)

Author : J. Boothroyd.

Performs the decomposition of a symmetric matrix A into lower and upper triangles L, L^T such that $LL^T=A$. The elements of the lower triangle of A are replaced by the elements of L. The strictly upper triangle of A is not used. If the decomposition fails, indicating that A is non positive definite the procedure exits through the label fail. The procedure uses procedure sigma (MS 01);

LE 17

INVERT LOWER TRIANGULAR MATRIX

procedure linv (a,n); value n; integer n; array a;

comment HUCC LIBRARY PROCEDURE LE 17 (revised)

Author : J. Boothroyd

Replaces the lower triangular portion L of $a[1:n,1:n]$ by the elements of L^{-1} . The procedure uses MS 01;

LE 18

MATRIX MULTIPLICATION

procedure mxmult (a,b,c,m,n,p); value m,n,p;

integer m,n,p; array a,b,c;

comment HUCC LIBRARY PROCEDURE LE 18

Author : J. Boothroyd.

$c[1:m,1:p] := a[1:m,1:n] * b[1:n,1:p];$

This procedure uses MS 01.

```
procedure choleski(a,n,fail); value n; integer n; array a; label fail;
begin integer i,j,k,iless1; real aii;
  for i:= 1 step 1 until n do
    begin iless1:= i-1;
      aii:=a[i,i]-sigma(a[i,k]*a[i,k],k,1,iless1);
      if aii<0 then goto fail;
      aii:=a[i,i]:=sqrt(aii);
      for j:= i+1 step 1 until n do
        a[j,i]:= (a[j,i]-sigma(a[i,k]*a[j,k],k,1,iless1))/aii
    end
endcholeski;
```

LE16

```
procedure linv(a,n); value n; integer n; array a;
comment inverts the lower triangle of a in situ;
begin integer i,j,k,iless1; real aii;
  for i:=1 step 1 until n do
    begin iless1:= i-1; aii:=a[i,i];
      for j:= 1 step 1 until iless1 do
        a[i,j]:= -sigma(a[i,k]*a[k,j],k,j,iless1)/aii;
      a[i,i]:= 1.0/aii
    end
end linv;
```

LE17

```
procedure mxmult(a,b,c,m,n,p); value m,n,p; integer m,n,p; array a,b,c;
comment c[m p]:= a[m n] b[n p];
begin integer i,j,k;
  for i:= 1 step 1 until m do
    for j:= 1 step 1 until p do c[i,j]:=sigma(a[i,k]*b[k,j],k,1,n)
endmxmult;
```

LE18

DECOMPOSE BANDMATRIX INTO LOWER AND UPPER 'TRIANGLES'

```
procedure bandlr (a,n,m,eps,escape); value n,m;  
  
integer n,m; array a; real eps; label escape;
```

comment HUCC LIBRARY PROCEDURE LE 19

Author : J. Boothroyd.

Decomposes the n^{th} order band matrix of width m stored in $a[1:n,1:m]$ into a lower 'triangle' L and a strictly upper triangle U . At entry to the procedure the elements of band matrix $A[1:n,1:n]$ must occupy the matrix $a[1:n,1:m]$ in accordance with the following rules :-

$A[i, j]$ occupies $a[i, j]$ $1 \leq i \leq (m+1) \text{ div } 2$

$A[i, j]$ occupies $a[i, j-i+(m+1)\text{div}2] \quad (m+1)\text{div}2 < i \leq n$

Diagrammatically

A11	A12	A13	0	0	0		A11	A12	A13	0	0
A21	A22	A23	A24	0	0		A21	A22	A23	A24	0
A31	A32	A33	A34	A35	0	→	A31	A32	A33	A34	A35
0	A42	A43	A44	A45	A46		A42	A43	A44	A45	A46
0	0	A53	A54	A55	A56		A53	A54	A55	A56	0
0	0	0	A64	A65	A66		A64	A65	A66	0	0

UNUSED ELEMENTS of matrix a must be set to zero.

This procedure does not incorporate a pivoting strategy.
eps is a tolerance which provides a criterion for deciding
whether any leading submatrix is singular in which event the
procedure exits to the label escape:

FORWARD SOLUTION AND BACK SUBSTITUTION FOR BAND EQUATION Ax=b

```
procedure bandsol(a,n,mb); value n,m;  
integer n,m; array a,b;
```

comment HUCC LIBRARY PROCEDURE LE 20

Author : J. Boothroyd

Solves the equation $Ax=b$ where A is a bandmatrix by the successive operations $Lf=b$ and $Ux=f$ where L and U are respectively a lower triangle and a unit upper triangle occupying the array $a[1:n;1:m]$ following the use of procedure LE 19. The solution vector replaces the vector $b[1:n]$.

Procedures LE 19 and LE 20 may together be used to solve $Ax=b$ for unlimited right-hand sides by some scheme similar to:-

Read Matrix A

bandlr (A,n,m,eps,escape) LE 19

→ Read vector b

bandsol (A,n,m,b) LE 20

Print solution vector b

```

procedure bandlr(a,n,m,eps,escape); value n,m; integer n,m; real array a; label escape; real eps;
begin
  integer i,j,k,jlim,klim,nless1; real pivot;
  jlim:=klim:=(m+1) div 2; nless1:=n-1;
  for i:=1 step 1 until nless1 do
    begin pivot:=a[i,1]; if abs(pivot)<eps then goto escape;
      for j:=2 step 1 until jlim do a[i,j]:=a[i,j]/pivot;
      for k:=i+1 step 1 until klim do
        begin pivot:=a[k,1];
          for j:=2 step 1 until jlim do a[k,j-1]:=a[k,j]-pivot*a[i,j];
          for j:=jlim +1 step 1 until m do a[k,j-1]:=a[k,j];
          a[k,m]:=pivot
        end;
      if klim<=n then klim:=klim+1
    end
  end;
end;

```

```

procedure bandsol(a,n,m,b); value n,m; integer n,m; real array a,b;
begin
  integer i,j,k,lim,jlim; real sum;
  lim:=(m+1) div 2; jlim:=1;
  for i:=1 step 1 until n do
    begin sum:=0.0; k:=0;
      for j:=i-1 step -1 until jlim do
        begin sum:=sum+b[j]*a[i,m-k]; k:=k+1 end;
        b[i]:=(b[i]-sum)/a[i,1];
        if i-jlim+2>lim then jlim:=jlim+1
    end forward solution;
  jlim:=n;
  for i:=n-1 step -1 until 1 do
    begin sum:=0.0; k:=2;
      for j:=i+1 step 1 until jlim do
        begin sum:=sum+b[j]*a[i,k]; k:=k+1 end;
        b[i]:=b[i]-sum;
        if jlim+2-i>lim then jlim:=jlim-1
    end back substitution
  end bandsol;

```

SOLVE BAND EQUATIONS Ax=b

```
procedure bandmx (a,n,m,b,eps,singular); value n,m,eps;
real eps; integer m,n; label singular; array a,b;
```

comment HUCC LIBRARY PROCEDURE LE 24

Author : J. Boothroyd

Solves the n^{th} order linear equation $Ax=b$ where the elements of A form a band matrix of width m . The band elements of $A[1:n,1:n]$ must occupy the array $a[1:n,1:m]$ in accordance with the rules:-

$A[i,j]$ occupies $a[i,j]$ $1 \leq i \leq (m+1) \text{div} 2$

$A[i,j]$ occupies $a[i,j-i+(m+1) \text{div} 2]$ $(m+1) \text{div} 2 < i \leq n$

Diagrammatically

A11	A12	A13	0	0	0		A11	A12	A13	0	0
A21	A22	A23	A24	0	0		A21	A22	A23	A24	0
A31	A32	A33	A34	A35	0	→	A31	A32	A33	A34	A35
0	A42	A43	A44	A45	A46		A42	A43	A44	A45	A46
0	0	A53	A54	A55	A56		A53	A54	A55	A56	0
0	0	0	A64	A65	A66		A64	A65	A66	0	0

Array $A[1:n,1:n]$

array $a[1:n,1:m]$

UNUSED ELEMENTS of matrix a must be set to zero.

The procedure uses maximum column pivots and the solution vector replaces the elements of vector $b[1:n]$. eps is a singularity tolerance and the procedure exits to the label singular if the matrix is singular or excessively ill conditioned;

```

procedure bandmx(a,n,m,b,eps,singular); value n,m,eps; real eps; integer m,n; label singular; array a,b;
begin
  integer i,j,k,nless1,lim,piv,ri,rk,iless1; real pivot,bri,ark1,sum; integer array r[1:n];
  for i:=1 step 1 until n do r[i]:=i; nless1:=n-1; lim:=(m+1) div 2;
  for i:=1 step 1 until nless1 do
    begin piv:=i; pivot:=a[r[i],1];
    for k:=i+1 step 1 until lim do
      if abs(a[r[k],1])>abs(pivot) then begin piv:=k; pivot:=a[r[k],1] end;
      if abs(pivot)<eps then begin print punch(3), '#1?MATRIX SINGULAR?; goto singular end;
      ri:=r[piv]; r[piv]:=r[i]; r[i]:=ri;
      bri:=b[ri]:=b[ri]/pivot;
      for j:=2 step 1 until m do a[ri,j]:=a[ri,j]/pivot;
      for k:=i+1 step 1 until lim do
        begin rk:=r[k]; ark1:=a[rk,1];
        b[rk]:=b[rk]-ark1*bri;
        for j:=2 step 1 until m do a[rk,j-1]:=a[rk,j]-ark1*a[ri,j];
        a[rk,m]:=0.0
        end;
      if lim<=n then lim:=lim+1
    end;
    ri:=r[n]; b[ri]:=b[ri]/a[ri,1];
    lim:=2;
    for i:=nless1 step -1 until 1 do
      begin iless1:=i-1; sum:=0.0; ri:=r[i];
      for j:=2 step 1 until lim do sum:=sum+a[ri,j]*b[r[iless1+j]];
      b[ri]:=b[ri]-sum; if lim<=m then lim:=lim+1
    end;
  for i:=1 step 1 until n do a[i,1]:=b[r[i]];
  for i:=1 step 1 until n do b[i]:=a[i,1]
end bandmx;

```

LE 22

REARRANGE PERMUTED TRIANGULAR MATRIX

procedure MOVEU(a)to:(b)order:(n)pivots:(r); value n;

integer n; array a,b; integer array r;

comment HUCC LIBRARY PROCEDURE LE 22

Author : J. Boothroyd

LIBRARY PROCEDURE LE 10 decomposes a matrix a into quasi triangular interlocking matrices L (a unit lower triangle) and U, an upper triangle which occupy the array a. This procedure separates L and U, replacing the elements of U formerly in a by zeroes and rearranging U so that it is a proper upper triangular matrix in array b. The permuted 'diagonal' elements of L are set equal to one;

LE 23

procedure MOVEL(a) to: (b) order: (n) pivots : (r); value n;

integer n; array a,b; integer array r;

comment HUCC LIBRARY PROCEDURE LE 23

Author : J. Boothroyd

LIBRARY PROCEDURE LE 03 decomposes a matrix a into quasi triangular interlocking matrices L, a lower triangle and U, a unit upper triangle, which together occupy the array a. This procedure separates L and U, replacing the elements of L formerly in array a by zeroes and rearranging L so that it becomes a proper triangular matrix in array b. The permuted 'diagonal' elements of U are at equal to one;

```

procedure MOVEU(a)to:(b)order:(n)pivots:(r);  value n;
integer n;  array a,b;  integer array r;
begin integer i,ri;
  for i:=1 step 1 until n do
    begin ri:=r[i];
      b[i,i]:=a[ri,i];  a[ri,i]:=1.0;
      for j:=i+1 step 1 until n do
        begin b[i,j]:=a[ri,j];  b[j,i]:=a[ri,j]:=0.0 end
    end
  end
endMOVE;

```

```

procedure MOVEL(a)to:(b)order:(n)pivots:(r);  value n;
integer n;  array a,b;  integer array r;
begin integer i,ri;
  for i:=1 step 1 until n do
    begin ri:=r[i];
      b[i,i]:=a[ri,i];  a[ri,i]:=1.0;
      for j:=i-1 step -1 until 1 do
        begin b[i,j]:=a[ri,j];  b[j,i]:=a[ri,j]:=0.0 end
    end
  end
end MOVE;

```

DECOMPOSE SYMMETRIC POSITIVE DEFINITE MATRIX

```
procedure SYMDET(a,n,symdet,fail); value n;
integer n; real symdet; array a; label fail;
```

comment HUCC LIBRARY PROCEDURE LE 24:

AUTHOR J. BOOTHROYD:

Procedure choleski (LE 16) performs the symmetric decomposition of a symmetric matrix stored in conventional form. This procedure performs an in-situ decomposition of the upper triangle of a symmetric matrix stored in vector form, permitting the solution, in Algol, of symmetric equations of order 100 in a 503 with 8K store. The time is approximately 90 seconds.

The upper triangle of $A[1:n,1:n]$ is stored in $a[1:n*(n+1)\text{div}2]$ by columns.

i.e. The upper triangle of

$\begin{matrix} A_{11} & A_{12} & A_{13} & A_{14} \\ A_{21} & A_{22} & A_{23} & A_{24} \\ A_{31} & A_{32} & A_{33} & A_{34} \\ A_{41} & A_{42} & A_{43} & A_{44} \end{matrix}$

$\begin{matrix} A_{11} & A_{12} & A_{13} & A_{14} \\ A_{21} & A_{22} & A_{23} & A_{24} \\ A_{31} & A_{32} & A_{33} & A_{34} \\ A_{41} & A_{42} & A_{43} & A_{44} \end{matrix}$

is stored as :-

a	A11	A12	A22	A13	A23	A33	A14	A24	A34	A44
---	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

and element $A_{i,j}$ is referenced by $a[i+j*(j-1)\text{div}2]$.

To avoid repeated evaluation of subscripts of this form it is preferable to set up a column index vector $c[1:n]$ such that $c[j]:=j*(j-1)\text{div}2$, permitting references to $A_{i,j}$ as $a[i+c[j]]$. This procedure assumes the existence of such a global vector which is best initialised by a recursion relation before entry to the procedure using :-

$r:=0$; for $j:=1$ step 1 until n do begin $c[j]:=r$; $r:=r+j$ end

If the procedure is non-positive-definite exit to the label fail occurs. On normal exit from the procedure symdet yields the determinant of the matrix A;

FORWARD AND BACK SOLUTION Ax=b, SYMMETRIC A

procedure SYMSOL(a,b,n); value n; integer n; array a,b;

comment HUCC LIBRARY PROCEDURE LE 25:

AUTHOR J. BOOTHROYD:

Solves the equation $Ax=b$ for symmetric $A[1:n,1:n]$ by the successive operations $Lf=b$, $L^T x=f$ where $A=LL^T$. The triangular matrix L^T should occupy the vector $a[1:n*(n-1)\text{div}2]$ using column storage, as it will if generated by LE 24, in the description of which procedure find details of the method of storage. The procedure assumes the existence of a global array $c[1:n]$ initialised before entry to the procedure so that $c[j]:=j*(j-1)\text{div}2$, permitting access to an element $L^T_{i,j}$ by reference to $a[i+c[j]]$. On exit from the procedure the elements of the solution vector occupy $b[1:n]$;

```

procedure SYMDET(a,n,symdet,fail); value n; integer n; real symdet; array a; label fail;
begin
    integer i,j,k,iless1,ci,ii,cj,ij;
    real det,aai,aki,aij;
    det:=1.0;
    for i:=1 step 1 until n do
        begin iless1:=i-1; ci:=c[i]; ii:=i+ci; aii:=a[ii];
        for k:=1 step 1 until iless1 do
            begin aki:=a[k+ci]; aii:=aai-aki*aki end;
            if aii<0.0 then goto fail;
            det:=det*aii;
            aii:=a[ii]:=sqrt(aii);
            for j:=i+1 step 1 until n do
                begin cj:=c[j]; ij:=i+cj; aij:=a[ij];
                for k:=1 step 1 until iless1 do aij:=aij-a[k+ci]*a[k+cj];
                a[ij]:=aij/aii
                end j
        end i;
    symdet:=det
end SYMDET;

```

LE 24

```

procedure SYMSOL(a,b,n); value n; integer n; real array a,b;
begin
    integer i,j,jless1,cj; real bi,bj;
    for j:=1 step 1 until n do
        begin bj:=b[j]; jless1:=j-1; cj:=c[j];
        for i:=1 step 1 until jless1 do bj:=bj-a[i+cj]*b[i];
        b[j]:=bj/a[j+cj]
        end;
    for i:=n step -1 until 1 do
        begin bi:=b[i];
        for j:=i+1 step 1 until n do bi:=bi-a[i+c[j]]*b[j];
        b[i]:=bi/a[i+c[i]]
        end
end SYMSOL;

```

LE 25

DECOMPOSE BAND MATRIX WITH PIVOTING

```
procedure pivrband(a,n,m,r,eps,singular); value n,m,eps;
integer n,m; real eps; array a; label singular; integer array r;
```

comment HUCC LIBRARY PROCEDURE LE 26:

AUTHOR J. BOOTHROYD:

Decomposes the n^{th} order bandmatrix of width m stored in columns 1:m of $a[1:n, 1:m+mdiv2]$ into a lower 'triangle' L and a strictly upper 'triangle' U. At entry to the procedure the elements of band matrix $A[1:n, 1:n]$ must occupy the first m columns of a in accordance with the rules:-

$A[i,j]$ occupies $a[i,j]$ $1 \leq i < (m+1)div2$

$A[i,j]$ occupies $a[i, j-i+(m+1)div2]$ $(m+1)div2 < i \leq n$

Diagrammatically :-

A11 A12 A13 0 0 0	A11 A12 A13 0 0 0 0
A21 A22 A23 A24 0 0	A21 A22 A23 A24 0 0 0
A31 A32 A33 A34 A35 0	A31 A32 A33 A34 A35 0 0
0 A42 A43 A44 A45 A46	→ A42 A43 A44 A45 A46 0 0
0 0 A53 A54 A55 A56	A53 A54 A55 A56 0 0 0
0 0 0 A64 A65 A66	A64 A65 A66 0 0 0 0

Array $A[1:n, 1:n]$

Array $a[1:n, 1:m+mdiv2]$

UNUSED ELEMENTS OF MATRIX a should be set to zero.

On exit from the procedure the elements of L occupy column 1 and the last $m div 2$ columns of a. The elements of U (omitting the unit diagonal elements) occupy the remainder of matrix a. The procedure utilises column pivoting without row interchanges with two consequences (a) the algorithm is fast and (b) the storage scheme for elements of L is complicated. All information relating to the succession of pivotal rows is however retained in the index vector $r[1:n]$ for subsequent use by LE 27. If the matrix is singular, eps is the tolerance, the procedure exits via the label parameter;

```

procedure pivlrband(a,n,m,r,eps,singular); value n,m,eps; integer n,m;
real eps; label singular; array a; integer array r;
begin
    integer i,j,k,p,ri,rk,piv,nless1,lim; real pivot,ark1;
    for i:=1 step 1 until n do r[i]:=i;
    lim:=(m+1) div 2; nless1:=n-1;
    for i:=1 step 1 until nless1 do
        begin piv:=i; ri:=r[i]; pivot:=a[ri,1];
        for k:=i+1 step 1 until lim do
            if abs(a[r[k],1])>abs(pivot) then begin piv:=k; pivot:=a[r[k],1] end;
            if abs(pivot)<eps then begin print punch(3),££1?MATRIX SINGULAR?; goto singular end;
            if piv≠i then begin ri:=r[piv]; r[piv]:=r[i]; r[i]:=ri end;
            for j:= 2 step 1 until m do a[ri,j]:=a[ri,j]/pivot;
            p:=m;
            for k:=i+1 step 1 until lim do
                begin p:=p+1; rk:=r[k]; ark1:=a[ri,p]:=a[rk,1];
                for j:=2 step 1 until m do a[rk,j-1]:=a[rk,j]-ark1*a[ri,j];
                a[rk,m]:=0.0
            end;
            if lim≠n then lim:=lim+1
        end;
    end pivlrband;

```

FORWARD AND BACK SOLUTION Ax=b, BANDMATRICES

```
procedure pivsolband(a,b,n,m,r); value m,n; integer m,n;  
array a,b; integer array r;
```

comment HUCC LIBRARY PROCEDURE LE 27:

AUTHOR J. BOOTHROYD:

Solves the equation $Ax=b$ where A is a bandmatrix of order n , width m by the successive operations $Lf=b$, $Ux=f$ where the 'triangular' matrices L and U , such that $LU=A$, occupy the array $a[1:n,1:m\text{div}2]$ following the use of the decomposition procedure LE 26. $r[1:n]$ is a pivotal index vector whose elements are generated by LE 26. These are subsequently used by LE 27 to effect the correct sequence of operations in the forward solution and back substitution;

```

procedure pivsolband(a,b,n,m,r);  value m,n;  array a,b;  integer array r;  integer n,m;
begin
    integer i,j,k,ri,rk,lim;  real bri;
    lim:=(m+1) div 2;
    begin integer array row[1:n];  switch s:=scan;
        for i:=1 step 1 until n do row[i]:=i;
        for i:=1 step 1 until n do
            begin ri:=r[i];  j:=i;
                if ri#row[i] then begin scan:  j:=j+1;  if row[j]#ri then goto scan;
                    row[j]:=row[i];  row[i]:=ri
                end;
                bri:=b[ri]:=b[ri]/a[ri,1];
                j:=m;
                for k:=i+1 step 1 until lim do
                    begin j:=j+1;  rk:=row[k];  b[rk]:=b[rk]-bri*a[ri,j]  end;
                    if lim#n then lim:=lim+1
                end
            end forward solution;
            begin real array bi[1:n];
                bi[n]:=b[r[n]];
                for i:=n-1 stop -1 until 1 do
                    begin ri:=r[i];  bri:=b[ri];
                        k:=1;
                        for j:=i+1 step 1 until lim do
                            begin k:=k+1;  bri:=bri-b[r[j]]*a[ri,k]  end;
                            bi[i]:=b[ri]:=bri;
                            if k=m then lim:=lim-1
                        end;
                    for i:=1 stop 1 until n do b[i]:=bi[i]
            end backsubstitution
    end pivsolband;

```

LE 28

DECOMPOSE SYMMETRIC POSITIVE DEFINITE MATRIX

procedure SYMDET(a,n,symdet,fail); value n; integer n;
real symdet; array a; label fail;

comment HUCC LIBRARY PROCEDURE LE 28:

AUTHOR J. BOOTHROYD:

Performs the same function as LE 24, using the same method of storing the elements of the upper triangle of $A[1:n,1:n]$ as elements of $a[1:n*(n+1)\text{div}2]$. This procedure does not, however, use the global addressing vector $c[1:n]$ necessary with LE 24. It is approximately 7% faster than LE 24.

LE 29

FORWARD AND BACKWARD SOLUTION OF $Ax=b$, SYMMETRIC

procedure SYMSOL (a,b,n); value n; integer n; array a,b;

comment HUCC LIBRARY PROCEDURE LE 29:

AUTHOR J. BOOTHROYD:

Performs the same function as LE 25 but avoids the need for the addressing vector $c[1:n]$;

```

procedure SYMDET(a,n,symdet,fail); value n; integer n; real symdet; array a; label fail;
begin
    integer i,j,ki,ii,k1,kiless1,ij,kj;
    real det,aaii,aki,aij;
    det:=1.0; ii:=1;
    for i:=1 step 1 until n do
        begin aii:=a[ii]; k1:=ii-i+1; kiless1:=ii-1;
        for ki:=k1 step 1 until kiless1 do
            begin aki:=a[ki]; aii:=aii-aki*aki end;
            if aii<0 then goto fail;
            det:=det*aii;
            aii:=a[ii]:=sqrt(aii);
            ij:=ii+i;
            for j:=i+1 step 1 until n do
                begin aij:=a[ij]; kj:=ij-i+1;
                for ki:=k1 step 1 until kiless1 do
                    begin aij:=aij-a[ki]*a[kj]; kj:=kj+1 end;
                    a[ij]:=aij/aii; ij:=ij+j
                end;
                ii:=ii+i+1;
            end;
            symdet:=det
    end SYMDET;;

```

LE 28

```

procedure SYMSOL(a,b,n); value n; integer n; array a,b;
begin integer i,j,k,ii; real sum;
ii:=1;
for i:=1 step 1 until n do
begin k:=ii-1; sum:=b[i];
for j:=i-1 step -1 until 1 do
begin sum:=sum-b[j]*a[k]; k:=k-1 end;
b[i]:=sum/a[ii]; ii:=ii+i+1
end forward solution;
k:=ii-1; ii:=ii-n-1;
for i:=n step -1 until 1 do
begin sum:=b[i];
for j:=i+1 step 1 until n do
begin sum:=sum-b[j]*a[k]; k:=k+j end;
b[i]:=sum/a[ii]; k:=ii-1; ii:=ii-i
end back substitution
end SYMSOL;

```

LE 29

REDUCE MATRIX TO UPPER HESSENBERG FORM

procedure hessenberg(a,n); value n; integer n;

comment HUCC LIBRARY PROCEDURE LE30:

AUTHOR : J. BOOTHROYD:

Transforms a real matrix a[1:n,1:n] to upper Hessenberg form by a sequence of elementary similarity transformations. The transformation is performed using a pivoting strategy which involves row and column interchanges. The method is similar to Gaussian elimination. For a compact CROUT-like method which converts a matrix to lower Hessenberg form see procedure TRINGLE in Parlett's paper "Languerre's Method Applied to the Matrix Eigenvalue Problem" Math. Comp v18 1964 469 - 485;

```

procedure hessenberg(a,n);  value n;  integer n;  array a;
begin integer i,j,k,iless1,iplus1,piv,nless1;  real pivot,mk,aki,temp;
array m[3:n];  boolean nonzero;  switch s:=pivsearch;
pivot:=0.0;  nless1:=n-1;
for i:=1 step 1 until nless1 do
begin iplus1:=i+1;
if pivot=0.0 then goto pivsearch;
if i#piv then
begin for j:=i+1 step 1 until n do begin temp:=a[i,j];  a[i,j]:=a[piv,j];  a[piv,j]:=temp end;
for k:=1 step 1 until n do begin temp:=a[k,i];  a[k,i]:=a[k,piv];  a[k,piv]:=temp end
end;
nonzero:= false;
for k:=iplus1 step 1 until n do
begin mk:=m[k]:=a[k,iless1]/pivot;
if mk#0.0 then begin nonzero:= true;  a[k,iless1]:=0.0 end
end;
if nonzero then
begin for k:=iplus1 step 1 until n do
begin mk:=m[k];
for j:=i step 1 until n do a[k,j]:=a[k,j]-mk*a[i,j]
end;
pivot:=iplus1;  pivot:=0.0;
for k:=1 step 1 until n do
begin aki:=a[k,i];
for j:=iplus1 step 1 until n do aki:=aki+m[j]*a[k,j];
a[k,i]:=aki;
if k>i then
begin if abs(aki)>abs(pivot) then begin piv:=k;  pivot:=aki end end
end;
end;
else
pivsearch: begin piv:=iplus1;  pivot:=a[piv,i];
for k:=i+2 step 1 until n do
begin aki:=a[k,i];
if abs(aki)>abs(pivot) then begin piv:=k;  pivot:=aki end
end;
end;
iless1:=i
end i
end hessenberg;

```

TRIDIAGONALIZE SYMMETRIC MATRIX GIVENS METHOD

procedure givens(a,s,n); value n; integer n; array a,s;

comment HUCC LIBRARY PROCEDURE LE31:

AUTHOR : J. BOOTHROYD:

Performs a similarity transformation of a full symmetric matrix $a[1:n,1:n]$ to tridiagonal form by GIVENS Method.

The orthogonal rotation matrices used have the typical form:

$$\begin{array}{cccc} P_r = & 1 & 0 & 0 \\ & 0 & c & s \\ & 0 & 0 & 1 \\ & 0 & s & 0 \end{array}$$

so that $P_r = P_r^{-1} = P_r^t$.

The continued product P_1, P_2, P_3, \dots of transformation matrices is finally available in array $s[1:n,1:n]$;

```

procedure givens(a,s,n);  value n;  integer n;  array a,s;
begin integer i,p,q,nless1,k;
  real akp,akq,app,apq,aqq,aip,aiq,fac,cos,sin,cs;
  for i:=1 step 1 until n do
  begin s[i,i]:=1.0;
    for j:=i+1 step 1 until n do s[i,j]:=s[j,i]:=0.0
  end;
  nless1:=n-1;
  for p:=2 step 1 until nless1 do
  begin k:=p-1;  akp:=a[k,p];
    for q:=p+1 step 1 until n do
    begin akq:=a[k,q];
      apq:=a[p,q];  app:=a[p,p];  aqq:=a[q,q];
      fac:=sqrt(akp*akp+akq*akq);
      if fac#0.0 then
      begin cos:=akp/fac;  sin:=akq/fac;
        for i:=1 step 1 until n do
        begin if i>p then
          begin aip:=a[i,p];  aiq:=a[i,q];
            a[i,p]:=a[p,i]:=cos*aip+sin*aiq;
            a[i,q]:=a[q,i]:=sin*aip-cos*aiq
          end;
          aip:=s[i,p];  aiq:=s[i,q];
          s[i,p]:=cos*aip+sin*aiq;
          s[i,q]:=sin*aip-cos*aiq
        end;
        akp:=fac;  a[k,q]:=a[q,k]:=0.0;
        cs:=cos*sin;  fac:=(cs+cs)*apq;
        cos:=cos*cos;  sin:=sin*sin;
        a[p,p]:=cos*app+sin*aqq+fac;
        a[q,q]:=sin*app+cos*aqq-fac;
        a[p,q]:=a[q,p]:=-(app-aqq)*cs-apq*(cos-sin)
      end;
    end;
    a[k,p]:=a[p,k]:=akp
  end;
end givens;

```

TRIDIAGONALIZE SYMMETRIC MATRIX GIVENS METHOD

procedure vecgivens(a,s,n); value n; integer n; array a,s;

comment HUCC LIBRARY PROCEDURE LE32:

AUTHOR : J. BOOTHROYD:

As for LE31 except that only the upper half of the symmetric matrix is stored, by columns, in the vector a[1:n*(n+1)div2]. The continued product P_1, P_2, P_3, \dots of orthogonal transformation matrices is available in s[1:n,1:n]. If this procedure is used as an alternative to LL09 and prior to LL10 and LL11 it will be necessary to transfer the diagonal and codiagonal elements of array a to vectors c[1:n] and b[1:n] respectively. This may be done by:-

```

jless1:=k:=1;  c[1]:=a[1];
for j:= 2 step 1 until n do
  begin k:=k+j;  c[j]:=a[k];
  b[jless1]:=a[k-1];  jless1:=j
  end;
b[n]:=0.0;

```

The eigenvectors of the tridiagonal system should be premultiplied by s[1:n,1:n] to yield the eigenvectors of a;

```

procedure vecgivens(a,s,n);  value n;  integer n;  array a,s;
begin integer i,p,q,nless1,k,cp,cq,ci,ip,iq;  integer array c[1:n];
  real akp,app,akq,apq,aqq,aip,aiq,fac,cos,sin,cs;
  for i:=1 step 1 until n do
    begin s[i,i]:=1.0
      for j:=i+1 step 1 until n do s[i,j]:=s[j,i]:=0.0
    end;
  p:=0;  for i:=1 step 1 until n do begin c[i]:=p;  p:=i+p  end;
  nless1:=n-1;
  for p:=2 step 1 until nless1 do
    begin k:=p-1;  cp:=c[p];  akp:=a[k+cp];
      for q:=p+1 step 1 until n do
        begin cq:=c[q];  akq:=a[k+cq];
          apq:=a[p+cq];  app:=a[p+cp];  aqq:=a[q+cq];
          fac:=sqrt(akp*akp+akq*akq);
          if fac#0.0 then
            begin cos:=akp/fac;  sin:=akq/fac;
              for i:=1 step 1 until n do
                begin if i>p then
                  begin ci:=c[i];  ip:=p+ci;
                    iq:= if i>q then q+ci else i+cq;
                    aip:=a[ip];  aiq:=a[iq];
                    a[ip]:=cos*aip+sin*aiq;
                    a[iq]:=sin*aip-cos*aiq
                  end;
                  aip:=s[i,p];  aiq:=s[i,q];
                  s[i,p]:=cos*aip+sin*aiq;
                  s[i,q]:=sin*aip-cos*aiq
                end i;
              akp:=fac;  a[k+cq]:=0.0;
              cs:=cos*sin;  fac:=(cs+cs)*apq;
              cos:=cos*cos;  sin:=sin*sin;
              a[p+cp]:=cos*app+sin*aqq+fac;
              a[q+cq]:=sin*app+cos*aqq-fac;
              a[p+cq]:=(app-aqq)*cs-apq*(cos-sin)
            end
          end q;
        a[k+cp]:=akp
      end p
    end vecgivens;

```

EIGENVALUES AND EIGENVECTORS OF SYMMETRIC MATRIX

```
procedure jacobi(a,s,n,rho);  value n,rho;  integer n;  real rho;
    array a,s;
comment    HUCC LIBRARY PROCEDURE LL01:
AUTHOR      ACM 85      :
```

The procedure finds all eigenvalues and eigenvectors of a real symmetric matrix by Jacobi's method as described in "Mathematical Methods for Digital Computers" edited by Ralston and Wilf. The eigenvectors of $a[1:n,1:n]$ are built up in $s[1:n,1:n]$ the k th eigenvector occupying column k . The corresponding eigenvalue occupies element $a[k,k]$ of the original matrix. rho is the precision tolerance for the process which is terminated when, for every off diagonal element $a[i,j]$, $abs(a[i,j]) < rho \times norm1/n$ where norm1 is the square root of the sum of the squares of the off diagonal elements of a ;

```
procedure jacobi(a,s,n,rho);
value n,rho;
integer n; real rho; array a,s;
begin real norm1,norm2,thr,mu,omega,sint,cost,int1,v1,v2,v3;
integer i,j,p,q,ind; switch ss:=main,main1;

for i:=1 step 1 until n do
for j:=1 step 1 until i do
if i=j then s[i,j]:=1.0 else s[i,j]:=s[j,i]:=0;

int1:=0;
for i:=2 step 1 until n do
for j:=1 step 1 until i-1 do int1:=int1+2*a[i,j]↑2;
norm1:=sqrt(int1); norm2:=(rho/n)*norm1;
thr:=norm1; ind:=0;

main: thr:=thr/n;
main1: for q:=2 step 1 until n do
for p:=1 step 1 until q-1 do
if abs(a[p,q]) > thr then
begin ind:=1; v1:=a[p,p]; v2:=a[p,q]; v3:=a[q,q]; mu:=.5*(v1-v3);

omega:=if mu=0.0 then -1.0 else -sign(mu)*v2/sqrt(v2*v2+mu*mu);
sint:=omega/sqrt(2*(1+sqrt(1-omega*omega)));
cost:=sqrt(1-sint*sint);
for i:=1 step 1 until n do
begin int1:=a[i,p]*cost-a[i,q]*sint;
a[i,q]:=a[i,p]*sint+a[i,q]*cost; a[i,p]:=int1;
int1:=s[i,p]*cost-s[i,q]*sint;
s[i,q]:=s[i,p]*sint+s[i,q]*cost; s[i,p]:=int1
end;
end;
```

```
for i:=1 step 1 until n do
  begin a[p,i]:=a[i,p]; a[q,i]:=a[i,q]
  end;

a[p,p]:=v1*cost*cost+v3*sint*sint-2*v2*sint*cost;
a[q,q]:=v1*sint*sint+v3*cost*cost+2*v2*sint*cost;
a[p,q]:=a[q,p]:=(v1-v3)*sint*cost+v2*(cost*cost-sint*sint)
end;

if ind=1 then
  begin ind:=0; go to main1
  end
else if thr > norm2 then go to main
end jacobi;
```

SOLUTION OF THE EIGENPROBLEM $(A-\lambda B)x=0$

```
procedure eigensolve (A,B,n,eigen,eps,nonposdef); value n,eps;
integer n; real eps; array A,B,eigen; label nonposdef;
```

comment HUCC LIBRARY PROCEDURE LL 02

Author : J. Boothroyd.

Solves the equation $(A-\lambda B)x=0$ for symmetric $A, B[1:n, 1:n]$ provided one of these is positive definite. The equation is transformed into $(C-\lambda I)y=0$ where $C=L^{-1}A(L^T)^{-1}$ is symmetric and $y=L^T x$. If B is non positive definite the Choleski decomposition fails in the attempted evaluation of a square root of a negative value. In this event the original equation is rearranged as $(B-\lambda A)x=0$ for which the eigenvalues are the reciprocals of $(A-\lambda B)x=0$ and the eigenvectors are unchanged. Failure of this second attempt causes exit to the label nonposdef.

At a successful exit the eigenvalues are in $eigen[1:n]$ and the vectors occupy A . This procedure uses MS 01, the revised issues of LE 16, LE 17, together with LE 18 and LL 01. The parameter eps is a precision tolerance used by LL 01 for details of which see the appropriate commentary;

```

procedure eigensolve(A,B,n,eigen,eps,nonposdef); value n,eps; integer n; real eps; array A,B,eigen; label nonposdef;
begin integer i,j; boolean recip; array C[1:n,1:n]; switch s:=newtry,swap,ok;

    comment temporary storage of the main diagonal of B;
    for i:= 1 step 1 until n do eigen[i]:=B[i,i];
    recip:=false; comment assumes B is positive definite;
newtry: choleski(B,n,swap);
    comment clear upper triangle of B;
    for i:= 1 step 1 until n do for j:= i+1 step 1 until n do B[i,j]:= 0.0;
    goto ok;
swap: if recip then goto nonposdef; recip:=true;
    comment B was nonpositive definite. Swap A,B and try again;
    for i:= 1 step 1 until n do
begin B[i,i]:=A[i,i]; A[i,i]:= eigen[i];
    for j:= i+1 step 1 until n do
        begin B[j,i]:=A[i,j]; A[i,j]:=A[j,i]:=B[i,j] end
end;
    goto newtry;
ok: linv(B,n); comment forms L-1 in B;
mxmult(B,A,C,n,n,n); comment C:=L-1A;
comment transpose L-1 and clear lower triangle of B;
for i:= 1 step 1 until n do
for j:= i+1 step 1 until n do begin B[i,j]:=B[j,i]; B[j,i]:=0.0 end;
mxmult(C,B,A,n,n,n); comment A:= C(Lt)-1 = L-1A(Lt)-1;
jacobi(A,C,n,eps); comment ACM85. Eigenvalues on A diagonal, y vectors in C;
for i:= 1 step 1 until n do eigen[i]:= if recip then 1.0/A[i,i] else A[i,i];
mxmult(B,C,A,n,n,n); comment x vectors = (Lt)-1y now in A;
end eigensolve;

```

SOLVE SYMMETRIC EIGENSYSTEM

These procedures are taken from the LINEAR ALGEBRA Handbook Series of Numerische Mathematique 4 pp 354 - 376 (1962). Their author, J.H. Wilkinson of the National Physical Laboratory, England, is the recognised authority on this subject and these procedures are included in the HUCC Library as examples of classical algorithms in this field.

The following program was used to test four of the five procedures and indicates how to use the various parameters.

```

TEST HOUSEHOLDER;
begin comment procedure declarations;
begin integer n; read n;
begin integer i,j,mr,m1,t,m1; real e,gamma,norm;
array z,a[1:n,1:n],c,b,w[1:n];
for i:=1 step 1 until n do
for j:=1 step 1 until i do read a[i,j];
householder(a,n,c,b);
tridibisection2(c,b,n,-20,n,0,30,10-16,w,norm,m1);
tridiinverseiteration(c,b,n,w,norm,m1,10-8,z);
backtransformation(a,b,z,n,m1);
print ff1?EIGENVALUES12?;
sameline; scaled(9);
for i:=1 step 1 until m1 do print w[i];
print ff15?EIGENVECTORS1?;
for i:=1 step 1 until n do
begin print ff12?;
for j:=1 step 1 until n do print z[j,i],ffs2?;
end i
end
end;

```

LL03

HOUSEHOLDER TRIDIAGONALIZATION, FULL MATRIX

procedure householder tridiagonalization(a,n,c,b);

value n; integer n; array a,b,c;

comment HUCC LIBRARY PROCEDURE LL03:

AUTHOR : J.H. WILKINSON:

The symmetric matrix a[1:n,1:n] is reduced to a symmetric tridiagonal matrix using only the lower triangle of a. The diagonal elements of the tridiagonal matrix are stored in c[1:n], the subdiagonal elements in b[1:n] with b[n]=0.

Sufficient details of the transformation are retained in a and b to enable the eigenvectors of a to be formed from the tridiagonal eigenvectors using procedure backtransformation;

LL04

EIGENVALUES OF SYMMETRIC TRIDIAGONAL MATRIX 1

procedure tridibisection 1(c,b,n,gu,go,t,gamma,w,norm,ml);

value n,gu,go,t,gamma; integer n,t,ml;

real gu,go,gamma,norm; array c,b,w;

comment HUCC LIBRARY PROCEDURE LL04:

AUTHOR : J.H. WILKINSON:

c[1:n] and b[1:n] are the diagonal and sub-diagonal of a symmetric tridiagonal matrix. The procedure determines the number ml of eigenvalues lying between gu and go and computes these eigenvalues in w[1:n] in decreasing order of magnitude by the method of bisection. t is the number of bisection steps (30 to 35 is appropriate on a 503), norm is the infinity norm of the matrix and gamma is the square of the relative machine precision (10^{-16} for a 503);

LL05

EIGENVALUES OF SYMMETRIC TRIDIAGONAL MATRIX 2

procedure tridibisection 2(c,b,n,e,mr,ml,t, gamma,w,norm,ml);
value n,e,mr,ml,t, gamma; real e, gamma, norm;
integer n, mr, ml, t, ml; array c,b,w;

comment HUCC LIBRARY PROCEDURE LL05:

AUTHOR : J.H. WILKINSON:

c[1:n] and b[1:n] are the diagonal and subdiagonal of a symmetric tridiagonal matrix. The mr eigenvalues to the left of e and the ml eigenvalues to the right of e are computed and stored in w[1:n]. t is the number of bisection steps(30to35 for a 503), norm is the infinity norm of the matrix and gamma is the square of the relative machine precision (10^{-16} for 503);

LL06

EIGENVECTORS OF TRIDIAGONAL MATRIX

procedure tridiinverseiteration(c,b,n,w,norm,w1,macheps,z);
value n,ml,norm,macheps; integer n,ml; real norm,macheps;
array c,b,w,z;

comment HUCC LIBRARY PROCEDURE LL06:

AUTHOR : J.H. WILKINSON:

c[1:n] and b[1:n] are the diagonal and subdiagonal of a symmetric tridiagonal matrix. w[1:n] contains the ml eigenvalues from w[1] to w[ml]. norm is the infinity norm of the matrix and macheps is the relative machine precision (10^{-8} on a 503). The ml eigenvectors are computed and stored in z[1:n,1:n], element z[i,j] being the jth component of the eigenvector corresponding to the eigenvalue w[i];

LL07

EIGENVECTORS OF SYMMETRIC MATRIX

procedure backtransformation(a,b,z,n,ml); value n,ml; integer n,ml; array a,b,z;

comment HUCC LIBRARY PROCEDURE LL07:

AUTHOR : J.H. WILKINSON:

b[1:n] is the subdiagonal of a symmetric tridiagonal matrix whose eigenvectors occupy z[1:n,1:n], by rows. The eigenvectors of the original matrix a[1:n,1:n] are derived and overwritten on array z;

```

procedure householder tridiagonalisation(a,n)result:(c,b);
value n; integer n; array a,b,c;
begin integer j,i,k; real ai,sigma,h,bj,bigk,bi; array q[1:n-1];
for i:=n step -1 until 3 do
begin sigma:=0;
  for k:=1 step 1 until i-1 do
    sigma:=sigma+a[i,k]*a[i,k];
    ai:=a[i,i-1];
    if ai > 0 then bi:=-sqrt(sigma) else bi:=sqrt(sigma);
    b[i-1]:=bi;
    if bi#0 then
    begin h:=sigma-ai*bi;
    a[i,i-1]:=ai-bi;
      for j:=i-1 step -1 until 1 do
      begin bj:=0;
        for k:=i-1 step -1 until j do
          bj:=bj+a[k,j]*a[i,k];
        for k:=j-1 step -1 until 1 do
          bj:=bj+a[j,k]*a[i,k];
        q[j]:=bj/h
      end j;
      bigk:=0;
      for j:=i-1 step -1 until 1 do
      begin bigk:=bigk+a[i,j]*q[j];
      bigk:=bigk/(2*h);
      for j:=i-1 step -1 until 1 do
        q[j]:=q[j]-bigk*a[i,j];
      for j:=i-1 step -1 until 1 do
      begin for k:=j step -1 until 1 do
        a[j,k]:=a[j,k]-a[i,j]*q[k]-a[i,k]*q[j];
      end j
      end
    end i;
    for i:=n step -1 until 1 do
    begin c[i]:=a[i,i];
    b[1]:=a[2,1];
    b[n]:=0
  end;

```

```

procedure tridibisection1(c,b,n,gu,go,t,gamma) result:(w,norm,m1);
value n,gu,go,t,gamma;
integer n,t,m1;
real gu,go,gamma,norm;
array c,b,w;
begin integer i,j,k,a1,a2,d;
    real l,g,h,lambda,p1,q1,y; switch ss:=noeigenvalue;
    array p[1:n];
    procedure sturms sequence;
    begin p1:=0; q1:=1; a1:=0;
        for i:=1 step 1 until n do
            begin y:=(c[i]-lambda)*q1-p[i]*p1;
                p1:=q1; q1:=y;
                if p1 ≥ 0 ≤ q1 ≥ 0 then a1:=a1+1
            end i;
        if q1=0 and p1>0 then a1:=a1-1
    end;
    if gu>go then
        begin g:=gu; gu:=go; go:=g end;
    norm:=abs(c[1])+abs(b[1]);
    for i:=2 step 1 until n do
        begin l:=abs(b[i-1])+abs(c[i])+abs(b[i]);
            if l>norm then norm:=l
        end;
    for i:=1 step 1 until n-1 do
        begin if b[i]=0 then p[i+1]:=gamma*norm*norm
            else p[i+1]:=b[i]*b[i]
        end;
    p[1]:=0;
    if gu>norm or go<-norm then
        begin m1:=0; goto noeigenvalue end;
    lambda:=gu;

    sturms sequence;
    a2:=a1;
    if q1=0 then a2:=a1+1;
    lambda:=go;
    sturms sequence;
    m1:=a2-a1;
    d:=a1;
    if go>norm then go:=norm;
    if gu<-norm then gu:=-norm;
    for k:=1 step 1 until m1 do
        begin d:=d+1;
            g:=go; h:=gu;
            for j:=1 step 1 until t do
                begin lambda:=(g+h)/2;
                    sturms sequence;
                    if a1 > d then h:=lambda else
                        g:=lambda
                end j;
            w[k]:=(g+h)/2
        end k;
noeigenvalue: end;

```

```

procedure tridibisection 2 (c,b,n,e,mr,ml,t,gamma) result:(w,norm,m1);
value n,e,mr,ml,t,gamma;
integer n,mr,ml,t,m1;
real e,gamma,norm;
array c,b,w;
begin      integer d1,d2,i,j,k,a1,d;
            real l,g,h,lambda,p1,q1,y;
            array p[1:n];
            procedure sturms sequence;
            begin      p1:=0; q1:=1; a1:=0;
                        for i:=1 step 1 until n do
                        begin      y:=(c[i]-lambda)*q1-p[i]*p1;
                                    p1:=q1; q1:=y;
                                    if p1 > 0 = q1 > 0
                                    then a1:=a1+1;
                        end i;
                        if q1=0 and p1>0 then a1:=a1-1
                        end;
            norm:=abs(c[1])+abs(b[1]);
            for i:=2 step 1 until n do
            begin      l:=abs(b[i-1])+abs(c[i])+abs(b[i]);
                        if l>norm then norm:=l
            end;
            for i:=1 step 1 until n-1 do
            begin      if b[i]=0 then p[i+1]:=gamma*norm*norm
                        else p[i+1]:=b[i]*b[i]
            end;
            p[1]:=0;
            lambda:=e;
            sturms sequence;

            d1:=a1-mr; d2:=a1+ml;
            if d1<0 then d1:=0;
            if d2 > n+1 then d2:=n;
            d:=d1; m1:=0;
            for k:=d1 step 1 until d2-1 do
            begin      d:=d+1;
                        g:=norm; h:=-norm;
                        for j:=1 step 1 until t do
                        begin      lambda:=(g+h)/2;
                                    sturms sequence;
                                    if a1 > d then h:=lambda else g:=lambda
                        end j;
                        m1:=m1+1;
                        w[m1]:=(g+h)/2
            end k
            end;

```

```

procedure tridiinverse iteration(c,b,n,w,norm,m1,macheeps) results:(z);
value n,m1,norm,macheeps; integer n,m1; real norm,macheeps; array c,b,w,z;
begin integer i,j; real bi,bi1,z1,lambda,u,s,v,h,eps,eta;
array m,p,q,r,int[1:n],x[1:n+2];
lambda:=norm; eps:=macheeps*norm;
for j:=1 step 1 until m1 do
begin lambda:=lambda-eps;
if w[j]<lambda then lambda:=w[j];
u:=c[1]-lambda; v:=b[1];
if v=0 then v:=eps;
for i:=1 step 1 until n-1 do
begin bi:=b[i];
if bi=0 then bi:=eps;
bi1:=b[i+1];
if bi1=0 then bi1:=eps;
if abs(bi) > abs(u) then
begin m[i+1]:=u/bi;
if m[i+1]=0 and abs(bi) < eps then m[i+1]:=1;
p[i]:=bi; q[i]:=c[i+1]-lambda;
r[i]:=bi1;
u:=v-m[i+1]*q[i];
v:=-m[i+1]*r[i];
int[i+1]:=1
end
else begin m[i+1]:=bi/u;
p[i]:=u; q[i]:=v;
r[i]:=0;
u:=-[i+1]-lambda-m[i+1]*v;
v:=bi1; int[i+1]:=-1
end
end i;
p[n]:=u; q[n]:=r[n]:=0;
x[n+1]:=x[n+2]:=0; h:=0; eta:=1/n;
for i:=n step -1 until 1 do
begin u:=eta-q[i]*x[i+1]-r[i]*x[i+2];
if p[i]=0 then x[i]:=u/eps
else x[i]:=u/p[i];
h:=h+abs(x[i])
end i;
h:=1/h;
for i:=1 step 1 until n do
x[i]:=x[i]*h;
for i:=2 step 1 until n do
begin if int[i]>0 then
begin u:=x[i-1];
x[i-1]:=x[i];
x[i]:=u-m[i]*x[i-1]
end
else x[i]:=x[i]-m[i]*x[i-1]
end i;
h:=0;
for i:=n step -1 until 1 do
begin u:=x[i]-q[i]*x[i+1]-r[i]*x[i+2];
if p[i]=0 then x[i]:=u/eps
else x[i]:=u/p[i];
h:=h+x[i]*x[i]
end i;
h:=1/sqrt(h);
for i:=1 step 1 until n do
z[j,i]:=x[i]*h
end j
end;

```

LL07

```
procedure backtransformation(a,b,z,n,m1);
value n,m1; integer n,m1; array a,b,z;
begin integer i,j,k; real s; for j:=1 step 1 until m1 do
  for k:=3 step 1 until n do
    if b[k-1]#0 then
      begin s:=0;
        for i:=1 step 1 until k-1 do
          s:=s+a[k,i]*z[j,i];
        s:=s/(b[k-1]*a[k,k-1]);
        for i:=1 step 1 until k-1 do
          z[j,i]:=z[j,i]+s*a[k,i]
      end
    end;
  end;
```

EIGENVALUES AND VECTORS OF SYMMETRIC MATRIX

```
procedure vecjacobi(a,s,n,rho); value n,rho; real rho;  
integer n; array a,s;
```

comment HUCC LIBRARY PROCEDURE LL08:

AUTHOR : J. BOOTHROYD:

A revised version of LL01, jacobi, such that only the upper half of the symmetric matrix is stored in the vector $a[1:n*(n+1)/2]$. Storage is by columns as described in LE24. On exit from the procedure the eigenvalues occupy elements $a[1]$ through $a[n]$ and the eigenvectors occupy the array $s[1:n,1:n]$. On a 20×20 matrix LL01 took 120 seconds, LL08 only 47 seconds. On the same matrix J.H. Wilkinson's procedures (LL03,04,06,07) took 27 seconds and their improved counterparts (LL09,10,11,12) took 23 seconds;

```

procedure vecjacobi(a,s,n,rho); value n,rho; real rho; integer n; array a,s;
begin integer array c[1:n]; integer i,j,ci,cj,p,q,apq,app,aqq,m,mu,lambda,cost,sint,aip,aiq,sip,siq,sincos;
switch ss:=main,main1;
real fac,aij,thr,norm1,norm2,apq,app,aqq,m,mu,lambda,cost,sint,aip,aiq,sip,siq,sincos;
boolean ind;
p:=0; fac:=0.0;
for i:=1 step 1 until n do
begin s[i,i]:=1.0; c[i]:=p; p:=p+i;
for j:=i+1 step 1 until n do s[i,j]:=s[j,i]:=0.0;
end;
for j:=2 step 1 until n do
begin cj:=c[j]; jless1:=j-1;
for i:=1 step 1 until jless1 do
begin aij:=a[i+cj]; fac:=fac+2.0*aij*aij end
end;
thr:=norm1:=sqrt(fac); norm2:=rho*norm1/n;
main: thr:=thr/n;
main1: ind:=false;
for q:=2 step 1 until n do
begin cq:=c[q]; qless1:=q-1;
for p:=1 step 1 until qless1 do
begin apq:=a[p+cq];
if abs(apq) > thr then
begin cp:=c[p]; ind:= true;
app:=a[p+cp]; aqq:=a[q+cq]; m:=app-aqq;
lambda:=sign(m)*apq; mu:=0.5*abs(m);
fac:=0.5/sqrt(lambda*lambda+mu*mu);
cost:=sqrt(0.5+mu*fac); sint:=lambda*fac/cost;
for i:=1 step 1 until n do
begin ci:=c[i];
if i < p then begin ip:=i+cp; iq:=i+cq end
else begin ip:=p+ci;
iq:=if i>q then q+ci else i+cq
end;
aip:=a[ip]; aiq:=a[iq];
sip:=s[i,p]; siq:=s[i,q];
s[i,p]:=cost*sip+sint*siq;
s[i,q]:=sint*sip-cost*siq;
a[ip]:=cost*aip+sint*aiq;
a[iq]:=sint*aip-cost*aiq
end i;
sincos:=sint*cost; fac:=(apq+apq)*sincos;
sint:=sint*sint; cost:=cost*cost;
a[p+cp]:=cost*app+sint*aqq+fac;
a[q+cq]:=sint*app+cost*aqq-fac;
a[p+cq]:=0.0
end
end;
if ind then goto main1 else if thr>norm2 then goto main;
for i:=2 step 1 until n do a[i]:=a[i+c[i]]
end vecjacobi;

```

```

procedure householder(a,n,c,b);  value n;  integer n;  real array a,b,c;
begin integer i,j,k,nless2,iplus1,cj,ck,kcj,cipplus1,nless1;  real s,aij,bi,h,rho,zj;
real array z[1:n];
nless2:=n-2;  nless1:=n-1;
for i:=1 step 1 until nless2 do
begin iplus1:=i+1;  s:=0.0;  k:=cipplus1:=col[iplus1];
for j:=iplus1 step 1 until n do
begin aij:=a[i+k];  k:=k+j;  s:=s+aij*aij end;
aij:=a[i+cipplus1];  bi:=sqrt(s);
if aij > 0.0 then bi:=-bi;  b[i]:=bi;
if bi ≠ 0.0 then
begin h:=s-aij*bi;  a[i+cipplus1]:=aij-bi;
for j:=iplus1 step 1 until n do
begin s:=0.0;  cj:=col[j];
for k:=iplus1 step 1 until n do
begin ck:=col[k];
s:=s+a[i+ck]*a[i+k] if k<j then k+cj else j+ck];
end;
z[j]:=s/h
end;
s:=0.0;  k:=cipplus1;
for j:=iplus1 step 1 until n do begin s:=s+a[i+k]*z[j];  k:=k+j end;
rho:=s/(h+h);  k:=cipplus1;
for j:=iplus1 step 1 until n do begin z[j]:=z[j]-rho*a[i+k];  k:=k+j end;
for j:=iplus1 step 1 until n do
begin cj:=col[j];  aij:=a[i+cj];  zj:=z[j];  ck:=cipplus1;
for k:=iplus1 step 1 until j do
begin kcj:=k+cj;
a[kcj]:=a[kcj]-a[i+ck]*zj-aij*z[k];  ck:=ck+k
end
end;
end;
c[i]:=a[i+col[i]]
end;
cj:=col[n];  c[nless1]:=a[nless1+col[nless1]];
b[nless1]:=a[nless1+cj];  b[n]:=0.0;  c[n]:=a[n+cj]
end householder;

```

LL09 to LL12

SYMMETRIC EIGENSYSTEM HALF MATRIX

These procedures are adapted from LL03 to LL07 to compute the eigenvalues and eigenvectors of an n^{th} order symmetric matrix of which only the upper half is stored by columns in a vector $a[1:n*(n+1)\text{div}2]$. In designing this revision the opportunity has been taken to improve, where possible, the efficiency of the original procedures and rationalise, to some extent, the parameter organisation.

The following program was used to test LL09 to LL12 and illustrates the use of the several parameters:-

```
TEST HOUSEHOLDER;
begin integer n,t; real eps; switch ss:=L;
L: read n,eps,t;
begin integer i,j,k; real norm;
array z[1:n,1:n],a[1:n*(n+1) div 2],c,b,w[1:n]; integer array col[1:n];
comment procedure declarations;
k:=0;
for j:=1 step 1 until n do
begin col[j]:=k;
for i:=1 step 1 until j do read a[i+k];
k:=k+j
end;
householder(a,n,c,b);
eigbise(c,b,n,eps,t,w,norm);
trivectors(c,b,n,w,norm,eps,z);
eigvectors(a,b,z,n);
freepoint(9); sameline;
for j:=1 step 1 until n do
begin print ffl??,w[j],ffl??;
for i:=1 step 1 until n do
begin if i-i div 10*10=1 then print ffl??; print z[i,j] end
end
end;
goto L
end;
```

LL09

HOUSEHOLDER TRIDIAGONALISATION - HALF MATRIX

procedure householder(a,n,c,b); value n; integer n; array a,c,b;

comment HUCC LIBRARY PROCEDURE LL09;

AUTHOR : J. BOOTHROYD:

Transforms an n^{th} order symmetric matrix whose upper half is stored by columns in $a[1:n*(n+1)\text{div}2]$ to tridiagonal form. The diagonal and super-diagonal stored in $c[1:n]$ and $b[1:n]$ respectively (with $b[n]=0$). Sufficient information is retained in a , and b for subsequent use by procedure eigvectors (LL12);

LL10

EIGENVALUES OF TRIDIAGONAL MATRIX

procedure eigbise(c,b,n,eps,t,w,norm); value n,eps,t;

integer n,t; real eps,norm; array c,b,w;

comment HUCC LIBRARY PROCEDURE LL10;

AUTHOR : J. BOOTHROYD:

Evaluates, in decreasing order of magnitude, using the method of bisection, the n eigenvectors of a tridiagonal matrix whose diagonal and co-diagonal occupy $c[1:n]$ and $b[1:n]$ respectively. The results occupy $w[1:n]$. eps is the relative machine precision (10^{-8} on a 503), t is the number of bisection steps (30 to 35 is appropriate) and norm is the computed infinity norm of the matrix;

LL11

EIGENVECTORS OF TRIDIAGONAL MATRIX

procedure trivector(c,b,n,w,norm,eps,z); value n,eps,norm;

integer n; real norm,eps; array c,b,w,z;

comment HUCC LIBRARY PROCEDURE LL11;

AUTHOR : J. BOOTHROYD:

Computes in the columns of $z[1:n,1:n]$ the eigenvectors of a symmetric tridiagonal matrix whose eigenvalues occupy $w[1:n]$ in decreasing order of magnitude. $c[1:n]$ and $b[1:n]$ are the diagonal and co-diagonal of the tridiagonal matrix. eps is the relative machine precision (10^{-8} on a 503), norm is the infinity norm of the matrix;

LL12

EIGENVECTORS OF SYMMETRIC MATRIX (UPPER HALF)

procedure eigvectors(a,b,z,n); value n;

integer n; array a,b,z;

comment HUCC LIBRARY PROCEDURE LL12;

AUTHOR : J. BOOTHROYD:

From the information retained in a[1:n*(n+1)div2] and b[1:n] following the use of LL09 and the eigenvectors of the tridiagonal matrix in z[1:n,1:n], computes the eigenvectors of the original matrix a;

```

procedure eigbise(c,b;n,ops,t,w,norm);  value n,eps,t;  integer n,t;
real eps,norm;  real array b,c,w;
begin integer i,k,a,j;  real g,h,p1,q1,bi,lim,lambda,y;  array p[1:n];
g:=abs(b[1]);  p1:=abs(c[1])+g;
for i:=2 step 1 until n do
begin h:=abs(b[i]);  q1:=g+h+abs(c[i]);
if q1>p1 then p1:=q1;  g:=h
end;
lambda:=norm:=lim:=p1;  bi:=b[1];
for i:=2 step 1 until n do
begin if bi=0.0 then bi:=eps*lim;
p[i]:=bi*bi;  bi:=b[i]
end;
for k:=1 step 1 until n do
begin g:=lambda;  h:=-lim;
for j:=1 step 1 until t do
begin lambda:=(g+h)/2.0;
p1:=0.0;  q1:=1.0;  a:=0;
for i:=1 step 1 until n do
begin y:=(c[i]-lambda)*q1-p[i]*p1;
p1:=q1;  q1:=y;
if p1>0.0 & q1 > 0.0 then a:=a+1
end i;
if q1=0.0 and p1>0 then a:=a-1;
if a>k then h:=lambda else g:=lambda;
if g=h then j:=t
end j;
w[k]:=(g+h)/2.0
end
end eigbise;

```

```

procedure trivectors(c,b,n,w,norm,eps,z);  value n,eps,norm;
integer n;  real norm,eps;  array c,b,w,z;
begin integer i,j,iplus1,iless1,nless1;  array m,p,q,r,int,x[1:n];
real bi,bi1,lambda,u,v,xi,miplus1,h,eta;
lambda:=norm;  eps:=eps*norm;  nless1:=n-1;
q[n]:=r[n]:=0.0;  eta:=1.0/n;
for j:=1 step 1 until n do
begin lambda:=lambda-eps;
if w[j]<lambda then lambda:=w[j];
u:=c[1]-lambda;  bi:=v:=b[1];
if bi=0.0 then bi:=v:=eps;
for i:=1 step 1 until nless1 do
begin iplus1:=i+1;
bi1:=b[iplus1];  if bi1=0.0 then bi1:=eps;
if abs(bi)>abs(u) then
begin miplus1:=m[iplus1]:=if u=0.0 and abs(bi)<eps then 1.0 else u/bi;
p[i]:=bi;  bi:=r[i]:=bi1;
bi1:=q[i]:=c[iplus1]-lambda;
u:=v-miplus1*bi1;  v:=-miplus1*bi;
int[iplus1]:=1.0
end
else
begin miplus1:=m[iplus1]:=bi/u;
p[i]:=u;  q[i]:=v;  r[i]:=0.0;
u:=c[iplus1]-lambda-miplus1*v;
bi:=v:=bi1;  int[iplus1]:=-1.0
end
end i;
p[n]:=if u#0.0 then u else eps;
v:=u:=h:=0.0;
for i:=n step -1 until 1 do
begin xi:=x[i]:=(eta-q[i]*u-r[i]*v)/p[i];
v:=u;  u:=xi;  h:=h+abs(xi)
end;
* for i:=1 step 1 until n do x[i]:=x[i]/h;
u:=x[1];  iless1:=1;
for i:=2 step 1 until n do
begin if int[i]>0.0 then begin v:=x[iless1]:=x[i];
u:=x[i]:=u-m[i]*v
end
else u:=x[i]:=x[i]-m[i]*u;
iless1:=i
end i;
v:=u:=h:=0.0;
for i:=n step -1 until 1 do
begin xi:=x[i]:=(x[i]-q[i]*u-r[i]*v)/p[i];
v:=u;  u:=xi;  h:=h+xi*xi;
end;
h:=1.0/sqrt(h);
for i:=1 step 1 until n do z[i,j]:=x[i]*h
end j
end trivectors;

```

LL12

```
procedure eigvectors(a,b,z,n); value n; integer n; array a,b,z;
begin integer i,j,k,iplus1,ck,cipplus1; real s,bi;
  for j:=1 step 1 until n do
    begin for i:=n-2 step -1 until 1 do
      begin bi:=b[i];
        if bi#0.0 then
          begin s:=0.0; iplus1:=i+1; ck:=cipplus1:=col[iplus1];
            for k:=iplus1 step 1 until n do begin s:=s+z[k,j]*a[i+ck]; ck:=ck+k end;
            s:=s/(bi*a[i+cipplus1]); ck:=cipplus1;
            for k:=iplus1 step 1 until n do begin z[k,j]:=z[k,j]+a[i+ck]*s; ck:=ck+k end
          end
        end
      end
    end
  end eigvectors;
```

GENERATE UNSYMMETRIC TEST MATRIX

procedure testmx(a,n); value n; integer n; array a;

comment HUCC LIBRARY PROCEDURE LZ 01:

AUTHOR J. BOOTHROYD:

The procedure generates matrices of the type described by T.J.Dekker,
Report No. MR 63, Mathematical Centre, Amsterdam.

The matrices have the following properties :-

- (a) elements $a[i,j]$ are integers
- (b) elements of the inverse, $(-1)^{i+j} * a[i,j]$, are also integers
- (c) the degree of ill condition increases rapidly with n
- (d) the determinant of all matrices is 1.

Computation of a matrix order 15 is possible on the 503 with a
39 bit register. As real matrices however the order of the largest
useable matrix should be restricted to that value of n(12) for which
all elements of the matrix have an exact floating point representation;

```

procedure testmx(a,n);  value n;  integer n;  array a;
begin integer i,j,k,fi,gi,d,q,r; boolean even;  integer array f,g[1:n];
comment first we compute  $F = \text{diag}(f_i)$ ;
fi:=f[1]:=n;  j:=n*n;
for i:=1 step 1 until n-1 do
begin d:=i*i;  k:=j-d;
q:=fi div d;  r:=fi-q*d;
f[i+1]:=fi:=q*k+(r*k) div d
end;
comment and now, using a modified prime factors algorithm to obtain  $G = \text{diag}(g_i)$  we compute
 $FG^{-1}$ , whose elements replace those of  $F$ ;
for i:=1 step 1 until n do
begin d:=gi:=1;  q:=fi:=f[i];  j:=2;
newj:  even:= false;
next:  if q > j then
begin q:=fi div j;
if fi≠q*j then begin j:=j+d;  d:=2;  goto newj end;
if even then gi:=gi*j;  even:= not even;
fi:=q;  goto next
end;
g[i]:=gi;  f[i]:=f[i] div gi
end;
comment finally, in one operation  $(FG^{-1})HG$  where  $H$  is a non-existent Hilbert
matrix whose reciprocal elements  $,i+j-1$ , are computed as we go;
for i:=1 step 1 until n do
begin fi:=f[i];
for j:=1 step 1 until n do
begin gi:=g[j];  k:=i+j-1;
q:=fi div k;  r:=fi-q*k;
a[i,j]:=q*gi+(r*gi) div k
end
end
end testmx;

```

LEAST SQUARES POLYNOMIAL FIT

procedure LSQFIT(x,y,N,a,n); value N,n; integer N,n;

real array x,y,a;

comment HUCC LIBRARY PROCEDURE MC01:

AUTHOR J. BOOTHROYD :

Given the $N+1$ sample pairs $(x_0, y_0), (x_1, y_1) \dots (x_N, y_N)$ in arrays $x, y[0:N]$ the procedure determines the coefficients $a_1, a_2, a_3, \dots, a_{n+1}$ of the n th order polynomial approximation

$$y = a_1 + a_2 x + a_3 x^2 + \dots + a_{n+1} x^n$$

obtained from applying the least squares principle to the given data. The coefficients a_1, a_2, \dots, a_{n+1} occupy the corresponding element positions of $a[1:n+1]$. The procedure uses SOLVE (LE02) to obtain the solutions of the normal equations. The matrix of coefficients of the normal equations becomes increasingly ill-conditioned as n increases, and it is recommended that the use of this procedure be restricted to values of $n < 6$;

```
procedure LSQFIT(x,y,N,a,n);  value N,n;  integer N,n;  real array x,y,a;
begin  integer i,j,k,nplus1,nplus2;  real p,q,xk;  real array b[1:n+1,1:n+2];
nplus1:=n+1;  nplus2:=n+2;
for i:=1 step 1 until nplus1 do
begin b[i,nplus2]:=0.0;
  for j:=1 step 1 until i do b[i,j]:=0.0
end;
for k:=0 step 1 until N do
begin p:=1.0;  xk:=x[k];
  for i:=1 step 1 until nplus1 do
begin q:=1.0;
    for j:=1 step 1 until i do
begin b[i,j]:=b[i,j]+p*q;  q:=q*xk end;
    b[i,nplus2]:=b[i,nplus2]+p*y[k];
    p:=p*xk
  end
end;
  for i:=2 step 1 until nplus1 do
    for j:=i-1 step -1 until 1 do b[j,i]:=b[i,j];
  SOLVEQ(b,nplus1);
  for i:=1 step 1 until nplus1 do a[i]:=b[i,nplus2]
end LSQFIT;
```

Least Squares Linear Fit

```
procedure Linfit(N,x,y,a0,a1,xm,ym); value N; integer N;  
real a0,a1,xm,ym; real array x,y;  
comment HUCC LIBRARY PROCEDURE MC02
```

AUTHOR : J.N. BAXTER

Given the $N+1$ sample pairs $(x_0, y_0), (x_1, y_1), \dots, (x_N, y_N)$
in arrays $x, y[0:N]$ the procedure determines the coefficients
a0 and a1 of the straight line approximation

$$y = a_0 + a_1 x$$

obtained from applying the least squares principle to
the given data. This is simpler than using procedure
LSQFIT(MC01) with $n=1$ for this purpose.

The procedure also evaluates the arithmetic means,
sm and ym, of the contents of the arrays x and y
respectively, for the reason that the straight line of
best fit may be expressed conveniently as

$$(y - y_m) = a_1(x - x_m)$$

Should this not be desired, and the values of xm and ym
be of no interest, the procedure call can be such as

Linfit(N, independent variable, dependent variable,
intercept, slope, slope, slope)

so that the values of xm and ym in turn are assigned to the
variable slope and then overwritten by the value of a1;

```
procedure Linfit(N,x,y,a0,a1,xm,ym);  value N;  integer N;  real a0,a1,xm,ym;  real array x,y;
  begin real n,xi,yi,sx,sy,sxx,sxy;  integer i;
  n:=N+1.0;  sx:=sy:=sxx:=sxy:=0.0;
  for i:=0 step 1 until N do
    begin xi:=x[i];  yi:=y[i];
    . . .  sx:=sx+xi;  sy:=sy+yi;  sxx:=sxx+xi*xi;  sxy:=sxy+xi*yi
    end;
  xm:=sx/n;  ym:=sy/n;
  a1:=(n*sxy-sx*sy)/(n*sxx-sx*sx);  a0:=(sy-a1*sx)/n
  end Linfit;
```

Weighted Linear Least Squares Fit

procedure wt linfit(N,xk,yk,k,wk,a0,a1,xm,ym); value N; integer N,k;
real xk,yk,wk,a0,a1,xm,ym;

comment HUCC LIBRARY PROCEDURE MC03

AUTHOR : J.N. BAXTER

In cases where a curve is to be fitted to sample pairs (x_i, y_i) by means of a transformation into $(f(x), f'(y))$ and a linear fit on the functions, it is necessary to weight the data so that the least squares approximation obtain is that for the original x and y. This is a Jensen procedure for the purpose, and k is the Jensen parameter which is utilised in the procedure call. For example, if $f(x) = 1/x$ and $f'(y) = \ln y$ with a chosen weighting for each point of y^*y the procedure call might be

wtlinfit(N,1.0/x[k],ln(y[k]),k,y[k]+2)

to find: (cept,slope,meanfx,meanfy).

The procedure is also capable of an unweighted fit, if the formal parameter wk is replaced in the procedure call by 1.0, but an unweighted fit is more simply carried out by procedure Linfit(MC02);

```
procedure wtlinfit(N,xk,yk,k,wk,a0,a1,xm,ym);  value N;  integer N,k;  real xk,yk,wk,a0,a1,xm,ym;
begin real w,x,yw,xw,sw,sxw,syw,sxxw,sxyw;
sxw:=syw:=sxxw:=sxyw:=sw:=0.0;
for k:=0 step 1 until N do
begin x:=xk;  w:=wk;  yw:=yk*w;  xw:=x*w;
sxw:=sxw+xw;  syw:=syw+yw;  sw:=sw+w;  sxxw:=sxxw+x*xw;  sxyw:=sxyw+yw*x
end;
xm:=sxw/sw;  ym:=syw/sw;
a1:=(sw*sxyw-sxw*syw)/(sw*sxxw-sxw*sxw);  a0:=(syw-a1*sxw)/sw
end wtlinfit;
```

Standard Deviation of Fitted Straight Line

```
real procedure Lindev(N,x,y,a0,a1); value N,a0,a1; integer N;  
real a0,a1; real array x,y;  
comment HUCC LIBRARY PROCEDURE MC04
```

AUTHOR : J.N. BAXTER

A procedure to evaluate the standard deviation of the N+1 data points held in x,y[0:N] from the fitted straight line $y=a0+a1x$. (The values of a0 and a1 supplied to this procedure may be those calculated by Linfit(MC02) or may be other values, chosen at will, or derived from other work).

To avoid the infinite value which would otherwise result from the use of this procedure with 2(N=1) data points the assignment Lindev:=5*10⁷⁵ occurs. The procedure is not protected from use with N=0 which case fails on a sqrt error;

MC04

```
real procedure Lindev(N,x,y,a0,a1); value N,a0,a1; integer N; real a ,a1; real array x,y;
begin integer i,n2; real sum,dev;
sum:=0.0;
for i:= 0 step 1 until N do begin dev:=a0+a1*x[i]-y[i]; sum:=sum+dev*dev end;
n2:=N-1;
Lindev:=(if n2#0 then sqrt(sum/n2) else 5.01075)
end Lindev;
```

LOCATE MINIMUM OF FUNCTION f(x)

```
real procedure MINX(a,b,eps,x,fx,fval); value a,b,eps;  
real a,b,eps,x,fx,fval;
```

comment HUCC LIBRARY PROCEDURE MD01:

AUTHOR J. BOOTHROYD :

A procedure to locate the minimum of fx in the interval $a \leq x \leq b$ by the method of trisection. fx must be monotonic decreasing from $x=a$ to the position of the minimum and thereafter monotonic increasing until $x=b$. eps is an absolute tolerance and the procedure aims to locate the minimum with an error less than eps. Whether it succeeds will depend on how well the minimum is defined taking into account the specified eps and the precision to which fx may be evaluated for any given real number representation. A relative tolerance delta may be specified using the call MINX(a,b,(b-a)*delta,x,fx,fval). On exit from the procedure fval yields the minimum value of the function;

```
real procedure MINX(a,b,eps,x,fx,fval);  value a,b,eps;  real a,b,eps,x,fx,fval;
begin real sep,fx1,fx2;  integer d;  switch s:=L,movea;
L:  sep:=(b-a)/3.0;  x:=a+sep;  fx1:=fx;  x:=b-sep;  fx2:=fx;
  if fx2=fx1 then begin d:=0; b:=b-sep; goto movea end;
  d:=1;
  if fx2>fx1 then b:=b-sep else movea: a:=a+sep;
  if sep>eps then goto L;
MINX:=x:= if d#0 then a+sep else (a+b)/2.0;  fval:= fx
end MINX;
```

MG 01

CONVERT SEXAGESIMAL ANGLES TO RADIANS

```
real procedure sexrad(deg,min,sec); value deg,min,sec;  
integer deg,min; real sec;
```

comment HUCC LIBRARY PROCEDURE MG 01:

AUTHOR J. BOOTHROYD:

Converts angular measure in the sexagesimal (360,60,60) system
to radians;

MG 02

CONVERT RADIAN ANGULAR MEASURE TO SEXAGESIMAL UNITS

```
procedure radtosex(rad,deg,min,sec); value rad;  
integer deg,min; real rad,sec;
```

comment HUCC LIBRARY PROCEDURE MG 02:

AUTHOR J. BOOTHROYD:

Converts angular measure in radians to degrees, minutes and
seconds in the sexagesimal system. The result is correctly
rounded to two decimal places of seconds;

MG 03

CONVERT CENTESIMAL ANGLES TO RADIANS

```
real procedure centrad (deg,min,sec); value deg,min,sec;  
integer deg,min; real sec;
```

comment HUCC LIBRARY PROCEDURE MG 03:

AUTHOR J. BOOTHROYD:

Converts angular measure from the centesimal (400,100,100) system
to radians;

MG 04

CONVERT RADIAN ANGULAR MEASURE TO CENTESIMAL UNITS

```
procedure radtocent(rad,deg,min,sec); value rad;
integer deg, min; real rad,sec;
```

comment HUCC LIBRARY PROCEDURE MG 04:

AUTHOR J. BOOTHROYD:

Converts angular measure from radians to centesimal (400,100,100) units. The result is correctly rounded to two decimal places of seconds;

MG 05

COMPUTE DISTANCE AND GRID BEARING

```
procedure bearing(x,y,r,t); value x,y; real x,y,r,t;
```

comment HUCC LIBRARY PROCEDURE MG 05:

AUTHOR J. BOOTHROYD:

Computes r, the length, and t, the clockwise radian angular bearing from north of a line joining the origin (0,0) to the point (x,y). For the bearing and length of a line joining (x1,y1) to (x2,y2) use the call bearing (x2-x1,y2-y1,r,t);

MG 06

COMPUTE POLAR COORDINATES OF POINT X,Y

```
procedure polar (x,y,r,t); value x,y; real x,y,r,t;
```

comment HUCC LIBRARY PROCEDURE MG 06:

AUTHOR NATIONAL PHYSICAL LABORATORY:

Computes the polar coordinates r,t corresponding to given x,y Cartesian coordinates;

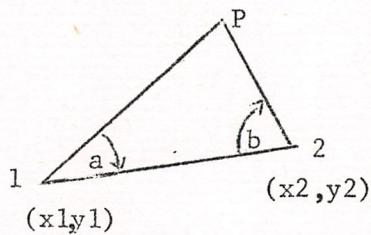
COMPUTE POINT OF INTERSECTION FROM ANGLES

```
procedure intang(x1,y1,a,x2,y2,b,x,y); value x1,y1,a,x2,y2,b;
real x1,y1,a,x2,y2,b,x,y;
```

comment HUCC LIBRARY PROCEDURE MG 07:

AUTHOR J. BOOTHROYD:

Computes the coordinates of the point of intersection of the lines through $(x_1, y_1), (x_2, y_2)$ which make angles, with respect to the line 12, of a and b radians respectively.



The following convention of signs must be observed.

Positive angles a are those swept out by closing the line 1P onto 12 clockwise.

Positive angles b are those swept out by closing the line 21 onto 2P clockwise.

The formulae used are:

$$x = \frac{((x_2 - x_1) \cot(a) - (y_2 - y_1))}{\cot(a) + \cot(b)} + x_1$$

$$y = \frac{((y_2 - y_1) \cot(a) + (x_2 - x_1))}{\cot(a) + \cot(b)} + y_1$$

```
real procedure sexrad(deg,min,sec); value deg,min,sec;
integer deg,min; real sec;
begin sec:= ((sec/60.0+min)/60.0+abs(deg))*0.0174532925;
  sexrad:=if deg<0 then -sec else sec
end sexrad;
```

MG 01

```
procedure rautosex(rad,deg,min,sec); value rad;
integer deg,min; real rad,sec;
begin real x;
  x:= abs(rad)/0.0174532925+0.00000139;
  deg:= entier(x); x:=(x-deg)*60.0;
  if rad<0.0 then deg:=-deg;
  min:= entier(x); sec:=(x-min)*60.0*100.0;
  sec:=entier(sec)/100.0
end radtosex;
```

MG 02

```
real procedure centrad(deg,min,sec); value deg,min,sec;
integer deg,min; real sec;
begin sec:=((sec/100.0+min)/100.0+abs(deg))*0.01570796327;
  centrad:= if deg<0 then -sec else sec
end centrad;
```

MG 03

```
procedure rautocent(rad,deg,min,sec); value rad;
integer deg,min; real rad,sec;
begin real x;
  x:= abs(rad)/0.01570796327+0.0000005;
  deg:= entier(x); x:=(x-deg)*100.0;
  if rad<0.0 then deg:=-deg;
  min:= entier(x); sec:=(x-min)*100.0*100.0;
  sec:=entier(sec)/100.0
end radtocent;
```

MG 04

```
procedure bearing(x,y,r,t); value x,y; real x,y,r,t;
begin real pi,piby2; pi:=3.141592654; piby2:=1.570796327;
r:=sqrt(x*x+y*y);
if r=0.0 then t:=0.0
else begin t:= if abs(x)<abs(y) then arctan(x/y)
else piby2-arctan(y/x);
y:=x+y;
if y<0.0 or x<0.0 then t:=t+pi;
if y>0.0 and x<0.0 then t:=t+pi
end
end bearing;
```

MG 05

```
procedure polar(x,y,r,t); value x,y; real x,y,r,t;
begin real pi,piby2; pi:=3.141592654; piby2:=1.570796327;
r:=sqrt(x*x+y*y);
if r=0.0 then t:=0.0
else begin t:= if abs(x)<abs(y) then piby2-arctan(x/y)
else arctan(y/x);
x:=x+y;
if x < 0.0 or y<0 then t:=t+pi;
if x > 0.0 and y<0 then t:=t+pi
end
end polar;
```

MG 06

```
procedure intang(x1,y1,a,x2,y2,b,x,y); value x1,y1,a,x2,y2,b;
real x1,y1,x2,y2,a,b,x,y;
begin a:=1.0/tan(a); b:=a+1.0/tan(b);
x2:=x2-x1; y2:=y2-y1;
x:=(x2*a-y2)/b+x1;
y:=(y2*a+x2)/b+y1
end intang;
```

MG 07

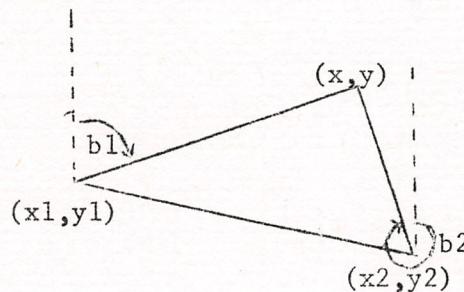
COMPUTE POINT OF INTERSECTION FROM BEARINGS

```
procedure intbrg (x1,y1,b1,x2,y2,b2,x,y); value x1,y1,b1,x2,y2,b2;  
real x1,y1,b1,x2,y2,b2,x,y;
```

comment HUCC LIBRARY PROCEDURE MG 08:

AUTHOR J. BOOTHROYD:

Computes the point of intersection of the line, bearing b_1 , through (x_1, y_1) with the line, bearing b_2 , through (x_2, y_2) .



This procedure uses MG 05, procedure bearing, to compute the bearing of the line joining (x_1, y_1) to (x_2, y_2) . The base angles of the triangle are then computed and the point of intersection evaluated by the method of MG 07;

THREE POINT RESECTION BY ANGLES

```

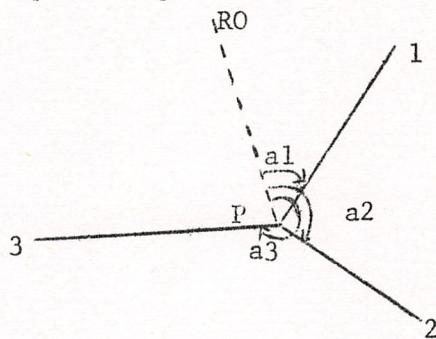
procedure resect(x1,y1,a1,x2,y2,a2,x3,y3,a3,x,y,scale,fail);
value x1,y1,a1,x2,y2,a2,x3,y3,a3,scale;
real x1,y1,a1,x2,y2,a2,x3,y3,a3,x,y,scale; label fail;

```

comment HUCC LIBRARY PROCEDURE MG 09:

AUTHOR J. BOOTHROYD:

Computes the coordinates x, y of the point of intersection P of the lines P_1, P_2, P_3 from the angles a_1, a_2, a_3 (in radians) where a_1, a_2, a_3 are measured clockwise from a reference observation line to the lines P_1, P_2 and P_3 respectively.



The RO line may coincide with the line P_1 , i.e. $a_1=0$.

The points may be in any order with respect either to themselves or to the resected point P . The procedure rejects all geometrically impossible configurations by an exit to the label fail, and special steps are taken to preserve accuracy in the numerically difficult configurations.

The procedure is based on the formulae:-

$$\begin{aligned}
 x-x_1 &= \tan b_1(y-y_1) \\
 x-x_2 &= \tan b_2(y-y_2) \\
 x-x_3 &= \tan b_3(y-y_3) \\
 b_2 &= b_1 + (a_2 - a_1), \quad b_3 = b_1 + (a_3 - a_1)
 \end{aligned}$$

where b_1, b_2, b_3 are the grid bearings of the lines P_1, P_2, P_3 .

From these may be derived the formula

$$\begin{aligned}
 \tan(b_1) &= \frac{((x_2-x_1)\cot(a_2-a_1)-(x_3-x_1)\cot(a_3-a_1)+(y_3-y_1))}{((y_2-y_1)\cot(a_2-a_1)-(y_3-y_1)\cot(a_3-a_1)-(x_3-x_1))}
 \end{aligned}$$

As is well known to surveyors, the case in which $P, 1, 2, 3$ together form a cyclic quadrilateral is indeterminate. Mathematically this case yields $\tan(b_1)=0/0$. Near cyclic cases also produce large inaccuracies and are detected by evaluating the numerator and denominator of the expression for $\tan(b_1)$ separately. The parameter scale has been provided for this case. The value given to scale should be twice the distance from the approximate centre of the known points to the most distant known point. The procedure computes 5% of this value and if both numerator and denominator are less than $0.05 * \text{scale}$ a cyclic case is assumed;

```

procedure resect(x1,y1,a1,x2,y2,a2,x3,y3,a3,x,y,scale,fail);
value x1,y1,a1,x2,y2,a2,x3,y3,a3,scale;
real x1,y1,a1,x2,y2,a2,x3,y3,a3,x,y,scale; label fail;
begin real t,r,eps; boolean a2zero,a3zero,absa2,absa3; switch s:=L2,L3,done;
procedure solve(x1,y1,x2,y2,a1,a2);
value x1,y1,x2,y2,a1,a2; real x1,y1,x2,y2,a1,a2;
begin y2:=(y2*a2-x2)/(a2-a1);
x:=y2*a1+x1; y:=y2+y1
end solve;
eps:=0.05*scale;
x3:=x3-x2; y3:=y3-y2; a3:=a3-a1;
x2:=x2-x1; y2:=y2-y1; a2:=a2-a1;
a3zero:=a3=0.0; a2zero:=a2=0.0;
if a2zero and a3zero then goto fail;
if a3zero then
begin a2:=1.0/tan(a2); if abs(a2)>10^5 then goto fail;
t:=y3+y2;
if t=0.0 then begin y:=0.0; x:=a2*y2+x2+x1; goto done end
else begin a1:=(x3+x2)/(y3+y2); goto L2 end
end;
if a2zero then
begin a3:=1.0/tan(a3); if abs(a3)>10^5 then goto fail;
if y2=0.0 then begin y:=0.0; x:=a3*y3+x3+x2+x1; goto done end
else begin a1:=x2/y2; goto L3 end
end;
a2:=1.0/tan(a2); a3:=1.0/tan(a3); absa2:=abs(a2)>10^5; absa3:=abs(a3)>10^5;
if absa2 and absa3 then goto fail;
if absa3 then
begin t:=y3+y2;
if t=0.0 then begin y:=0.0; x:=a2*y2+x2+x1; goto done end
else begin a1:=(x3+x2)/(y3+y2); goto L2 end
end;
if absa2 then
begin if y2=0.0 then begin y:=0.0; x:=a3*y3+x3+x2+x1; goto done end
else begin a1:=x2/y2; goto L3 end
end;
t:=(x2*a2-(x3+x2)*a3+y3); r:=(y2*a2-(y3+y2)*a3-x3);
if abs(t)<eps and abs(r)<eps then goto fail;
a1:=t/r;
L2: a2:=(1.0+a2*a1)/(a2-a1); solve(x1,y1,x2,y2,a1,a2); goto done;
L3: a3:=(1.0+a3*a1)/(a3-a1); solve(x1,y1,x3+x2,y3+y2,a1,a3);
done: end resect;

```

MULTIPLE POINT RESECTION FROM ANGLES

```
procedure avresect(obs,n,scale,x,y);  value n,scale;
real scale,x,y;  integer n;  array obs;
```

comment HUCC LIBRARY PROCEDURE MG 10:

AUTHOR J. BOOTHROYD:

The minimum number of known points necessary for the evaluation of the coordinates of a resected point is 3. Where the number of known points exceeds 3 some method of obtaining the best estimate of x,y from all computed values of x,y is required. This procedure computes the weighted mean of all calculations based on the selection of all possible combinations of 3 points selected from n points, using the procedure resect (MG 09) $\binom{n}{3}$ times. The array obs[1:n,1:4] contains the observational data for the n observations. For the rth observation x_r , y_r , a_r and weight w_r should be provided respectively in obs[r,1], obs[r,2], obs[r,3] and obs[r,4]. For a resection based on points i,j,k with respective weights wti,wtj,wtk a group weight wti+wtj+wtk is assumed. The parameter scale is used to reject near cyclic cases for details of which see the description of MG 09. For angular errors of the order of 1 second the procedure accuracy over eight (difficult) points is of the order of .0001 unit with a scale of 100 units. This means an accuracy of .1 yd for resections over distances of 100,000 yds. The errors are approximately linear with angular error and will increase if the number of observations decreases;

```

procedure intbrg(x1,y1,b1,x2,y2,b2,x,y);  value x1,y1,b1,x2,y2,b2;
real x1,y1,b1,x2,y2,b2,x,y;
begin      real r,t;
           x2:=x2-x1;  y2:=y2-y1;
           bearing(x2,y2,r,t);
           b1:=1.0/tan(t-b1);
           t:=b1+1.0/tan(3.141592654+b2-t);
           x:=(x2*b1-y2)/t+x1;
           y:=(y2*b1+x2)/t+y1
end intbrg;

```

MG 08

```

procedure avresect(obs,n,scale,x,y);  value n,scale;
real scale,x,y;  integer n;  array obs;
begin integer i,j,k;
real sn,sx,sy,x1,y1,a1,x2,y2,a2,xx,yy,wti,wtj,wt;
switch s:=loop;
sn:=sx:=sy:=0.0;
for i:=1 step 1 until n do
begin x1:=obs[i,1];  y1:=obs[i,2];  a1:=obs[i,3];  wti:=obs[i,4];
for j:=i+1 step 1 until n do
begin x2:=obs[j,1];  y2:=obs[j,2];  a2:=obs[j,3];  wtj:=obs[j,4];
for k:=j+1 step 1 until n do
begin resect(x1,y1,a1,x2,y2,a2,obs[k,1],obs[k,2],obs[k,3],xx,yy,scale,loop);
wt:=wti+wtj+obs[k,4];
sn:=sn+wt;
sx:=sx+wt*xx;
sy:=sy+wt*yy;
loop: end k
end j
end i;
if sn<0 then begin x:=sx/sn;  y:=sy/sn  end
else print punch(3),#1?RESECT ERROR?,wait
end avresect;

```

MG10

THREE POINT RESECTION FROM DISTANCES

```
procedure distresect(x1,y1,d1,x2,y2,d2,x3,y3,d3,x,y,fail);
value x1,y1,d1,x2,y2,d2,x3,y3,d3;
real x1,y1,d1,x2,y2,d2,x3,y3,d3,x,y; label fail;
```

comment HUCC LIBRARY PROCEDURE MG 11:
AUTHOR J. BOOTHROYD:

Computes the coordinates of a resected point P from the distances d1,d2,d3 from P to known points (x1,y1),(x2,y2),(x3,y3). This procedure has been written to provide assistance to those making use of telurometry.

The equations are $(x-x1)^2 + (y-y1)^2 = d1^2$
 $(x-x2)^2 + (y-y2)^2 = d2^2$
 $(x-x3)^2 + (y-y3)^2 = d3^2$

from which two linear equations in x and y may be obtained by subtraction in pairs. If the data provided results in a zero determinant of the two linear equations the procedure exits to the label fail;

MULTIPLE POINT RESECTION USING DISTANCES

```
procedure avdistresect(obs,n,x,y);  value n;
integer n;  real x,y;  array obs;
```

comment HUCC LIBRARY PROCEDURE MG 12:

AUTHOR J. BOOTHROYD:

This procedure evaluates the coordinates x,y of a resected point as the weighted mean of all computations based on the use of all combinations of 3 points selected from n observations, using distances in all cases. The procedure disresect (MG 11) is called n^3 times. For the r^{th} observation the data $x_r, y_r, d_r, \text{weight}_r$ should occupy $\text{obs}[r,1], \text{obs}[r,2], \text{obs}[r,3], \text{obs}[r,4]$ respectively of the array $\text{obs}[1:n,1:4]$.

For a computation using observations i,j,k with respective weights $\text{wti}, \text{wtj}, \text{wtk}$, a group weight $\text{wti}+\text{wtj}+\text{wtk}$ is assumed;

```

procedure distresect(x1,y1,d1,x2,y2,d2,x3,y3,d3,x,y,fail);
value x1,y1,d1,x2,y2,d2,x3,y3,d3;
real x1,y1,d1,x2,y2,d2,x3,y3,d3,x,y; label fail;
begin real a11,a12,a21,a22,b1,b2,det;
a11:=(x2-x1); a12:=(y2-y1);
a21:=(x3-x1); a22:=(y3-y1); d1:=d1*d1;
b1:=(a11*(x2+x1)+a12*(y2+y1)-d2*d2+d1)/2.0;
b2:=(a21*(x3+x1)+a22*(y3+y1)-d3*d3+d1)/2.0;

det:=a11*a22-a21*a12;
if det=0.0 then goto fail;
x:=(a22*b1-a21*b2)/det;
y:=(a11*b2-a21*b1)/det
end distresect;

```

MG 11

```

procedure aviistresect(obs,n,x,y); value n;
real x,y; integer n; array obs;
begin integer i,j,k;
real sn,sx,sy,x1,y1,d1,x2,y2,d2,xx,yy,wti,wtj,wt;
switch s:=loop;
sn:=sx:=sy:=0.0;
for i:=1 step 1 until n do
begin x1:=obs[i,1]; y1:=obs[i,2]; d1:=obs[i,3]; wti:=obs[i,4];
for j:=i+1 step 1 until n do
begin x2:=obs[j,1]; y2:=obs[j,2]; d2:=obs[j,3]; wtj:=obs[j,4];
for k:=j+1 step 1 until n do
begin distresect(x1,y1,d1,x2,y2,d2,obs[k,1],obs[k,2],obs[k,3],xx,yy,loop);
wt:=wti+wtj+obs[k,4];
sn:=sn+wt;
sx:=sx+wt*xx;
sy:=sy+wt*yy;
loop: end k
end j
end i;
if sn=0 then begin x:=sx/sn; y:=sy/sn end
else print punch(3),#1?DISTRESECT ERROR?,wait
end avdistresect;

```

MG 12

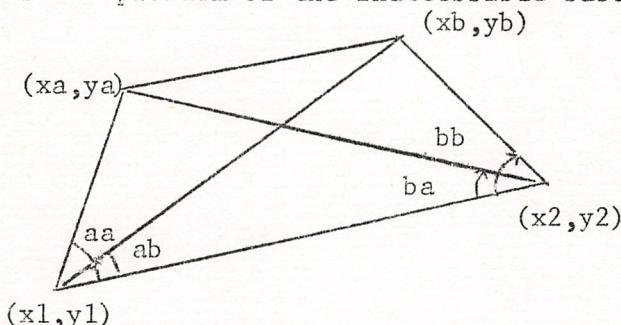
THE INACCESSIBLE BASE PROBLEM

```
procedure farbase(xa,ya,xb,yb,aa,ab,ba,bb,x1,y1,x2,y2);
value xa,ya,xb,yb,aa,ab,ba,bb;
real xa,ya,xb,yb,aa,ab,ba,bb,x1,y1,x2,y2;
```

comment HUCC LIBRARY PROCEDURE MG 13:

AUTHOR J. BOOTHROYD:

Solves the problem of the inaccessible base.



(x_1, y_1) and (x_2, y_2) are stations whose coordinates are required.
 (x_a, y_a) , (x_b, y_b) are stations whose coordinates are known.
 aa, ab, ba, bb are angular observations from (x_1, y_1) and (x_2, y_2) on (x_a, y_a) and (x_b, y_b) using the angular sign convention of procedure MG 07. Standard survey computational methods would compute the tangent of the angle of inclination of line a, b to line $1, 2$ and, from the computed bearing of $1, 2$ and the measured angles deduce the bearings of all other lines. This would provide sufficient information to evaluate x_1, y_1, x_2, y_2 using procedure MG 08. This is not the best method if automatic computing facilities are available. Using the formulae described in procedure MG 07 we obtain the four equations :-

$$\begin{aligned} x_a &= ((x_2 - x_1) \cot aa - (y_2 - y_1)) / (\cot aa + \cot ba) + x_1 \\ x_b &= ((x_2 - x_1) \cot ab - (y_2 - y_1)) / (\cot ab + \cot ba) + x_1 \\ y_a &= ((y_2 - y_1) \cot aa + (x_2 - x_1)) / (\cot aa + \cot ba) + y_1 \\ y_b &= ((y_2 - y_1) \cot ab + (x_2 - x_1)) / (\cot ab + \cot ba) + y_1 \end{aligned}$$

which, rearranged yields four linear equations in the four unknowns x_1, y_1, x_2, y_2 . This procedure computes the coefficients of the linear system and calls on LE 02 (SOLVEQ) to obtain the solution;

```
procedure farbase(xa,ya,xb,yb,aa,ab,ba,bb,x1,y1,x2,y2);
value xa,ya,xb,yb,aa,ab,ba,bb;
real xa,ya,xb,yb,aa,ab,ba,bb,x1,y1,x2,y2;
begin real aaba,abbb; array A[1:4,1:5];
aa:=1.0/tan(aa); ab:=1.0/tan(ab);
ba:=1.0/tan(ba); bb:=1.0/tan(bb);
aaba:=aa+ba; abbb:=ab+bb;
A[1,1]:=A[3,2]:=ba; A[2,1]:=A[4,2]:=bb;
A[1,3]:=A[3,4]:=aa; A[2,3]:=A[4,4]:=ab;
A[1,4]:=A[2,2]:=A[3,3]:=A[4,3]:=1.0;
A[1,5]:=xa*aaba; A[2,5]:=xb*abbb;
A[3,5]:=ya*aaba; A[4,5]:=yb*abbb;
SOLVEQ(A,4);
x1:=A[1,5]; y1:=A[2,5]; x2:=A[3,5]; y2:=A[4,5]
end farbase;
```

3-DIMENSIONAL COORDINATE GEOMETRY AND VECTOR ARITHMETIC PACKAGE

The procedures of this package are collectively designated as MG14. No provision is made for copying individual procedures into user's tapes since these are brief enough to be re-written as required.

comment HUCC LIBRARY PROCEDURE MG14:

AUTHOR : B. ROBINSON:

The point $P(x,y,z)$ in three dimensional space has one-to-one correspondence with the vector \vec{V} which has components x,y,z parallel to the three mutually perpendicular coordinate axes respectively. In these procedures the vector \vec{V} (which may be visualised as the directed line-segment OP from the origin to the point P) is represented by an ALGOL array of three elements, which is declared thus

real array $V[1:3]$;

The elements of this array are the components of the vector,

$V[1]=x$

$V[2]=y$

$V[3]=z$

In procedure calls the vector is referred to as $V[j]$.

The subscript j must be declared as an integer and the declaration must be valid for all blocks in which these procedures are employed. Inside the procedures j is assigned the values 1,2, or 3 as needed for operations on the components of the vector. Outside the procedures j should not be used.

The procedures make extensive use of the ALGOL facility of calling by name. This is convenient because it makes possible the use of linear expressions of vectors as actual parameters in procedure calls, and the use of arrays of vectors or individual vectors. For example, a two-dimensional lattice of vectors with typical element L_{mn} would be declared as

real array $L[\ell_1:u_1, \ell_2:u_2, 1:3]$

and any individual vector of this lattice could be referred to in a procedure call as $L[m,n,j]$.

The result of this is that the text of the program will be conveniently reminiscent of vector algebra.

VECTOR ASSIGNMENT:

procedure make(yj) equal to:(xj); real xj,yj;

Examples : make(V2[j],V1[j]); effect is "V2:=V1"

make(V3[j],V1[j] + V2[j]); "V3:=V1+V2"

make(V2[j],V1[j]*sc); "V2:=V1*sc"

make(M[j],(A[j]+B[j])/2.0);

This finds the midpoint of the line-segment AB

effect is "M:=(A+B)/2"

make(G[j],(A[j]+B[j]+C[j])/3.0);

This finds the centroid of the triangle ABC,

effect is "G:=(A+B+C)/3"

MULTIPLE VECTOR ASSIGNMENT:

procedure make both(yj) and:(zj)equal to:(xj); real xj,yj,zj;

Example: makeboth(V[6,j],temp[j],V2[j]-V1[j];

effect is "V[6]:=temp:=V2-V1"

DOT PRODUCT OF TWO VECTORS:

real procedure dot(xj,yj); real xj,yj;

The procedure takes the value of the scalar product of the two vectors.

Example: r:=dot(V1[j],V2[j]);

If the dot product is zero, then the two vectors are perpendicular or else one is a zero vector.

LENGTH OF A VECTOR:

```
real procedure length(xj); real xj;
```

NOTE: uses procedure dot.

Examples : The length of the vector V from the origin is found by

```
r:=length(V[j]);
```

The distance between the ends of the vectors V_1 and V_2 is found by $r:=length(V2[j]-V1[j]);$

COSINE OF THE ANGLE BETWEEN TWO VECTORS:

```
real procedure cosang(xj,yj) escape label : (fail);
```

```
real xj,yj; label fail;
```

The procedure takes the value of the cosine of the angle subtended at the origin by the two vectors. If either vector is of zero length then the angle is indeterminate and the procedure exits to a label.

Example: $cst\theta:=cosang(V1[j],V2[j],fail);$

To find the cosine of an angle not at the origin the call by name is so arranged as to simulate a change of origin to the apex of the angle. Thus where the angle $\angle ABC$ is represented by the three vectors $a[j]$, $b[j]$ and $c[j]$,

```
cst $\theta$ :=cosang( $a[j]-b[j]$ ,  $c[j]-b[j]$ , fail);
```

NOTE: uses procedures dot and length.

VECTOR OF UNIT LENGTH PARALLEL TO A GIVEN VECTOR:

```
procedure unit(nj)parallel to:(xj)escape label: (fail);
```

```
real nj,xj; label fail;
```

Example: $unit(V[j],V[j],fail);$ converts the vector V into a unit vector in the same direction. The procedure will exit by the label if the vector given to it has zero length.

NOTE: uses procedures dot and length

CROSS PRODUCT OF TWO VECTORS:

procedure cross(xj,yj) is written into:(zj); real xj,yj,zj;

Example: cross(V1[j], V2[j], V3[j]); effect is "V3:=V1xV2"

The direction of V3 is perpendicular to the plane in which V1 and V2 lie.

If V1 and V2 are parallel vectors then V3 will be a zero vector. If V1 lies along the x-axis and V2 lies along the y-axis, then V3 will lie along the z-axis in accordance with the convention, this may be expressed by the statement that right-handed coordinates are used (remember that $V1xV2=-V2xV1$).

VECTOR OF UNIT LENGTH PERPENDICULAR TO TWO GIVEN VECTORS:

procedure perpunit(nj) perpendicular to:(xj,yj) escape label:(fail);

real nj,xj,yj; label fail;

Example: perpunit(n[j],V1[j],V2[j]);

NOTE: uses procedures dot, length, unit and cross.

BOX PRODUCT OF THREE VECTORS:

real procedure box(xj,yj,zj); real xj,yj,zj;

The procedure takes the value of the scalar triple product of the three vectors. If the three vectors are coplanar with the origin then the product is zero. It can therefore be used to test whether any four points are coplanar, by evaluation of

q:=box(a[j]-b[j], c[j]-b[j], d[j]-b[j]);

NOTE: uses procedures dot and cross.

comment vector procedures;

```

procedure make(yj,xj); real xj,yj; for j:=1,2,3 do yj:=xj;

procedure makeboth(yj,zj,xj); real xj,yj,zj; for j:=1,2,3 do yj:=zj:=xj;

real procedure dot(xj,yj); real xj,yj;
  begin real t; t:=0.0; for j:=1,2,3 do t:=t+xj*yj; dot:=t
  end dot product;

real procedure length(xj); real xj; length:=sqrt(dot(xj,xj));

real procedure cosang(xj,yj,fail); real xj,yj; label fail;
  begin real d; d:=length(xj)*length(yj); if d=0.0 then goto fail;
    cosang:= dot(xj,yj)/d
  end cos angle;

procedure unit(nj,xj,fail); real nj,xj; label fail;
  begin real d; d:=length(xj); if d=0.0 then goto fail;
    for j:=1,2,3 do nj:=xj/d
  end unit;

procedure cross(xj,yj,zj); real xj,yj,zj;
  begin real x1,x2,x3,y1,y2,y3;
    j:=1; x1:=xj; y1:=yj;
    j:=2; x2:=xj; y2:=yj;
    j:=3; x3:=xj; y3:=yj;
    zj:=x1*y2-x2*y1; j:=2; zj:=x3*y1-x1*y3;
    j:=1; zj:=x2*y3-x3*y2
  end cross product;

procedure perpunit(nj,xj,yj,fail); real nj,xj,yj; label fail;
  begin array a[1:3]; cross(xj,yj,a[j]); unit(nj,a[j],fail) end;

real procedure box(xj,yj,zj); real xj,yj,zj;
  begin array a[1:3]; cross(xj,yj,a[j]); box:= dot(a[j],zj)
  end box product;

```

COMPUTE FOURIER COEFFICIENTS

procedure fourier(f,a,b,n); value n; integer n; real array f,a,b;

comment HUCC LIBRARY PROCEDURE MH01:

AUTHOR: J. BOOTHROYD

Evaluates the coefficients of the Fourier expansion of a function $f(x)$ specified at the n sample points $f[0], f[1], \dots, f[n-1]$. n may be odd or even and on exit from the procedure the arrays $a, b[0:n \text{ div } 2]$ will contain the coefficients appropriate to either case as follows:—

n even (n = 2N)

$$f(x) = a[0] + \sum_{k=1}^{N-1} (a[k] * \cos(k\pi x/N) + b[k] * \sin(k\pi x/N)) + a[N] \cos \pi x$$

with $b[0] = b[N] = 0$.

n odd (n = 2N+1)

$$f(x) = a[0] + \sum_{k=1}^N (a[k] * \cos(2\pi kx/n) + b[k] * \sin(2\pi kx/n)),$$

with $b[0] = 0$.

The algorithm uses the recurrence relations described in R. W. Hamming - Numerical Methods for Scientists and Engineers, page 72.

```
procedure fourier(f,a,b,n); value n; integer n; real array f,a,b;
begin integer ndiv2,k,m; real c,s,vk,v1,u0,u1,um,t,ck,nby2;
  ndiv2:= n div 2; nby2:= n/2; t:= 6.2831853/n;
  c:= cos(t); s:= sin(t); vk:= 0.0; v1:= -1;
  for k:= 0 step 1 until ndiv2 do
    begin t:= c*vk; ck:= t-v1; v1:= vk; vk:= ck+t;
    t:= ck+ck; u1:=0.0; um:= f[n-1];
    for m:= n-2 step -1 until 1 do
      begin u0:= u1; u1:= um; um:= t*u1-u0+f[m] end;
    a[k]:= (ck*um-u1+f[0])/nby2;
    b[k]:= s*v1*um/nby2
    end;
  a[0]:= a[0]/2.0; if ndiv2*2=n then
    begin a[ndiv2]:= a[ndiv2]/2.0;
    b[ndiv2]:=0.0
    end
  end fourier;
```

SUM FOURIER SERIES

```
real procedure sumfourier(a,b,n,x); value n,x; real x;  
integer n; real array a,b;
```

comment HUCC LIBRARY PROCEDURE MH02:

AUTHOR: J. BOOTHROYD

Evaluates the function $f(x)$ for given argument x
by summing the Fourier series of the function, the
coefficients of which occupy arrays $a,b[0:n \text{ div } 2]$
as a result of using procedure MH01.

```
real procedure sumfourier(a,b,n,x);  value n,x;  real x;  integer n;  real array a,b;
comment evaluates the function f(x) for given argument x by summing the Fourier series whose
coefficients are the elements of a,b[0:ndiv2] derived by the use of procedure fourier, i.e.
n even (n=2N)
f(x)= a[0]+sigma(a[k]*cos(pi*k*x/N)+b[k]*sin(pi*k*x/N),k,1,N-1) + a[N]*cos(pi*x), b[0]=b[N]=0
n odd (n=2N+1)
f(x)=a[0]+sigma(a[k]*cos(2*pi*k*x/n)+b[k]*sin(2*pi*k*x/n),k,1,N),  b[0] = 0;
begin real vk,v1,sum,c,s,ck;  integer k,ndiv2;
  ndiv2:=n div 2;  x:=6.2831853*x/n;  c:=cos(x);  s:=sin(x);
  sum:=a[0];  vk:=1.0;  v1:=0.0;
  for k:=1 step 1 until ndiv2 do
    begin x:=c*vk;  ck:=x-v1;
      v1:=vk;  vk:=ck+x;
      sum:=sum+a[k]*ck+b[k]*s*v1
    end;
  sumfourier:=sum
end sumfourier;
```

LOCATE ROOT OF $f(x)=0$

real procedure bisec($f, x, a, b, \text{eps}, \text{error}$); value a, b, eps ;

real f, x, a, b, eps ; label error ;

comment HUCC LIBRARY PROCEDURE MRO1:

AUTHOR: NOT KNOWN

A procedure to find a root of $f(x)=0$ in the interval a to b by continued bisection. If $f(a)$ and $f(b)$ have the same sign the procedure exits to the label error . The error escape facility may be used to determine the interval (a, b) within which some root of interest lies. Suppose that $f(x)$ has several positive roots $x_1 x_2 x_3 \dots$ and it is known that the root separation $(x_2 - x_1)$, $(x_3 - x_2)$ etc. is at least z but the location of x_1 is in some doubt.

The routine

$xa := -z;$

strobe: $xa := xa + z;$

$y := \text{bisec}(f(x), x, xa, xa+z, 10^{-4}, \text{strobe})$

will increment xa until $f(xa)$ and $f(xa+z)$ have opposite signs when the bisec routine will evaluate the required root to an accuracy of 10^{-4} .

real procedure biseq(f,x,a,b,eps,error); value a,b,eps; real f,x,a,b,eps; label error;
comment the procedure finds a root of $f(x)=0$ in the interval a to b by
continued bisection. If $f(a)$ and $f(b)$ have the same sign the procedure exits
through the label error;

```
begin real q; x:=a; q:=f; x:=b;  
  if f*q>0 then goto error;  
  q:=(b-a)*sign(q)/2.0; x:=(a+b)/2.0; eps:=eps/2.0;  
  for q:=q/2.0 while abs(q)>eps do x:=x+(if f>0 then q else -q);  
  biseq:=x  
end biseq;
```

GENERAL SUM SERIES

real procedure sigma(tk,k,a,b); value a,b; integer k,a,b; real tk;

comment HUCC LIBRARY PROCEDURE MS01:

AUTHOR: J. BOOTHROYD

A Jensen procedure to evaluate the sum over k from a to b of tk with positive unit increments in k.

The parameters tk and k are called by name. As examples of the flexibility of this procedure the calls,

1. sigma(a+i*d,i,0,3)
2. sigma(a*r^i,i,0,n-1)
3. sigma(a[i]*b[i],i,1,N)

will evaluate respectively

1. The sum of 4 terms of an arithmetic progression with first term a and common difference d.
2. The sum of n terms of a geometric progression.
3. The inner product of the vectors a,b[1:N].

These examples illustrate the dependence of the actual first parameter on the variable actually used as the second parameter.

```
real procedure sigma(tk,k,a,b); value a,b; real tk; integer k,a,b;
comment the sum over k from a to b of tk, with positive unit increments in k;
begin real sum;
    sum:=0.0;
    for k:=a step 1 until b do sum:=sum+tk;
    sigma:=sum
end sigma;
```

20 ROMBERG INTEGRATION

```

real procedure romberg1(x,fx,a,b,order,eps); value a,b,eps,order;
      real x,fx,a,b,eps; integer order;
begin integer i,j,mark; real ti,k,h,n,m,t,fac,del,abstm,abstol,absti;
      array aa[1:order]; switch s:=test,newi,done;
      x:=a; ti:=fx; x:=b; ti:=aa[1]:=(ti+fx)/2.0;
      k:=h:=b-a;
      if k=0.0 then begin i:=0; goto done end;
      n:=1.0; abstol:=10^-8/abs(k);
      for i:=2 step 1 until order do
      begin h:=h/2.0; m:=0.0;
          for x:=a+h step k until b do :=m+fx;
          m:=m/n; t:=ti; j:=i;
          absti:=abs(ti);
      test: abstm:=abs(m-t);
          if ti≠0.0 then
              begin if abstm/absti<eps then goto done end
              else
                  begin if abstm<abstol then goto done end;
                  if i ≠ j then
                      begin if mark=0 then
                          begin t:=aa[j]:=m+(m-t)/fac;
                          mark:=1; del:=3.0*(fac+1.0);
                          fac:=fac+del
                          end
                          else
                          begin j:=j-1; if j=0 then goto newi;
                          m:=t; t:=aa[j]; mark:=0
                          end;
                      end
                      else
                      begin mark:=0; fac:=3.0; del:=0.0;
                      ti:=m:=aa[j]:=(m+t)/2.0; j:=j-1
                      end;
                      goto test;
      newi: k:=h; n:=n+n
      end i;
      i:=0;
done: romberg:=(b-a)*(if i>0 then (t+m)/2.0 else aa[1])
end romberg;

```

ADAPTIVE SIMPSON INTEGRATION

```
real procedure simps (f, x, a, b, eps); value a, b, eps; real
f, x, a, b, eps;
```

comment HUCC LIBRARY PROCEDURE MS02 :

AUTHOR J. BOOTHROYD : :

A modification of Algorithm 233 COMM. A.C.M. VOL 7 NO. 6
 JUNE 1964 p 348. The procedure evaluates the integral of
 f with respect to x from lower bound a to upper bound b .

$$\text{simps} := \int_a^b f \, dx$$

The parameter eps determines the acceptable relative error between successive evaluations of the integral I . If I_r and I_{r-1} are successive approximations, the process will terminate when

$$\text{abs} ((I_r - I_{r-1}) / I_r) < \text{eps}$$

unless $I_r = 0$ in which case, to avoid division by zero, the process terminates when the absolute difference between successive values of I differ by less than 10^{-8} . As an example of the use of simps , the call

$y := \text{simps} (1.0/x, x, 1.0, B, .0001)$
 will assign to y the value $\int_1^B (1/x) \, dx = \ln[B]$

with a tolerance of .0001.

The procedure operates correctly for the cases $a < b$, $a = b$ and $a > b$.

```
real procedure simps(f,x,a,b,eps);  value a,b,eps;  real f,x,a,b,eps;
begin real Z1,Z2,Z3,h,k;  switch s:=again,done;
  if a=b then begin h:=0.0;  goto done  end;
  x:=a;  Z1:=f;  x:=b;  Z1:=Z1+f;  k:=(b-a)/2.0;
  x:=a+k;  Z2:=f;  Z3:=Z1+4.0*Z2;  Z1:=Z1+2.0*Z2;
again:  Z2:=0;  h:=k/2;
  for x:=a+h step k until b do Z2:=Z2+f;
  Z1:=Z1+4.0*Z2;
  if Z1#0 then begin if abs((Z1-2.0*Z3)/Z1)<eps then goto done end
    else if abs(Z1-2.0*Z3)*h/3.0<10^-8 then goto done ;
  Z3:=Z1;  Z1:=Z1-2.0*Z2;  k:=h;  goto again;
done:   simps:=h*Z1/3.0
end simps;
```

EVALUATE DEFINITE INTEGRAL

```
real procedure havie (x,a,b,eps,integrand,m,mask);
value a,b,eps,m,mask; integer m; real a,b,eps,integrand,x,mask;
comment HUCC LIBRARY PROCEDURE MS 03
```

Author : ACM 257

Performs numerical integration of integrand with respect to x over the interval a to b

$$\text{havie} := \int_a^b \text{integrand} \, dx$$

eps is the convergence criterion, m is the maximum order of approximation to be considered in attempting to satisfy the eps criterion and mask is a value supplied by the user which will be returned by havie if convergence is not obtained.

The method is basically the trapezium formula with higher order corrections (m specifies an upper limit to the number of these performed).

Examples of calls are :-

$$(a) \text{ To evaluate } y = \int_0^{\pi/2} \cos x \, dx$$

$$y := \text{havie} (x, 0.0, 1.5707963, 10^{-6}, \cos(x), 10, 9.99999);$$

$$(b) \text{ To evaluate } z = \int_0^{4.3} e^{-y^2} dy$$

$$z := \text{havie} (y, 0.0, 4.3, 10^{-6}, \exp(-y^2), 10, 9.99999)$$

```
real procédure havie(x,a,b,eps,integrand,m,mask); value a,b,eps,m,mask; integer m; real a,b,x,eps,integrand,mask;
begin real h,endpts,sumt,sumu,d; integer i,j,k,n; switch s:=estimate,test,exit;
real array t,u,tprev,uprev[1:m];
x:=a; endpts:=integrand; x:=b; endpts:=0.5*(integrand+endpts);
sumt:=0.0; i:=n:=1; h:=b-a;
estimate: t[1]:=h*(endpts+sumt); sumu:=0.0; x:=a-h/2.0;
for j:=1 step 1 until n do begin x:=x+h; sumu:=sumu+integrand end;
u[1]:=h*sumu; k:=1;
test: if abs(t[k]-u[k]) < eps then begin havie:=0.5*(t[k]+u[k]); goto exit end;
if k ≠ i then
begin d:=2.0↑(k+k); t[k+1]:=(d*t[k]-tprev[k])/(d-1.0); tprev[k]:=t[k];
u[k+1]:=(d*u[k]-uprev[k])/(d-1.0); uprev[k]:=u[k]; k:=k+1;
if k=m then begin havie:=mask; goto exit end; goto test
end;
h:=h/2.0; sumt:=sumt+sumu; tprev[k]:=t[k]; uprev[k]:=u[k];
i:=i+1; n:=n+n; goto estimate;
exit: end havie;
```

LAGRANGE INTERPOLATION

real procedure Lagrange(x,y,arg,n,m); value arg,m,n;

array x,y; real arg; integer m,n;

comment HUCC LIBRARY PROCEDURE MT01:

AUTHORS: J. BOOTHROYD, G. WALKER.

The array $y[0:n]$ contains function values $y(x)$ at the sample points of $x[0:n]$ which is sorted in ascending order. The procedure finds an approximation to the function at the specified point arg by evaluating an m th order polynomial ($m \leq n$), choosing the appropriate subset of $m+1$ sample points so that these are evenly distributed about arg . If the requested m exceeds n the assignment $m:=n$ occurs.

For $arg < x[0]$ the procedure extrapolates using

$x[0] \dots x[m]$

For $arg > x[n]$ the procedure extrapolates using

$x[n-m] \dots x[n]$

```
real procedure Lagrange(x,y,arg,n,m);
value arg,m,n; array x,y; real arg; integer n,m;
begin integer i,j,min;
    real term,fac,yest; switch SS:=out;
    integer procedure setmin(L); label L;
    begin integer i;
        switch S:=found;
        for i:=0 step 1 until n do
            if arg<x[i] then goto found;
            i:=n;
        found: if arg=x[i] then begin yest:=y[i]; goto L end;
            i:=i-m div 2-1;
            setmin:= if i<0 then 0 else if i+m>n then n-m else i
        end setmin;
        if m>n then m:=n;
        min:= setmin(out);
        fac:=1.0;
        for j:=min+m step -1 until min do
            fac:=fac*(arg-x[j]);
        yest:=0;
        for i:=min+m step -1 until min do
            begin term:=y[i]*fac/(arg-x[i]);
                for j:=min+m step -1 until i+1,i-1 step -1 until min do
                    term:=term/(x[i]-x[j]);
                yest:=yest+term
            end;
    out:   Lagrange:=yest
end Lagrange;
```

AITKEN UNEQUAL INTERVAL INTERPOLATION

```
real procedure a Aitken (x,y,arg,n,m); value arg,n,m;
array x,y; real arg; integer m, n;
```

```
comment HUCC LIBRARY PROCEDURE MT 02;
```

Author : J. Boothroyd.

This procedure performs the same function as MT 01 but is more efficient and economical. The array $y[0:n]$ contains function values $y(x)$ corresponding to the sample points of $x[0:n]$ which is assumed sorted in ascending order.

The procedure estimates the value of the function corresponding to a specified point arg by evaluating an m^{th} order polynomial ($m \leq n$), choosing the appropriate subset of $m+1$ sample points evenly distributed about the value of arg .

If the requested value m exceeds n the assignment $m:=n$ occurs.

For $arg < x[0]$ the procedure extrapolates using $x[0] \dots x[m]$.

For $arg > x[n]$ the procedure extrapolates using $x[n-m] \dots x[n]$;

```
real procedure aitken(x,y,arg,n,m); value arg,m,n; array x,y; real arg; integer m,n;
begin integer i,j,mless1; real fi,zi; real array z,f[0:m]; switch s:=out;
integer procedure setmin(L); label L;
begin integer i; switch s:=found;
for i:=0 step 1 until n do if arg<x[i] then goto found;
i:=n;
if arg=x[i] then begin f[m]:=y[i]; goto L end;
i:=i-m div 2-1;
setmin:= if i<0 then 0 else if i+m>n then n-m else i
end setmin;
if m>n then m:=n; j:=setmin(out);
for i:=0 step 1 until m do begin z[i]:=arg-x[j]; f[i]:=y[j]; j:=j+1 end; ]
```

found:

```
mless1:=m-1;
for i:= 0 step 1 until mless1 do
begin fi:=f[i]; zi:=z[i];
for j:= i+1 step 1 until m do
f[j]:= fi+zi*(f[j]-fi)/(zi-z[j])
```

end;

out: aitken:=f[m]

end;

AITKEN EQUAL INTERVAL INTERPOLATION

```
real procedure equipol (xbase,y,arg,n,m,h); value xbase,arg,m,n,h;
real xbase,arg,h; array y; integer m,n;
```

comment HUCC LIBRARY PROCEDURE MT 03

Author : J. Boothroyd.

An equal interval version of MT 02. Array y[0:n] contains values of a function $y(x)$ corresponding to the equal interval values of the argument $x_{base}, x_{base} + h, x_{base} + 2h \dots \dots, x_{base} + nh$. The procedure estimates the value of the function for some specified value arg by evaluating an m^{th} order polynomial ($m \leq n$), choosing the appropriate subset of $m+1$ sample points evenly distributed about the value arg . If m exceeds n the assignment $m:=n$ occurs.

For $arg < x_{base}$ the procedure extrapolates using $y[0] \dots y[m]$.

For $arg > x_{base}+nh$ the procedure extrapolates using $y[n-m] \dots y[n]$;

AITKEN Nth ORDER INTERPOLATION

```
real procedure ait(z,f,arg,n); value arg,n; integer n;
real arg; array z,f;
```

comment HUCC LIBRARY PROCEDURE MT 04

Author : J. Boothroyd

An abridged version of MT 02. Array f[0:n] contains values of a function corresponding to the values of the argument in z[0:n] which may be in any order. The procedure estimates the value of $f(z)$ corresponding to some specified value of $z=arg$ by evaluating an n^{th} order polynomial, using all the values of z,f[0:n];

```

real procedure equipol(xbase,y,arg,n,m,h); value xbase,arg,n,m,h; real xbase,arg,h; array y; integer m,n;
begin integer i,j,mless1; real jh,fi; array f[0:m];
if m>n then m:=n; i:=entier((arg-xbase)/h)-m div 2;
j:= if i<0 then 0 else if i+m>n then n-m else i;
for i:= 0 step 1 until m do f[i]:=y[i+j];
arg:=arg-j*h;
mless1:=m-1;
for i:=0 step 1 until mless1 do
begin fi:=f[i]; jh:=h;
for j:=i+1 step 1 until m do
begin f[j]:=fi+arg*(f[j]-fi)/jh; jh:=jh+h end ;
arg:=arg-h
end;
equipol:=f[m]
end equipol;

```

```

real procedure ait(z,f,arg,n); value arg,n; integer n; real arg; array z,f;
begin integer i,j,nless1; real fi,zi,u; nless1:=n-1;
for i:=0 step 1 until nless1 do
begin fi:=f[i]; zi:=z[i]; u:=arg-zi;
for j:=i+1 step 1 until n do
f[j]:=fi+u*(f[j]-fi)/(z[j]-zi)
end;
ait:=f[n]
end;

```

ESTIMATION OF DERIVATIVE OF GIVEN FUNCTION - UNEQUAL INTERVALS

```
real procedure dydx(x,y,arg,n,m,est); value arg,m,n;  
array x,y; real est,arg; integer m,n;
```

comment HUCC LIBRARY PROCEDURE MT 05

Author : J. Boothroyd.

Array $y[0:n]$ contains values of a function $y(x)$ corresponding to the sample values of the argument in $x[0:n]$ which is assumed sorted in ascending order. The procedure estimates the value of the derivative of $y(x)$ at some specified value of $x=arg$ by evaluating the derivative of an m^{th} order polynomial ($m \leq n$), choosing the appropriate subset of $m+1$ sample points evenly distributed about the value of arg .

If $m > n$ the assignment $m:=n$ occurs.

For $arg < x[0]$ the procedure extrapolates using $x[0] \dots x[m]$

For $arg > x[m]$ the procedure extrapolates using $x[n-m] \dots x[n]$.

The output parameter est delivers an estimation of $y(arg)$ yielding a value which is the same as would be obtained from the use of MT 02.

In so far as derivative estimation is a numerical process susceptible to large errors this procedure should be used with caution using values of m not exceeding 5;

```
real procedure dydx(x,y,arg,n,m,est); value arg,m,n; array x,y; real arg,est; integer m,n;
begin integer i,j,mless1; real fi,zi,diffi,zizj,fjfi; real array z,f,diff[0:m];
integer procedure setmin;
begin integer i; switch s:=found;
for i:=0 step 1 until n do if arg<x[i] then goto found;
i:=n;
i:=i-m div 2-1;
setmin:= if i<0 then 0 else if i+m>n then n-m else i
end setmin;
if m>n then m:=n; j:=setmin;
for i:=0 step 1 until m do begin z[i]:=arg-x[j]; f[i]:=y[j]; diff[i]:=0.0; j:=j+1 end;

mless1:=m-1;
for i:= 0 step 1 until mless1 do
begin fi:=f[i]; zi:=z[i]; diffi:=diff[i];
for j:= i+1 step 1 until m do
begin zizj:=zi-z[j]; fjfi:= f[j]-fi;
diff[j]:=diffi+(fjfi+zi*(diff[j]-diffi))/zizj;
f[j]:= fi+zi*fjfi/zizj
end
end;
est:=f[m];
dydx:=diff[m]
end dydx;
```

ESTIMATION OF DERIVATIVE OF GIVEN FUNCTION EQUAL INTERVALS

```
real procedure equidydx (xbase,y,arg,n,m,h,est);value xbase,arg,m,n,h;
real xbase,arg,est,h; array y; integer m,n;
```

comment HUCC LIBRARY PROCEDURE MT 06

Author : J. Boothroyd.

An equal interval version of MT 05. Array $y[0:n]$ contains values of a function $y(x)$ corresponding to the sample values of the argument $xbase, xbase+h, \dots, xbase+nh$. The procedure estimates the value of the derivative of the function for some specified value $x=arg$ by evaluating the derivative of an m^{th} order polynomial ($m \leq n$) choosing the appropriate subset of $m+1$ sample points evenly distributed about arg . If $m > n$ the assignment $m:=n$ occurs.

If $arg < xbase$ the procedure extrapolates using $y[0] \dots y[m]$

If $arg > xbase+nh$ the procedure extrapolates using $y[n-m] \dots y[n]$.

The output parameter est delivers an estimation of the function value $y(arg)$ and yields a value which is the same as would be obtained from MT 03.

This procedure should be used with discretion and for values of m not exceeding 5;

```
real procedure equidydx(xbase,y,arg,n,m,h,est); value xbase,arg,m,n,h; real xbase,arg,est,h; array y; integer m,n;
begin integer i,j,mless1; real jh,fi,diffi,fjfi; array f,diff[0:m];
  if m>n then m:=n; i:=entier((arg-xbase)/h)-m div 2;
  j:= if i<0 then 0 else if i+m>n then n-m else i;
  for i:= 0 step 1 until m do begin f[i]:=y[i+j]; diff[i]:=0.0 end;
  arg:=arg-j*h;
  mless1:=m-1;
  for i:=0 step 1 until mless1 do
    begin fi:=f[i]; jh:=h; diffi:=diff[i];
      for j:=i+1 step 1 until m do
        begin fjfi:=f[j]-fi;
          diff[j]:=diffi+(fjfi+arg*(diff[j]-diffi))/jh;
          f[j]:=fi+arg*fjfi/jh; jh:=jh+h
        end;
      arg:=arg-h
    end;
  est:=f[m];
  equidydx:=diff[m]
end equidydx;
```

REARRANGE ELEMENTS OF VECTOR

procedure PERMB(b,r,n); value n; real array b; integer array r;
integer n;

comment HUCC LIBRARY PROCEDURE NC01 :

AUTHOR J. BOOTHROYD :

A procedure which rearranges the elements of b[1:n] so that
b[i] := b[r[i]] i = 1,2,....n. The procedure should be used
after solving Ax = b by the use of LEO3 and LEO4 in those cases
where further processing of the solutions x_i requires x_i to
occupy b[i] i = 1,2,....n;

```
procedure PERMB(b,r,n); value n; real array b; integer array r; integer n;
comment rearranges the elements of b[1:n] so that b[i]:=b[r[i]],i=1,2,...,n;
begin integer i,k; real w; switch S:= L;
  for i:= n step -1 until 2 do
    begin k:= r[i];
      L: if k then
        begin if k>i then begin k:=r[k]; goto L end;
          w:= b[i]; b[i]:= b[k]; b[k]:= w
        end
      end
    end
  end PERMB;
```

PERMUTE ROWS OR COLUMNS OF MATRIX

comment mxperm(a,b,j,k,s,d,n,p); value n; real a,b; integer array s,d;
integer j,k,n,p;

comment HUCC LIBRARY PROCEDURE NC02:

AUTHOR: J. BOOTHROYD

A procedure using Jensen's device which exchanges rows or columns of a matrix to achieve a rearrangement specified by the permutation vectors s,d[1:n]. Elements of s specify the original source locations while elements of d specify the destination locations.

Normally a and b will be called as subscripted variables of the same array. The parameters j,k nominate the subscripts of the dimension affected by the permutation, p is Jensen's parameter. As an example of the use of this procedure suppose r,c[1:n] to contain the row and column subscripts of the successive matrix pivots following an in-situ matrix inversion. The two calls

mxperm(a[j,p],a[k,p],j,k,r,c,n,p)
and mxperm(a[p,j],a[p,k],j,k,c,r,n,p)

will respectively perform the required rearrangement of rows and columns.

```
procedure mxperm(a,b,j,k,s,d,n,p);  value n;  real a,b;  integer array s,d;  integer j,k,n,p;
begin integer array tag,loc[1:n];  integer i,t,tagj;  real w;
    comment set up initial vector tag number and address arrays;
    for i:=1 step 1 until n do tag[i]:=loc[i]:=i;
    comment start permutation;
    for i:=1 step 1 until n do
        begin t:=s[i];  j:=loc[t];  k:=d[i];
        if j≠ k then
            begin for p:=1 step 1 until n do
                begin w:=a;  a:=b;  b:=w end;
                tag[j]:=tag[k];  tag[k]:=t;  tagj:=tag[j];
                loc[t]:=loc[tagj];  loc[tagj]:=j
            end jk conditional
        end i loop
    end mxperm;
```

PERMUTATION OF ELEMENTS OF A VECTOR

procedure vectorperm(m,d,n,mode,endperm); value n,mode; integer n,mode
array m; integer array d;

comment HUCC LIBRARY PROCEDURE NC03:

AUTHOR: J. BOOTHROYD

A procedure for generating, at each entry, a permutation of the elements $m[1], m[2], \dots, m[n]$ of $m[1:n]$. $(n-1)!$

successive entries generate all $n!$ permutations.

The permutation is controlled by a variable radix counter, the array $d[2:n]$, with digit positions $d[2], d[3] \dots d[n]$ where the subscript indicates the radix value. Starting with $d=(0,0,0 \dots 0)$ one is added to the counter at each entry to the procedure. One and only one element of d increases each time and all element positions below this are reset to zero.

Denoting by k the subscript position of the counter element which increases the permutation rules are:-

(k odd) or (k even and $d[k] < 2$) interchange $m[k], m[k-1]$

k even for $2 < d[k] < k$ interchange $m[k], m[k-d[k]]$

A call of vectorperm initialises d to $(0,0,0,0,\dots,0)$ with mode=1 preparatory to further calls with mode=2. At the $n!$ th call the counter d resets to zero and the procedure exits to the label endperm;

```
procedure vectorperm(m,d,n,mode,endperm);  value n,mode;  integer n,mode;
integer array d;  array m;  label endperm;
begin integer k,j;  real temp;  switch s:=set,run,swap,exit;
  goto s[mode];
set:   for k:=2 step 1 until n do d[k]:=0;  goto exit;
run:   j:=-1;
        for k:=2 step 1 until n do
          begin if d[k]#k-1 then goto swap;  d[k]:=0;  j:=-j end;
  goto endperm;
swap:  d[k]:=d[k]+1;
        if j#1 then j:= if 2>d[k] then 1 else d[k];
        temp:=m[k];  m[k]:=m[k-j];  m[k-j]:=temp;
exit: end vectorperm;
```

PERMUTE ROWS OR COLUMNS OF MATRIX

procedure mxperm1(a,b,j,k,s,d,n,p); value n; real a,b; integer array s,d;
integer j,k,n,p;

comment HUCC LIBRARY PROCEDURE NC04:

AUTHOR: J. BOOTHROYD

A procedure which is the same as, but avoids the use of the local arrays of, NC02. It is, therefore, more economical of space but slower in operation. The identifiers of NC02 and NC04 agree over the first six characters, the parameter specifications are identical and NC04 is thus a direct replacement for NC02.

```
procedure mxperm1(a,b,j,k,s,d,n,p);  value n;  real a,b;  integer array s,d;  integer j,k,n,p;
begin integer i; switch ss:= scan; real w;
  for i:= n step -1 until 2 do
    begin j:= s[i]; k:= d[i];
    scan: if j#k then
      begin for p:= i+1 step 1 until n do
        if j=d[p] then begin j:= s[p]; goto scan end;
        for p:= 1 step 1 until n do begin w:= a; a:= b; b:= w end
      end
    end
  endmxperm;
```

PRE OR POST MULTIPLY MATRIX BY PERMUTED IDENTITY MATRIX

```
procedure permx (a,b,j,k,r,n,p,inv); value n,inv;
real a,b; integer j,k,n,p,inv; integer array r;
```

comment HUCC LIBRARY PROCEDURE NC 05

Author : J. Boothroyd

A procedure using Jensen's device whereby a matrix $A[1:n,1:n]$ may be either pre or post multiplied by either a permuted identity matrix Ir or by the inverse of Ir . All row (or column) exchanges are performed on A without the use of an array of the same size. The inverse of Ir is its transpose and both Ir and Ir^{-1} may be defined by a permutation vector $r[1:n]$ as follows :-

$$Ir[r[i],j] = \begin{cases} 0 & i \neq j \\ 1 & i = j \end{cases} \quad Ir^{-1}[i,r[j]] = \begin{cases} 0 & i \neq j \\ 1 & i = j \end{cases}$$

for $i=1,2,\dots,n$, $j=1,2,\dots,n$.

Values of the parameter $inv \geq 0$ accomplish multiplication by Ir . Otherwise, for $inv < 0$ ($inv = -1$ is appropriate) multiplication is by Ir^{-1} .

The following calls illustrate how to achieve pre or post multiplication

Premultiplication by Ir `permx(A[j,p],A[k,p],j,k,r,n,p,0)`

Post multiplication by Ir^{-1} `permx(A[p,j],A[p,k],j,k,r,n,p,-1);`

```
procedure permx(a,b,j,k,r,n,p,inv); value n,inv; real a,b; integer j,k,n,p,inv; integer array r;
begin integer i; real w;
procedure itori;
begin switch s:=L;
  for i:=n step -1 until 2 do
    begin k:=r[i]; j:=i;
    L: if j≠k then
      begin if k>j then begin k:=r[k]; goto L end end;
      for p:=1 step 1 until n do
        begin w:=a; a:=b; b:=w end
    end
  end itori;
procedure ritoi;
begin switch s:=scan;
  for i:=n step -1 until 2 do
    begin j:=i; k:=r[i];
    scan: if j≠k then
      begin for p:=i+1 step 1 until n do
        if j=r[p] then begin j:=p; goto scan end;
        for p:=1 step 1 until n do
          begin w:=a; a:=b; b:=w end
      end
    end
  end ritoi;
  if inv>0 then itori else ritoi
end permx;
```

NC 06

INVERSEPERMUTATION OF INTEGER VECTOR

procedure inversepermb(P,n); value n; integer n; integer array P;

comment HUCC LIBRARY PROCEDURE NC 06

Author : J. Boothroyd

Rearranges the natural number integer elements of $P[1:n]$
to effect an inversepermutation

	1	2	3	4	5	6	7	
i.e. Given	P	6	2	1	3	7	5	4

the procedure produces P

3	2	4	7	6	1	5
---	---	---	---	---	---	---

This process is best described by :-

for i:=1 step 1 until n do Q[P[i]]:=i;

for i:=1 step 1 until n do P[i]:=Q[i]

though NC 06 accomplishes the transformation without the use of the auxiliary array Q;

```
procedure inverseperm(P,n); value n; integer n; integer array P;
begin integer i,j; switch s:=next;
  for i:=1 step 1 until n do P[i]:=-P[i];
  for n:=n step -1 until 1 do
    begin j:=n;
    next: i:=P[j]; if i>0 then begin j:=i; goto next end;
          P[j]:=P[-i]; P[-i]:=n
    end
end inverseperm;
```

CRITICAL PATH SCHEDULING

```
procedure CRITICALPATH(n,I,J,DIJ,ES,LS,EF,LF,TF,FF,i,out);  
integer i,n; integer array I,J,DIJ,ES,LS,EF,LF,TF,FF; label out;  
comment HUCC LIBRARY PROCEDURE OP01:  
AUTHOR ACM 74 (improved):
```

Given the total number of jobs n of a project, the pair I[k],J[k] representing the kth job as a directed line joining event I[k] to event J[k] ($I[k] < J[k], k=1,2,\dots,n$) and a duration vector DIJ, the procedure determines the earliest starting time ES[k], latest starting time LS[k], earliest completion time EF[k], latest completion time LF[k] total float TF[k] and free float FF[k]. I[1] must be 1 and the I[k],J[k] must be in ascending order with no values in the sequence 1 to n-1 omitted from the I[k] sequence. The critical path is the chain of all jobs with zero total float. The following tests are included:

- (a) That $I[k] < J[k]$,
- (b) All I[k] are in ascending sequence
- (c) No I[k] is missing.

Failure to meet these requirements causes exit to the label parameter out. This exit is so arranged that by the use of an actual parameter OUT[i], for example, the respective failure paths are (a) OUT[1], (b) OUT[2], (c) OUT[3];

```
procedure CRITICALPATH(n,I,J,DIJ,ES,LS,EF,LF,TF,FF,i,out);
  integer i,n;  integer array I,J,DIJ,ES,LS,EF,LF,TF,FF;  label out;
  begin integer k,index,m,tiik,tejk,dijk;  integer array ti,te[1:n];
    index:=1;
    for k:=1 step 1 until n do
      begin if I[k]>J[k] then begin i:=1;  goto out end;
         if I[k]<index then begin i:=2;  goto out end;
         if I[k]>index and I[k]≠ index+1 then begin i:=3;  goto out end;
         if I[k]=index+1 then index:=I[k]
      end;
    for k:=1 step 1 until n do begin ti[k]:=0;  te[k]:=9999 end;
    for k:=1 step 1 until n do
      begin m:=ti[I[k]]+DIJ[k];
        if ti[J[k]]<m then ti[J[k]]:=m
      end ti;
    te[J[n]]:=ti[J[n]];
    for k:=n step -1 until 1 do
      begin m:=te[J[k]]-DIJ[k];
        if te[I[k]]>m then te[I[k]]:=m
      end te;
    for k:=1 step 1 until n do
      begin tiik:=ti[I[k]];  tejk:=te[J[k]];  dijk:=DIJ[k];
        ES[k]:=tiik;  EF[k]:=tiik+dijk;
        LS[k]:=tejk-dijk;  LF[k]:=tejk;
        TF[k]:=tejk-tiik-dijk;  FF[k]:=ti[J[k]]-tiik-dijk
      end
  end CRITICALPATH;
```

REAL RANDOM NUMBER GENERATOR

real procedure random;

Comment HUCC LIBRARY PROCEDURE SS01:

AUTHOR J. BOOTHROYD :

A procedure to generate a stream of uniformly distributed real random numbers in the interval $0 \leq \text{random} \leq 1$. Random integers xr are formed by the operation.

$$xr := ((2^{20} - 3) \times r - 1) \text{ modulo } 2^{38}$$

and each real value random is formed from a corresponding value of xr by discarding the least significant nine bits. The resulting 30 bit truncated integer is then treated as a fraction f in the interval $0 \leq f < 1$ and an appropriate adjustment made in the integer to real conversion. The period of the sequence is thus 2^{38} though, by virtue of the truncation, each value of random will repeat 512 times in a complete cycle.

To obtain a uniform distribution in the interval $[a, b]$ use the expression $a + (b-a) * \text{random}$. To obtain a symmetrical triangular distribution with mean =1 in the range $0 \leq y < 2$ use the statement $y := \text{random} + \text{random}$.

Note that random is a procedure without parameters. The initial value of random is thus indeterminate. In circumstances where this might be inconvenient, as for instance in successive program testing runs, it would be preferable to use LIBRARY SS02;

SS02

Real Random Number Generator

real procedure randG2 (x, start); value x, start; integer x, start;
comment HUCC LIBRARY PROCEDURE SS02:

AUTHOR J. BOOTHROYD :

A procedure having the same function as LG1 but with facilities for determining the starting value of the stream of random numbers. A call of the procedure with start = 0 and x = X will cause the first number to be generated from

$$((2^{20} - 3) X - 1) \text{ modulo } 2^{38}$$

Subsequent calls of the form randG2 (0, 1) will generate numbers from the value of randG2 at the immediately previous call;

comment LG1;

```
real procedure random;
begin own integer xr; integer three8,cc; real f; switch S:=out;
cc:=274877906432; three8:=38;
elliott(0,2,xr,0,0,1,0);
elliott(5,0,1,0,0,5,xr);
elliott(5,0,19,0,0,4,xr);
elliott(5,0,18,0,5,7,0);
elliott(2,0,xr,0,0,3,cc);
elliott(6,5,4096,0,0,5,three8);
elliott(2,0,f,0,4,3,out);
out: random :=f
end random;
```

SS01

comment LG2; real procedure rand G2(x,start); value x,start; integer x,start;

```
begin own integer xr,cc, three8; real f; switch S:=out;
if start =0 then begin xr:=x; three8:=38; cc:=274877906432 end;
elliott(0,2,xr,0,0,1,0);
elliott(5,0,1,0,0,5,xr);
elliott(5,0,19,0,0,4,xr);
elliott(5,0,18,0,5,7,0);
elliott(2,0,xr,0,0,3,cc);
elliott(6,5,4096,0,0,5,three8);
elliott(2,0,f,0,4,3,out);
out: rand G2:=f
end randG2;
```

SS02

RANDOM NUMBER GENERATOR - POISSON DISTRIBUTION

integer procedure poisson(lambda); value lambda; real lambda;

comment HUCC LIBRARY PROCEDURE SS03:

AUTHOR J. BOOTHROYD :

A procedure which generates integers x having the
Poisson probability distribution

$$P(x) = \frac{e^{-\lambda} \lambda^x}{x!}$$

The numbers are generated from a distribution uniform
in $0 \leq y < 1$ by forming x terms of a continued product
satisfying

$$y_1 y_2 y_3 \dots y_x < e^{-\lambda}$$

For large values of λ the method is therefore not
particularly efficient, and it may be preferable to use
the fact that for large λ the Poisson distribution
tends to a normal distribution and employ SS04;

```
integer procedure poisson(lambda); value lambda; real lambda;
begin own integer xr; integer three8,n; real elambda,try; switch s:= L;
    three8 := 38; n:=-1; try:= 1.0; elambda:=exp(-lambda);
    L: elliott(0,2,xr,0,0,1,0);
    elliott(5,0,1,0,0,5,xr);
    elliott(5,0,19,0,0,4,xr);
    elliott(5,0,18,0,5,7,0);
    elliott(1,6,xr,0,5,4,29);
    elliott(5,5,9,0,6,5,4096);
    elliott(0,5,three8,0,6,3,try);
    elliott(2,0,try,0,6,2,elambda);
    elliott(2,2,n,0,4,1,L);
    elliott(7,3,4,1,4,3,1);
    poisson:= n
endpoisson;
```

RANDOM NUMBER GENERATOR - STANDARD NORMAL DISTRIBUTION.

real procedure normal;

comment HUCC LIBRARY PROCEDURE SS04:

AUTHOR J. BOOTHROYD :

Generates random numbers with a standard normal distribution (mean zero and unit standard deviation). The numbers are generated from a distribution uniform in $0 \leq y < 1$ by summing twelve values of y and adjusting the mean by subtracting 6.0

$$\text{normal} := \sum_{i=1}^{12} y_i - 6.0$$

To generate numbers having a mean u and standard deviation σ use the transformation

$\text{musigma} := \sigma * \text{normal} + u$;

```
real procedure normal;
begin own integer xr; integer three8,count; real sum; switch s:= L;
    three8:=38; count:= -11; sum:= -6.0;
    L: elliott(0,2,xr,0,0,1,0);
    elliott(5,0,1,0,0,5,xr);
    elliott(5,0,19,0,0,4,xr);
    elliott(5,0,18,0,5,7,0);
    elliott(1,6,xr,0,5,4,29);
    elliott(5,5,9,0,6,5,4096);
    elliott(0,5,three8,0,6,0,sum);
    elliott(2,0,sum,0,3,2,count);
    elliott(4,1,L,0,0,0,0);
    elliott(7,3,4,1,4,3,1);
    normal:= sum
endnormal;
```

SS05

GENERATE LARGE INTEGER RANDOM NUMBERS

integer procedure bigrn;

comment HUCC LIBRARY PROCEDURE SS05:

AUTHOR J. BOOTHROYD :

A procedure to generate, equiprobably, integers in the range $0 \leq \text{bigrn} < 2^{38}-1$. There is thus a 0.5 probability that $\text{bigrn} > 2^{37}$. The numbers are formed by the recurrence

$$\text{xr} := ((2^{20}-3)\text{xr}-1) \bmod 2^{38}.$$

For repeatable program testing purposes it may be preferable to use SS06;

SS06

GENERATE LARGE INTEGER RANDOM NUMBERS

integer procedure bigrnx(x,start); value x,start; integer x,start;

comment HUCC LIBRARY PROCEDURE SS06:

AUTHOR J. BOOTHROYD :

Performs the same function as SS05 but with facilities for initialising the starting value of the stream of random numbers. With start=0 the value of bigrnx is generated from the value supplied by the parameter x. For start=1 the value of x is ignored and operation of SS06 is the same as SS05;

SS07

GENERATE SMALL INTEGER RANDOM NUMBERS

integer procedure rn(n); value n; integer n;

comment HUCC LIBRARY PROCEDURE SS07:

AUTHOR J. BOOTHROYD :

A procedure to generate, equiprobably, integers in the range $0 \leq rn < 2^n$. The basic generation scheme is the recurrence

$$xr := ((2^{20}-3) \times r - 1) \bmod 2^{38}$$

The value given by the top n bits of each generated xr is assigned to rn . Where repeatable program testing results are required, it may be preferable to use SS08;

SS08

GENERATE SMALL INTEGER RANDOM NUMBERS

integer procedure rn(x,start,n); value x,start,n; integer x,start,n;

comment HUCC LIBRARY PROCEDURE SS08:

AUTHOR J. BOOTHROYD :

A procedure performing the same function as SS07 but with facilities for determining the starting value of the stream of random numbers. With $start=0$ the value of the first generated number is determined by the value of x . For $start=1$, x is ignored, and the procedure is identical with SS07;

SS05

```
integer procedure bigrn;
begin own integer xr;
    elliott(0,2,xr,0,0,1,0);
    elliott(5,0,1,0,0,5,xr);
    elliott(5,0,19,0,0,4,xr);
    elliott(5,0,18,0,5,7,0);
    elliott(2,0,xr,0,0,0,0);
    elliott(7,3,4,1,4,3,1);
    bigrn:= xr
endbigrn;
```

SS06

```
integer procedure bigrnx(x,start); value x,start; integer x,start;
begin own integer xr;
    if start=0 then xr:= x;
    elliott(0,2,xr,0,0,1,0);
    elliott(5,0,1,0,0,5,xr);
    elliott(5,0,19,0,0,4,xr);
    elliott(5,0,18,0,5,7,0);
    elliott(2,0,xr,0,0,0,0);
    elliott(7,3,4,1,4,3,1);
    bigrnx:= xr
endbigrnx;
```

SS07

```
integer procedure rn(n); value n; integer n;
begin own integer xr;
    elliott(0,2,xr,0,0,1,0);
    elliott(5,0,1,0,0,5,xr);
    elliott(5,0,19,0,0,4,xr);
    elliott(5,0,18,0,5,7,0);
    elliott(1,6,xr,0,6,7,n);
    elliott(5,4,0,0,2,0,n);
    elliott(7,3,4,1,4,3,1);
    rn:=n
endrn;
```

SS08

```
integer procedure rnx(x,start,n); value x,start,n; integer x,start,n;
begin own integer xr;
    if start=0 then xr:= x;
    elliott(0,2,xr,0,0,1,0);
    elliott(5,0,1,0,0,5,xr);
    elliott(5,0,19,0,0,4,xr);
    elliott(5,0,18,0,5,7,0);
    elliott(1,6,xr,0,6,7,n);
    elliott(5,4,0,0,2,0,n);
    elliott(7,3,4,1,4,3,1);
    rnx:=n
endrn;
```

OUTPUT BINARY VALUE

```
procedure binary(i,g);
value i,g; integer i,g;
comment HUCC LIBRARY PROCEDURE ZM01:
AUTHOR: W. G. WARNE:
```

Prints on current device the binary value of the integer i in groups of g digits separated by a space starting from the most significant binary digit. No character is output before the first digit. If g = 0 no spaces are output to cause grouping.

OUTPUT OCTAL VALUE

```
procedure octal(i,g);
value i,g; integer i,g;
comment HUCC LIBRARY PROCEDURE ZM02:
AUTHOR: W. G. WARNE:
```

Outputs on current device the octal value of the integer i in groups of g octal digits, separated by a space, starting from the most significant octal digit. No character is output before the first octal digit.

If g = 0 the integer is output as a single group.

ZM01

```
procedure binary(i)in groups of:(g);
value i,g; integer i,g;
begin integer n,gc;
switch sw := neg, cont;
gc:=0;
for n:=1 step 1 until 39 do
begin elliot(3,0,i,0,4,1,neg);
print f0?;
goto cont;
neg:
print f1?;
gc:=gc+1;
if gc=g then begin print ffs??; gc:=0 end;
elliot(3,0,i,0,5,5,1);
elliot(7,3,4,1,4,3,1);
elliot(2,0,i,0,0,0,0);
end
end binary;
```

```
procedure octal(i) in groups of:(g);
value i,g;  integer i,g;
begin      integer n,gc,temp,seven;
           gc:=0;
           seven:=7;
           for n:=36 step -3 until 0 do
               begin elliot(3,0,i,0,6,7,n);
                   elliot(5,1,0,0,0,3,seven);
                   elliot(2,0,temp,0,2,2,gc);
                   print sameline,special(1),digits(1),temp;
                   if gc=g then begin print ffs??;  gc:=0 end
               end
           end octal;
```

SENSE NUMBER GENERATOR KEY

boolean procedure ng(i); value i; integer i;

comment HUCC LIBRARY PROCEDURE ZM03:

AUTHOR ELLIOTT AUTOMATION:

This procedure, when called as ng(x) where x is an integer, $1 \leq x \leq 39$, takes the value true if the key numbered x on the word generator is depressed at the time of call. If key x is not depressed, the value of ng is false. A typical call would then be

if ng(x) then ;

INPUT NEWLINE STRINGS

procedure charead(array); integer array array;

comment HUCC LIBRARY PROCEDURE ZM04:

AUTHOR P. W. FORD :

A call of charead causes the input tape to be searched for the first occurrence of a non newline non blank character immediately following a newline \backslash , i.e. blank tape and a succession of newlines is ignored. This first character and all subsequent non blank characters up to but not including the next newline are stored in the integer array array in a form consistent with subsequent use of the procedure outstring. The character string is stored starting in array[1] and adequate bounds for this array must be declared (see 503 TECH MANUAL 2.1.3.2 section 2.6). There is no protection on the bounds of array and the input of a string longer than can be accommodated is likely to be discovered by a subscript overflow failure when outstring is called;

ZM03

```
boolean procedure ng(i); value i; integer i;
begin switch s:=1;
    elliott(0,2,0,0,0,0,0);
    elliott(0,0,i,1,5,4,8191);
    elliott(1,6,i,0,7,0,0);
    elliott(2,3,i,0,4,3,1);
    l: ng:=i ≠ 0
end ng(i);
```

ZM04

```
procedure charead(array);
integer array array;
begin integer i,j,k,l,m;
    switch ss:=search,next,read,store,shift,test,out,back,on;
    m:=2;
    l:=64;
    elliott(0,2,0,0,5,5,38);
    elliott(2,0,k,0,4,3,search);
    search: elliott(0,6,0,0,7,1,0);
    elliott(0,5,m,0,4,2,next);
    elliott(4,0,search,0,0,0,0);
    next:   elliott(0,2,0,0,0,7,array);
    elliott(1,6,j,0,0,0,0);
    read:   elliott(7,1,0,0,4,2,read);
    elliott(0,5,m,0,4,2,read);
    elliott(0,4,m,0,5,0,7);
    back:   elliott(3,0,1,0,0,0,0);
    shift:  elliott(5,4,7,0,1,6,i);
    elliott(7,1,0,0,0,5,m);
    elliott(4,2,out,0,0,4,m);
    elliott(5,0,7,0,3,0,i);
    elliott(4,3,store,0,4,0,shift);
    store:  elliott(2,2,j,1,1,6,0);
    elliott(4,0,back,0,0,0,0);
    out:   elliott(3,0,i,0,0,0,0);
    test:   elliott(4,3,on,0,5,4,7);
    elliott(4,0,test,0,0,0,0);
    on:    elliott(0,4,k,0,0,0,0);
    elliott(2,2,j,1,1,6,0)
end charead;
```

ZM05
ZM06

Algol A.D.T. Read and Punch Binary

Two procedures designed to meet the requirements of programs using paper tape as temporary storage so that partial results output by one program are used solely as subsequent input data for the same or another program. In these cases the most appropriate information format is pure binary and further efficiency has been obtained by taking advantage of the interrupt facilities of the 503 so that input and output takes place simultaneously with computation or the use of peripheral devices other than those used by ZM05 or ZM06, which may operate together.

Both procedures use a boolean parameter flag to indicate when a transfer initiated by a call of ZM05 or ZM06 is complete. Programs using these procedures should not access the data transfer locations unless flag = true.

Realignment of a tape, output by ZM05, at subsequent use by ZM06 may be accomplished in several ways:-

- (a) By arranging to output a correctly terminated number which is read in and checked before a call of ZM06
- (b) By arranging to output an identifiable character and, on reinput, to execute a buffer search for the same character, or more simply but without any check, execute a call of the procedure advance which ignores blank tape and erase characters.

ZM05

procedure out(a,n,flag); value n; integer n; array a; boolean flag;

comment HUCC LIBRARY PROCEDURE ZM05

AUTHOR : W.G. WARNE

A procedure which outputs, in binary, on punch 1, the first n(>1) locations of array a, i.e. by rows. The parameter flag is assigned the value true when the transfer is complete.

Total time is approximately 60n milliseconds of which 99% is available for simultaneous computation;

ZM06

procedure in(a,n,flag); value n; integer n; array a; boolean flag;

comment HUCC LIBRARY PROCEDURE ZM06

AUTHOR : W.G. WARNE

A procedure which reads, from reader 1, tapes previously output by ZM05. (see the notes above concerning realignment of tape). Data from the tape is input to the first $n(>1)$ locations of array a. Total time is approximately $6n$ milliseconds of which 90% is available for simultaneous computation;

```

procedure out(a,n,flag);
value n;
array a;
boolean flag;
integer n;
begin      own integer aa,i,ab,temp,hold,x,mask,imask;
            switch s:= fin,L1,L2,L5;
            mask := 127;
            x := 32;
            imask := 223;
            flag := false;
            aa := address(a);
            ab := aa + n - 1;
            goto L1;
L2:        i := i - 262144;
            elliot(3,0,7886,0,0,3,imask);
            elliot(0,4,x,0,2,0,7886);
            elliot(6,7,7886,0,7,2,0);
            elliot(0,2,i,0,5,5,20);
            elliot(2,6,7895,0,1,0,7892);
            elliot(2,0,hold,0,2,6,7893);
            elliot(2,6,7894,0,3,0,fin);
            elliot(1,6,4,0,0,2,0);
            elliot(5,5,3,0,2,0,7896);
            for aa := aa step 1 until ab do
            begin      elliot(0,6,aa,1,0,4,0);
                        elliot(2,0,temp,0,0,0,0);
                        for i := 35 step -7 until 0 do
                        begin      elliot(3,0,temp,0,6,7,i);
                                    elliot(5,1,0,0,0,3,mask);
                                    elliot(2,0,x,1,7,4,0);
                                    elliot(6,6,7893,0,0,0,0);
L1:        elliot(7,3,i,0,4,0,L2);
            elliot(7,5,7168,0,5,5,2);
            elliot(7,3,x,1,4,3,1);
            elliot(4,1,L5,0,0,0,0);
            end;
            end;
            flag := true;
            elliot(3,0,hold,0,2,0,7892);
            elliot(3,0,imask,0,0,3,7886);
            elliot(2,0,7886,1,7,2,0);
L5:        elliot(6,6,7893,0,0,0,0);
fin: end;

```

```
procedure in (a,n,flag);
value n; integer n;
array a;
boolean flag;
    begin own integer aa,i,ab,temp,hold,five,imask;
        switch s:= fin,L1,L2,L3,L4,L5;
        flag:= false;
        aa:=address(a);
        ab:=aa+n-1;
        five := 5;
        temp := 64;
        imask := 191;
        goto L1;
L2:   i:=i-262144;
        elliott(3,0,7886,0,0,3,imask);
        elliott(0,4,temp,0,2,0,7886);
        elliott(6,7,7886,0,7,2,0);
        elliott(0,2,i,0,5,5,20);
        elliott(2,6,7888,0,1,0,7887);
        elliott(2,0,hold,0,2,6,7889);
        elliott(3,0,five,0,2,1,i);
        elliott(2,6,7890,0,3,0,fin);
        elliott(1,6,4,0,0,2,0);
        elliott(5,5,3,0,2,0,7891);
        elliott(0,6,0,0,4,0,L3);
L4:   elliott(2,6,temp,0,3,0,five);
        elliott(2,1,i,0,6,6,7888);
L1:   elliott(7,3,i,0,4,0,L2);
        elliott(7,5,7168,0,4,1,L5);
        elliott(3,0,temp,0,5,5,7);
L3:   elliott(7,1,0,0,2,0,temp);
        elliott(3,2,i,0,4,1,L5);
        elliott(3,0,temp,0,6,7,aa);
        elliott(2,0,0,0,3,2,aa);
        elliott(0,5,ab,0,4,1,L4);
        flag:=true;
        elliott(3,0,hold,0,2,0,7887);
        elliott(3,0,imask,0,0,3,7886);
        elliott(2,0,7886,1,7,2,0);
L5:   elliott(6,6,7888,0,0,0,0);
fin: end;
```

ZM07

COPY LEGEND FROM DATA TAPE

procedure copy(n); value n; integer n;

comment HUCC LIBRARY PROCEDURE ZM07:

AUTHOR : W.G. WARNE:

Copies characters from reader (1) to punch (1) up to but not including the first character having binary value n where $0 \leq n \leq k27$ in 8-hole mode or $0 \leq n \leq 31$ in 5-hole mode;

ZM08

Read ALGOL BUFFER

integer procedure BUFFER (d); value d; integer d;

comment HUCC LIBRARY PROCEDURE ZM08:

AUTHOR : W.G. WARNE:

Takes the value of the character in buffer(d) where d(1,2 or 3) is the device number. If d is not 1,2 or 3 the message "BUFFER error" is displayed followed by a data wait. Continuation assigns the value zero to BUFFER. The procedure has been tested with Algol 1/3 Tapes 1 and 2 only on a basic machine. There is thus no guarantee that it is operative with the backing store version of Algol 1;

ZM07

```
procedure copy(n);
value n; integer n;
comment Copies characters from reader 1 to punch 1 until a character with
binary value n is met. This character is not copied.
Note :- 0<=n<=127 in 8 hole mode and 0<=n<=31 in 5 hole mode. ;
begin
    integer t;
    switch L := L1,L2;
    L1:   elliott(0,6,0,0,7,1,0);
          elliott(2,0,t,0,0,5,n);
          elliott(4,6,L2,0,6,7,t);
          elliott(7,4,0,0,4,4,L1);
    L2:
end copy;
```

```
integer procedure BUFFER(d);
comment Takes binary value of character in buffer for device d (d=1,2 or 3).;
value d; integer d;
if d < 1 or d > 3 then
  begin print punch(3), ££1?BUFFER error?,wait;
  BUFFER := 0
  end
else begin elliott(3,0,7920,0,2,4,d);
  elliott(0,6,d,1,0,4,79);
  elliott(2,0,d,0,0,0,0);
  BUFFER := d
end BUFFER;
```

PRINT MATRIX

```

procedure matprint(fields,rows,cols,i,j,margin,gap,var);
value fields,rows,cols; integer fields,rows,cols,i,j; real var;
string margin,gap;
comment HUCC LIBRARY PROCEDURE ZPO2:
AUTHOR: J. BOOTHROYD      :

```

A Jensen procedure for printing the elements of either one or two dimensional arrays with facilities for specifying

- (a) fields the number of numbers printed per line
- (b) margin the prefix required at the start of each line
- (c) gap the number of spaces (or other symbols) printed between numbers.

Notes 1. i and j are the Jensen parameters, with ranges $i,1(1)rows,j,1(1)cols$.

2. var must be called as a real variable, $A[i,j]$ for a matrix $A[1:rcws,1:cols]$ or $a[j]$ for a vector $a[1:cols]$. For arrays with other than unit lower bound the appropriate adjustments should be made to the subscript expressions and to the values used for rows and cols. For example, a in the case of TABLE[0:6,0:10] the appropriate parameters would be

```

rows = 7           cols = 11
var = TABLE[i-1,j-1]

```

3. The procedure does not specify the style of number printing. The user has the option of any of the following methods of determining the required format:

- (a) By calling the required setting procedure (aligned(m,n), digits(n)etc) globally.
- (b) By calling the required setting procedure locally in a print statement.

e.g. print aligned(2,3),matprint();

3 (cont.)

(c) By declaring a private format procedure to permit optional setting procedures depending on circumstances

e.g. real procedure format(x); value x; real x;
begin if abs(x)<0.1 then scaled(7) else aligned(5,6)
format:= x
end;

Within matprint this procedure would be used as an actual parameter as in

matprint(5,10,10,i,j,ffs5??,ffs4??,format(A[i,j]))

4. It is the users responsibility to ensure that the total character requirements of the setting procedure, margin and gap are consistent with the requested value of fields and the limit of 120 characters per line.

Examples: To print A[1:20,1:20] with 5 spaces as margin, 3 spaces between numbers and 6 numbers per line.

matprint(6,20,20,i,j,ffs5??,ffs3??,A[i,j])

To print LIST[1:n] with each line starting one tab from the paper margin, 7 numbers per line, 4 spaces between numbers.

matprint(7,1,n,i,j,fft??,ffs4??,LIST[j])

```
procedure matprint(fields,rows,cols,i,j,margin,gap,var);  value fields,rows,cols;
    integer fields,rows,cols,i,j;  real var;  string margin,gap;
begin integer k,line,np;  switch s:=more;
for i:=1 step 1 until rows do
    begin print ffl12??;  j:=0;  np:=cols;
more:   np:=np-fields;  line:=fields;
        if np < 0 then line:=np+fields;
        print margin;  sameline;
        for k:=1 step 1 until line do
            begin j:=j+1;
                if k ≠ 1 then print gap;
                print var
            end;
        print ffl1??;  if np > 0 then goto more
end i
end matprint;
```

CONTROL FLEXOWRITER PAGE PRINTING

```
procedure pageout(string,type,resultline,array,start);  
value type,resultline,start; string string;  
integer type,resultline,start; integer array array;  
comment HUCC LIBRARY PROCEDURE ZP03:  
AUTHOR P. W. FORD :  
:
```

A procedure to control the automatic pageing of results printed on Lamson Paragon Paraflo paper (Form No. 1112). The number of resultlines is internally set at 50 and change to a new page occurs if the next batch of output would cause this limit to be exceeded. Page numbering is provided and facilities are provided for including on each page indicative information such as job identification, result headings, etc. A call of pageout must occur before each print statement or group of print statements;

Parameters

type an integer with permissible values 0,1,2.
type = 0 prepares the procedure by initialising the the line and page counters.
type = 1 for normal page control.
type = 2 initiates an immediate page change and resets the page counter.
resultline an integer specifying the number of lines of output in the immediately following print statement(s). This parameter is ignored, and may be 0, if type = 0 or 2.

Parameters (cont.)

array the identifier of an array containing some job or other indicative string to be output at the head of each newpage on the line above the page number. If this facility is not required (start = 0, see below) use any declared integer array identifier if one exists otherwise declare a minimal array, e.g. A[1:1] for this purpose.

start an integer parameter which controls the outstring facility used for job identification. If start = 0 the outstring facility is ignored, while start ≠ 0 causes the string starting at array[start] to be output at the head of each newpage. Note that start may be an integer constant, and does not have to be an integer variable identifier.

string Array string (including the empty string f?). This is output on each newpage following the line bearing the page number. This facility permits the printing of result headings, etc.

NOTES

STRINGS should NOT contain newline characters.

This procedure does not control line changing. This must be done in the print statements which pageout monitors. The procedure must be initialised once with type = 0, and this is preferably done in the housekeeping.

```
procedure pageout(string,type,resultline,array,start);
  value type,resultline,start;
  string string;
  integer type,resultline,start;
  integer array array;
begin  own integer line,page;
  switch sss:=L1,L2,L3;
  if type=2 then goto L1;
  if type=1 then
    begin if line+resultline>50 then
      L1:begin for line:=line+1 while line<62 do print ff1??
        if type=2 then goto L3;
        goto L2
      end
      else line:=line+resultline
    end;
  if type=0 then
    L3:begin page:=1;
      L2:print ff1??
        if start ≠ 0 then outstring(array,start);
        print ff1t8?page?, sameline, digits(3), page, ff12??, string;
        page:=page+1;
        line:=if type=1 then resultline else 0
    end
end procedure pageout;
```